

## VTT Technical Research Centre of Finland

### Spine Deliverable 3.1 The Spine Model and its documentation

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<b>Spine</b>		
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## Deliverable 3.1 The Spine Model and its documentation

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Description of the related task and the deliverable. Extract from DoA		<p>This work package is central to the whole project as it will design and implement the Spine Model that will optimize integrated energy systems (by minimizing the value of the objective function while respecting constraint functions). The work will be carried out in strong collaboration as all partners are familiar with developing energy system models using mathematical programming languages and bring their expertise in specific areas (energy sectors and methodologies) to the table.</p> <p>Task 3.1 designs the core model structures. Task 3.2 implements the data conversion from the Spine Toolbox’s generic data structure to the format used by the modelling language. The Spine Model will allow multiple ways to present energy transfers and conversions and will formulate the equations efficiently (Task 3.3). The Spine Model will also enable parallel computing (Task 3.4). Finally, the results will be postprocessed and sent to the Spine Toolbox (Task 3.5). A working model with just core elements will be provided by M11 and a full model by M19. Tasks 3.3 and 3.5 will continue almost through the whole project in order to keep the model updated based on the input from the case studies.</p> <p>D3.1 The Spine Model and its documentation:</p> <ul style="list-style-type: none"><li>- A bare bone version of the model that passes data through, M07</li><li>- A version for the case studies A1–A5, M10</li><li>- A version for the case studies B1–B5, M19</li><li>- A version with all the capabilities implemented during the project, M43</li></ul>										
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## SPINEOPT DOCUMENTATION

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The purpose of this document is to give an all-encompassing documentation of the Spine Model. During the project, the name “SpineOpt.jl” was adopted to reflect the Julia package naming convention. This document describes SpineOpt and all its capabilities implemented during the project and applied to the case studies. In line with the ambition to provide open-source and easily accessible software, the documentation of SpineOpt is hosted online and open-source on github. The link to the latest version of the documentation is <https://spine-project.github.io/SpineOpt.jl/latest/>. The link to the latest version of SpineOpt itself is <https://github.com/Spine-project/SpineOpt.jl>. For the sake of completeness, the online documentation is attached to the document at hand (accessed: 2021-04-30).

## ANNEX: ONLINE DOCUMENTATION OF SPINEOPT

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The following pages contain a full copy of the SpineOpt online documentation as they were on 30th of April 2021. The reader is instructed to go to the online version for the latest changes. SpineOpt and its documentation will continue to evolve also after the Spine project ends.

# Introduction

*SpineOpt.jl* is an integrated energy systems optimization model created as part of the [Spine project](#), striving towards adaptability for a multitude of modelling purposes. The data-driven model structure allows for highly customizable energy system descriptions, as well as flexible temporal and stochastic structures, without the need to alter the model source code directly. The methodology is based on mixed-integer linear programming (MILP), and *SpineOpt* relies on [JuMP.jl](#) for interfacing with the different solvers.

While, in principle, it is possible to run *SpineOpt* by itself, it has been designed to be used through the [Spine toolbox](#), and take maximum advantage of the data and modelling workflow management tools therein. Thus, we highly recommend installing *Spine toolbox* as well, as outlined in the [Installation](#) guide.

## Contents

In order to make it easier for you to familiarize yourself with the documentation, here's a list of all the different chapters, as well as descriptions of what they're about.

## Getting Started

As the name implies, this chapter contains guides for starting to use *SpineOpt.jl* for the first time. The [Installation](#) section contains a step-by-step guide for how to install *SpineOpt.jl* and *Spine Toolbox* on your computer. The [Setting up a workflow for SpineOpt in Spine Toolbox](#) section explains how to set up and run *SpineOpt.jl* from *Spine Toolbox*. The [Creating Your Own Model](#) section explains how to create a new model from scratch. This includes a list of the necessary [Object Classes](#) and [Relationship Classes](#), but for more information, you will probably need to consult the **Concept Reference** chapter.

## Concept Reference

This chapter lists and explains all the important *data and model structure related concepts* to understand in *SpineOpt.jl*. For a mathematical modelling point of view, see the **Mathematical Formulation** chapter instead. The [Basics of the model structure](#) section briefly explains the general purpose of the most important concepts, like [Object Classes](#) and [Relationship Classes](#). Meanwhile, the [Object Classes](#), [Relationship Classes](#), [Parameters](#), and [Parameter Value Lists](#) sections contain detailed explanations of each and every aspect of *SpineOpt.jl*, organized into the respective sections for clarity.

## Mathematical Formulation

This chapter provides the mathematical view of *SpineOpt.jl*, as some of the methodology-related aspects of the model are more easily understood as math than Julia code. The [Variables](#) section explains the purpose of each variable in the model, as well as how the variables are related to the different [Object Classes](#) and [Relationship Classes](#). The [Constraints](#) section contains the mathematical formulation of each constraint, as well as explanations to their purpose and how they are controlled via different [Parameters](#). Finally, the [Objective](#) section explains the default objective function used in *SpineOpt.jl*.

## Advanced Concepts

This chapter explains some of the more complicated aspects of *SpineOpt.jl* in more detail, hopefully making it easier for you to better understand and apply them in your own modelling. The first few sections focus on aspects of *SpineOpt.jl* that most users are likely to use, or which are more or less required to understand for advanced use. The [Temporal Framework](#) section explains how defining *time* works in *SpineOpt.jl*, and how it can be used for different purposes. The [Stochastic Framework](#) section details how different stochastic structures can be defined, how they interact with each other, and how this impacts writing [Constraints](#) in *SpineOpt.jl*. The [Unit commitment](#) section explains how clustered unit-commitment is defined, while the [Ramping and Reserves](#) section explains how to enable these operational details in your model. The [Investment Optimization](#) section explains how to include investment variables in your models, while the [Unit Constraints](#) section details how to include generic data-driven custom constraints.

The last few sections focus on highly specialized use-cases for *SpineOpt.jl*, which are unlikely to be relevant for simple modelling tasks. The [Decomposition](#) section explains the Benders decomposition implementation included in *SpineOpt.jl*, as well as how to use it. The remaining sections, namely [PTDF-Based Powerflow](#), [Pressure driven gas transfer](#), [Lossless nodal DC power flows](#), and [Representative days with seasonal storages](#), explain various use-case specific modelling approaches supported by *SpineOpt.jl*.

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[Installation »](#)

# Compatibility

This package requires Julia 1.2 or later.

## Installation

SpineOpt is cross-platform (Linux, Mac and Windows) and uses other cross-platform tools. The installation process includes several steps, since there are two other pieces of software that make the use of SpineOpt more convenient (Spine Toolbox and Conda) and two programming languages that are needed (Python for Spine Toolbox and Julia for SpineOpt). Python will be installed with Conda while Julia will be setup for Spine Toolbox (explained below).

You may skip parts of the following installation process if you already have some of these software available - but please make sure they are in a clean Conda environment to avoid compatibility issues between different package versions.

SpineOpt and Spine Toolbox are under active development and the getting started process could change. If you notice any problems with these instructions, please check if it is a known issue, and if not, then report an [issue](#) or start a [discussion](#) if you're unsure whether it is an actual issue.

- The recommended interface to SpineOpt is [Spine Toolbox](#). Install Spine Toolbox following instructions from here: [Spine Toolbox installation](#)
- Setup Julia for Spine Toolbox: [Start Spine Toolbox](#). Go to *File -> Settings -> Tools*. Either select an existing Julia installation or press *Install Julia* and follow the instructions.
- Select a Julia Kernel spec. If there is none, you may need to define and install a Kernel specification using the dialog under *Kernel spec editor*. Use the newly installed Julia, give it a name and a path to a directory where it should put the files related to the Julia project (SpineOpt in this case). The process may also install iJulia, which allows you to interact with Julia code inside Spine Toolbox using the Julia console.
- Install SpineOpt from a Julia console (you can use the Julia console in Spine Toolbox: Go to *Consoles -> Start Julia Console*)

```
julia> using Pkg

julia> pkg"registry add https://github.com/Spine-project/SpineJuliaRegistry"

julia> pkg"add SpineOpt"
```

This may take a while and nothing seems to happen, but the installation process should be ongoing.

- Add SpineOpt tool icons to Spine Toolbox. Go to *Plugins* -> *Install plugins* and select and install SpineOpt.
- You should get a new ribbon in the toolbar with *Run SpineOpt* and *Load template*



After this, you have SpineOpt available as a tool in Spine Toolbox, but next you need to setup a workflow including input and output databases. Instructions are in the next section [here](#).

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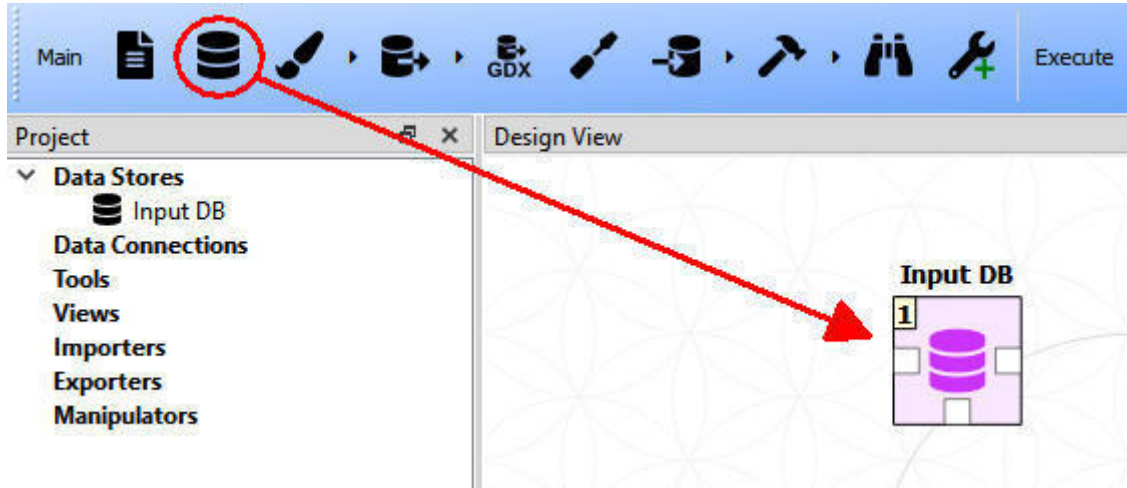
[« Introduction](#)

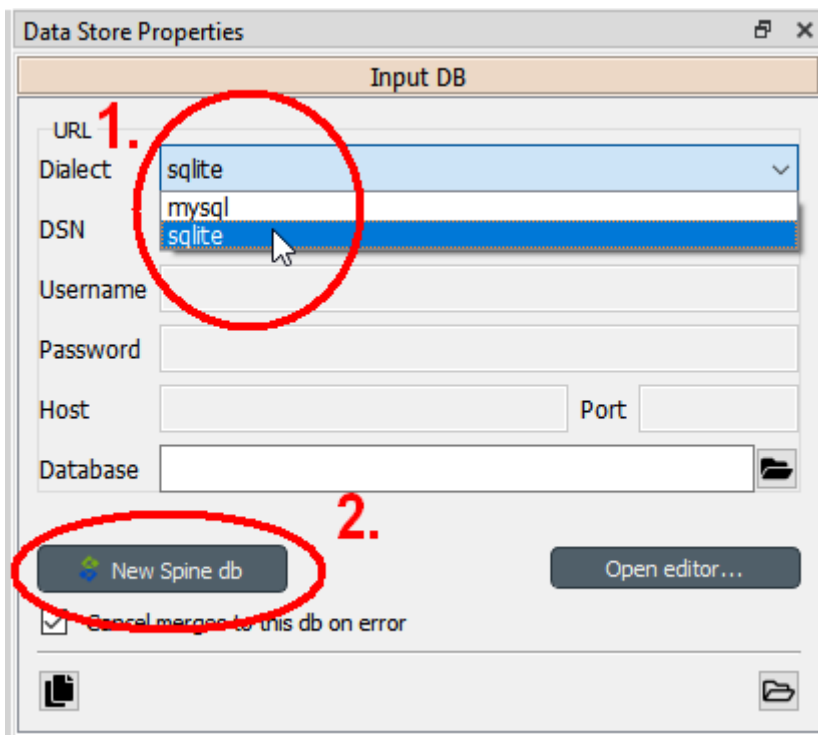
[Setting up a workflow »](#)

# Setting up a workflow for SpineOpt in Spine Toolbox

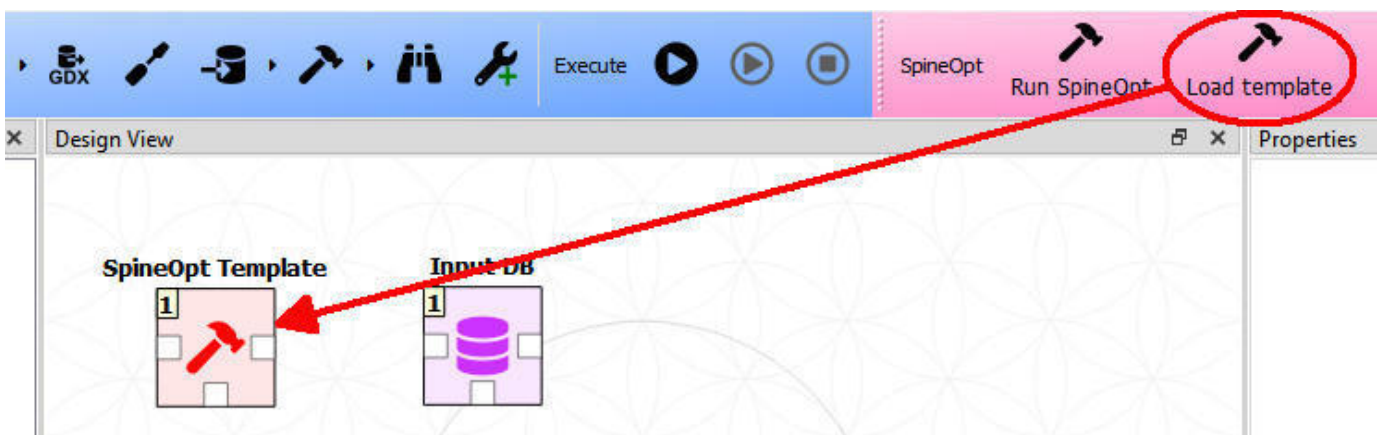
The next steps will set up a SpineOpt specific input database by creating a new Spine database, loading a blank SpineOpt template, connecting it to a SpineOpt instance and setting up a database for model results.

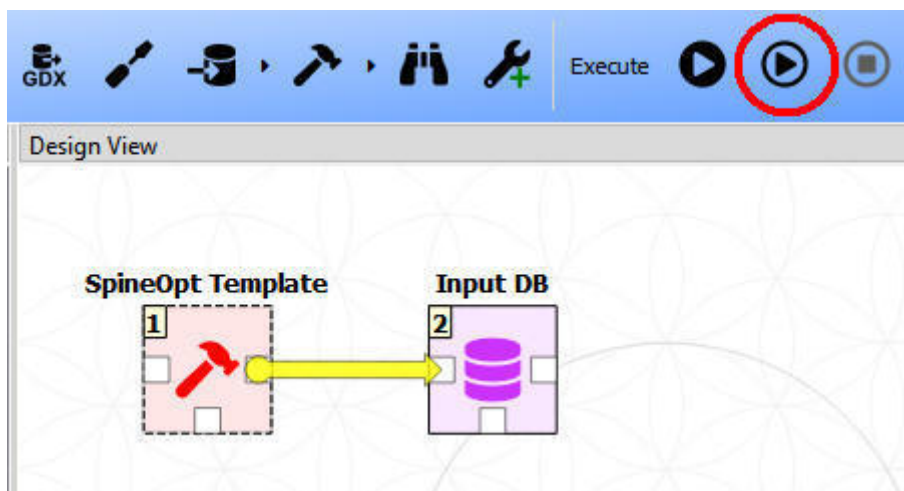
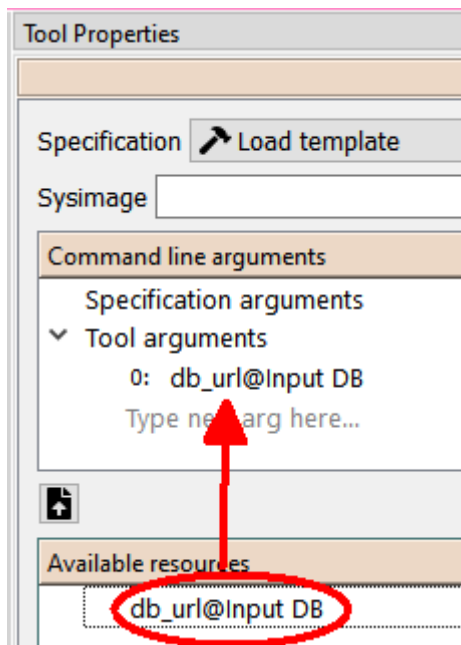
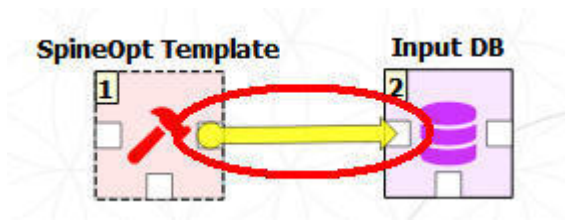
- Create a new Spine Toolbox project in an empty folder of your choice: *File -> New project...*
- Create the input database
  - Drag an empty *Data store* from the toolbar to the *Design View*.
  - Give it a name like "Input DB".
  - Select SQL database dialect (sqlite is a local file and works without a server).
  - Click *New Spine DB* in the *Data Store Properties* window and create a new database (and save it, if it's sqlite).
  - For more information about creating and managing Spine Toolbox database, see the [documentation](#)





- Fill the *Input DB* with SpineOpt data format **either** by:
  - Drag a tool *Load template* from the SpineOpt ribbon to the *Design View*.
  - Connect an arrow from the *Load template* to the new *Input DB*.
  - Make sure the *Load template* item from the Design view is selected (then you can edit the properties of that workflow item in the *Tool properties* window).
  - Add the url link in *Available resources* to the *Tool arguments* - you are passing the database address as a command line argument to the load\_template.jl script so that it knows where to store the output.
  - Then execute the *Load template* tool. Please note that this process uses SpineOpt to generate the data structure. It takes time, since everything is compiled when running a tool in Julia for the first time in each Julia session. You may also see lot of messages and warnings concerning the compilation, but they should be benign.





- ...or by:

- Start Julia (you can start a separate Julia console in Spine Toolbox: go to *Consoles* -> *Start Julia Console*).
- Copy the URL address of the Data Store from the 'Data Store Properties' -> a copy icon at the bottom.
- Then run the following script with the right URL address pasted. The process uses SpineOpt itself to build the database structure. Please note that 'using SpineOpt' for the first time for each Julia session takes time - everything is being compiled.

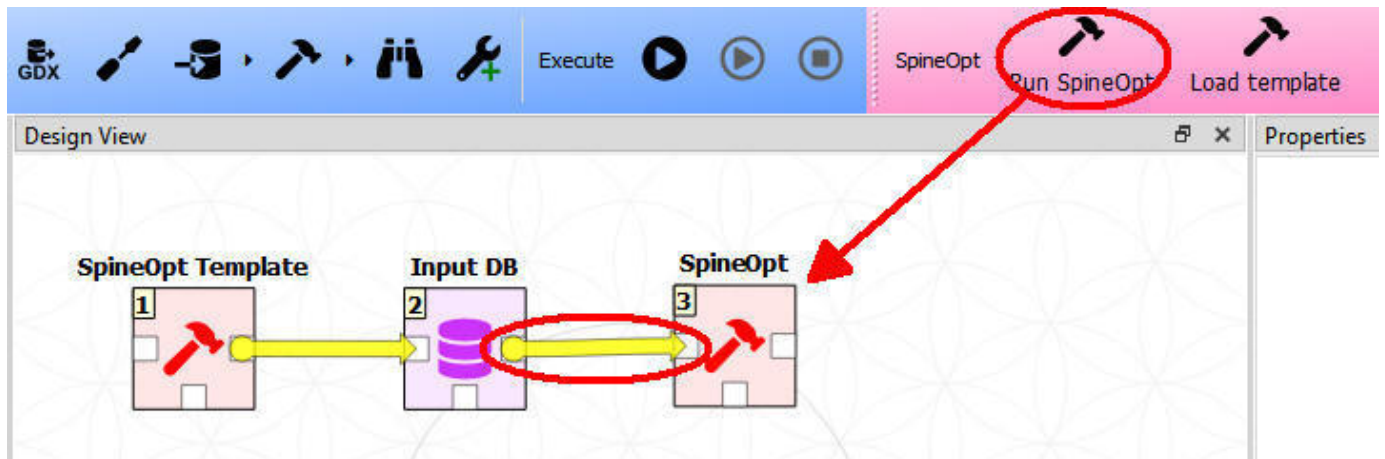
```
```julia julia> using SpineOpt
```

```
julia> SpineOpt.import_data("copied URL address, inside these quotes", SpineOpt.template(), "Load SpineOpt template")``` Known issue: On Windows, the backslash between directories need to be
```

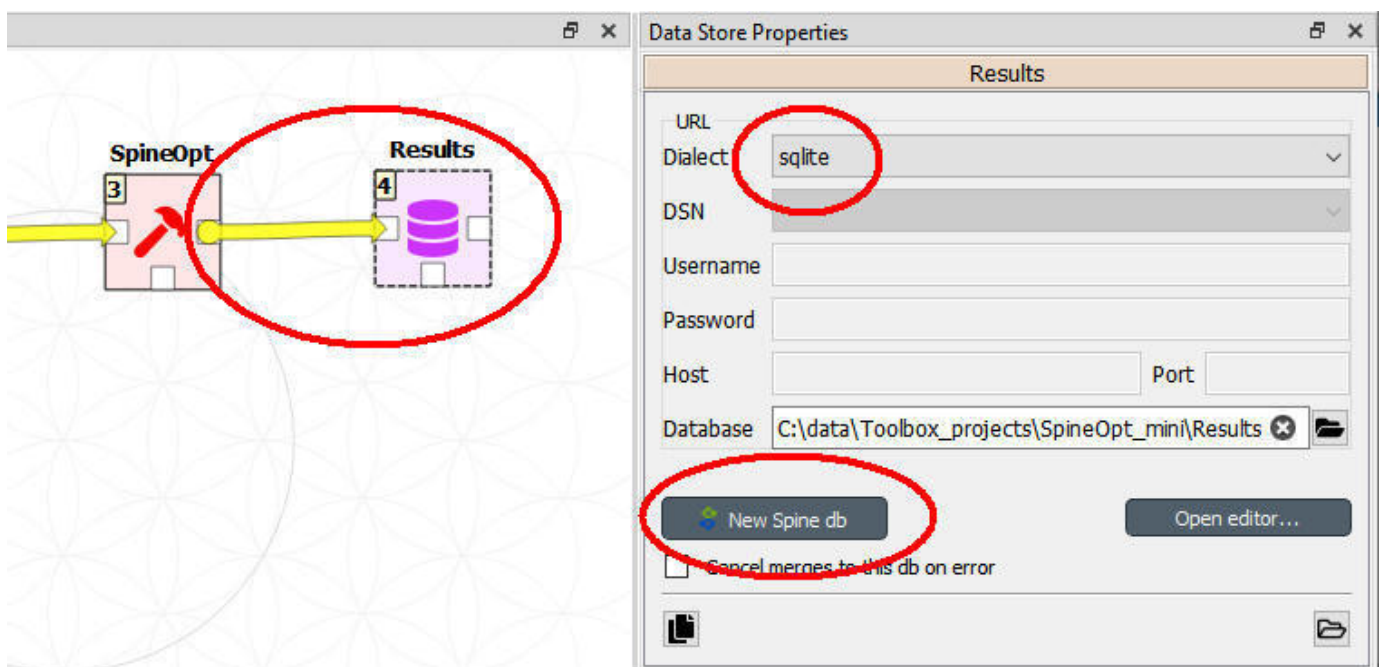


changed to a double forward slash.

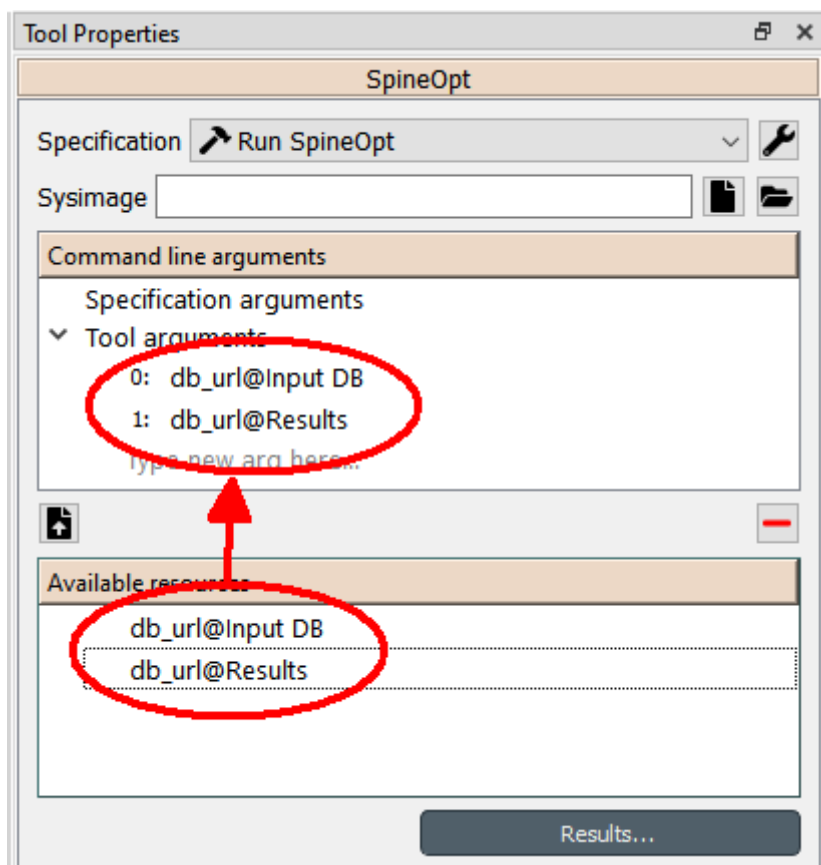
- Drag SpineOpt tool icon to the *Design view*.
- Connect an arrow from the *Input DB* to *SpineOpt*.



- Create a database for results
  - Drag a new *Data store* from the toolbar to the *Design View*.
  - You can rename it to e.g. *Results*. Select SQL database dialect (sqlite is a local file and works without a server).
  - Click *New Spine DB* in the *Data Store Properties* window and create a new database (and save it, if it's sqlite).
  - Connect an arrow from the *SpineOpt* to *Results*.



- Select *SpineOpt* tool in the *Design view*.
- Add the url link for the input data store and the output data store from *Available resources* to the *Tool arguments* (in that order).



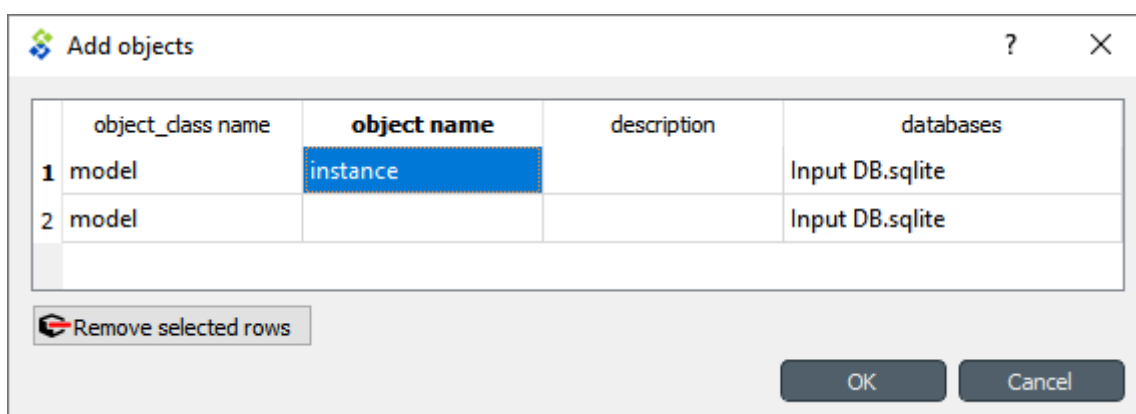
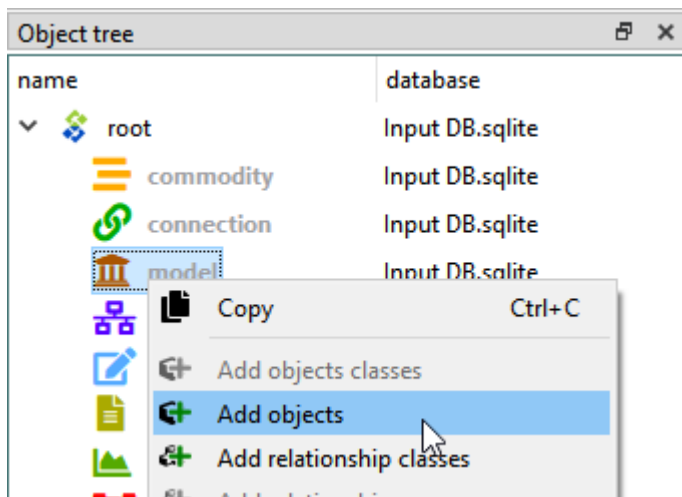
SpineOpt would be ready to run, but for the *Input DB*, which is empty of content (it's just a template that contains a SpineOpt specific data structure). The [next step](#) goes through setting up and running a simple toy model.

# Creating Your Own Model

This part of the guide shows first an example how to insert objects and their parameter data. Then it shows what other objects, relationships and parameter data needs to be added for a very basic model. Lastly, the model instance is run.



## Creating a SpineOpt model instance

- First, open the database editor by double-clicking the *Input DB*.
- Right click on *model* in the *Object tree*.
- Choose *Add objects*.
- Then, add a model object by writing a name to the *object name* field. You can use e.g. *instance*.
- Click ok.
- The **model** object in SpineOpt is an abstraction that represents the model itself. Every SpineOpt database needs to have at least one **model** object.
- The model object holds general information about the optimization. The whole range of functionalities is explained in **Advanced Concepts** chapter - in here a minimal set of parameters is used.



# Add parameter values to the model instance

- Select the model object `instance` from the object tree.
- Go to the `Object parameter value` tab.
- Every parameter value belongs to a specific alternative. This allows to hold multiple values for the same parameter of a particular object. The alternative values are used to create scenarios. Choose, `Base` for all parameter values (`Base` is required in Spine Toolbox - all other alternatives can be chosen freely).
- Then define a `model_start` time and a `model_end` time.
  - Double-click on the empty row under `parameter_name` and select `model_start`.
  - A `None` should appear in `value` column.
  - To assign a start date value, right-click on `None` and open the editor (cannot be entered directly, since the datatype needs to be changed).
  - The parameter type of `model_start` is of type `Datetime`.
  - Set the value to e.g. `2019-01-01T00:00:00`.
  - Proceed accordingly for the `model_end`.

Object parameter value					
object_class_name	object_name	parameter_name	alternative_name	value	database
 model	instance	model_end	Base	2019-01-01 23:00:00	Input DB
 model	instance	model_start	Base	2019-01-01 00:00:00	Input DB
model				None	Input DB

Further reading on adding parameter values can be found [here](#).

## Add other necessary objects and parameter data for the objects.

- Add all objects and their parameter data by replicating what has been done in the picture below. Do it the same way as explained above with the following caveats.
- Whilst most object names can be freely defined by the user, there is one object name in the example below that needs to be written exactly since it is used internally by SpineOpt: `unit_flow`.
- The `parameter_name` can be selected from a drop down menu.
- The date time and time series parameter data can be added by using right-click to access the *Edit...* dialog. When creating the time series, use the fixed resolution with `Start time` of the model run and with 1h resolution. Then only values need to be entered (or copy pasted) and time stamps come automatically.
- Parameter `balance_type` needs to have value `balance_type_none` in the gas node, since it allows the node to create energy (natural gas) against a price and therefore the energy balance is not maintained.

Object tree		Object parameter value					
name	database	object_class_name	object_name	parameter_name	alternative_name	value	database
root	Input DB	model	instance	model_end	Base	2019-01-01 23:00:00	Input DB
commodity	Input DB	model	instance	model_start	Base	2019-01-01 00:00:00	Input DB
connection	Input DB	node	node_elec	demand	Base	Time series	Input DB
model	Input DB	node	node_gas	balance_type	Base	balance_type_none	Input DB
instance	Input DB					None	Input DB
node	Input DB						
node_elec	Input DB						
node_gas	Input DB						
output	Input DB						
unit_flow	Input DB						
report	Input DB						
results	Input DB						
stochastic_scenario	Input DB						
realization	Input DB						
stochastic_structure	Input DB						
deterministic	Input DB						
temporal_block	Input DB						
hourly	Input DB						
unit	Input DB						
unit_constraint	Input DB						

## Define temporal and stochastic structures

- To specify the temporal structure for SpineOpt, you need to define **temporal\_block** objects. Think of a **temporal\_block** as a distinctive way of 'slicing' time across the model horizon.
- To link the temporal structure to the spatial structure, you need to specify **node\_\_temporal\_block** relationships, establishing which **temporal\_block** applies to each node. This relationship is added by right-clicking the **node\_\_temporal\_block** in the relationship tree and then using the **add relationships...** dialog. Double clicking on an empty cell gives you the list of valid objects. The relationship name is automatically formed, but you can change it if that is desirable.
- To keep things simple at this point, let's just define one **temporal\_block** for our model and apply it to all nodes. We add the object **hourly\_temporal\_block** of type **temporal\_block** following the same procedure as before and establish **node\_\_temporal\_block** relationships between **node\_gas** and **hourly\_temporal\_block**, and **electricity\_node** and **hourly\_temporal\_block**.
- In practical terms, the above means that there energy flows over **gas\_node** and **electricity\_node** for each 'time-slice' comprised in **hourly\_temporal\_block**.
- Similarly with the stochastic structure, each node is assigned a **deterministic stochastic\_structure**.

## Define the spatial structure

- To specify the spatial structure for SpineOpt, you will need to use the **node**, **unit**, and **connection** objects added before.

- Nodes can be understood as spatial aggregators. In combination with units and connections, they form the energy network.
- Units in SpineOpt represent any kind of conversion process. As one example, a unit can represent a power plant that converts the flow of a commodity fuel into an electricity and/or heat flow.
- Connections on the other hand describe the transport of goods from one location to another. Electricity lines and gas pipelines are examples of such connections. This example does not use connections.
- The database should have an object `gas_turbine` for the `unit` object class and objects `node_gas` and `node_elec` for the `node` object class.
- Next, define how the `unit` and the `nodes` interact with each other: create a `unit_from_node` relationship between `gas_turbine` and `node_gas`, and `unit_to_node` relationships between `gas_turbine` and `node_elec`.
- In practical terms, the above means that there is an energy flow going from `node_gas` into `node_elec`, through the `gas_turbine`.

## Add remaining relationships and parameter data for the relationships.

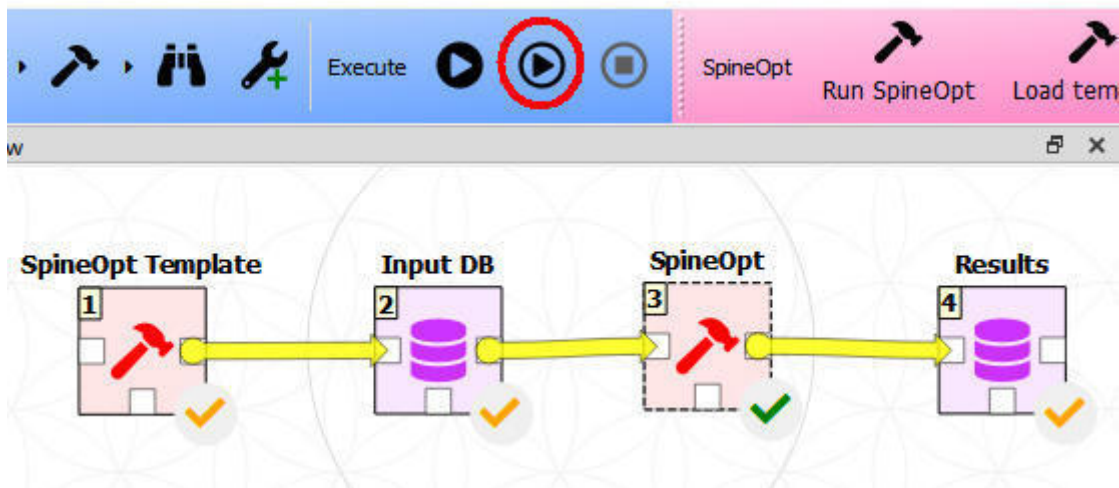
- Similar to adding the objects and their parameter data, add the relationships and their parameter data based on the picture below.
- The capacity of the `gas_turbine` has to be sufficient to meet the highest demand for electricity, otherwise the model will be infeasible (it is possible to set penalty values, but they are not included in this example).
- The parameter `fix_ratio_in_out_unit_flow` forces the ratio between an input and output flow to be a constant. This is one way to establish an efficiency for a conversion process.

Relationship tree		Relationship parameter value				
name		relationship_class_name	object_name_list	parameter_name	native_r	value
▼ root		unit_from_node	gas_turbine   node_gas	operating_cost	Base	20.0
connection_from_node		unit_from_node	gas_turbine   node_gas	unit_capacity	Base	200.0
connection_from_node_unit_constraint		unit_node_node	gas_turbine   node_gas   node_elec	fix_ratio_in_out_unit_flow	Base	0.5
connection_investment_stochastic_structure		unit_to_node	gas_turbine   node_elec	unit_capacity	Base	100.0
connection_investment_temporal_block						None
connection_node_node						
connection_to_node						
connection_to_node_unit_constraint						
model_default_investment_stochastic_structure						
model_default_investment_temporal_block						
model_default_stochastic_structure						
model_default_temporal_block						
▼ model_report						
instance   results						
▼ model_stochastic_structure						
instance   deterministic						
▼ model_temporal_block						
instance   hourly						
node_commodity						
node_investment_stochastic_structure						
node_investment_temporal_block						
node_node						
▼ node_stochastic_structure						
node_elec   deterministic						
node_gas   deterministic						
▼ node_temporal_block						
node_elec   hourly						
node_gas   hourly						
node_unit_constraint						
parent_stochastic_scenario_child_stochastic_scenario						
▼ report_output						
results   unit_flow						
▼ stochastic_structure_stochastic_scenario						
deterministic   realization						
unit_commodity						
▼ unit_from_node						
gas_turbine   node_gas						
unit_from_node_unit_constraint						
unit_investment_stochastic_structure						
unit_investment_temporal_block						
▼ unit_node_node						
gas_turbine   node_gas   node_elec						
▼ unit_to_node						
gas_turbine   node_elec						
unit_to_node_unit_constraint						
unit_unit_constraint						
▼ units_on_stochastic_structure						
gas_turbine   deterministic						
▼ units_on_temporal_block						
gas_turbine   hourly						

Run the model



- Select *SpineOpt*
- Press *Execute selection*.

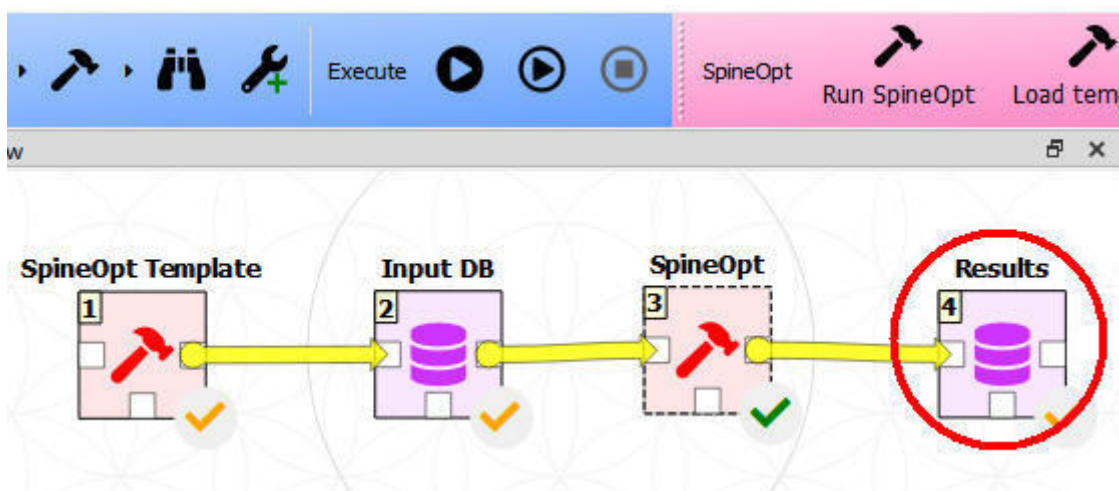


## If it fails

- Double-check that the data is correct
- Try to see what the problem might be
- Ask help from the [discussion forum](#)

## Explore the results

- Double-clicking the *Results* database.



## Create and run scenarios and build the model further

- Create a new alternative
- Add parameter data for the new alternative

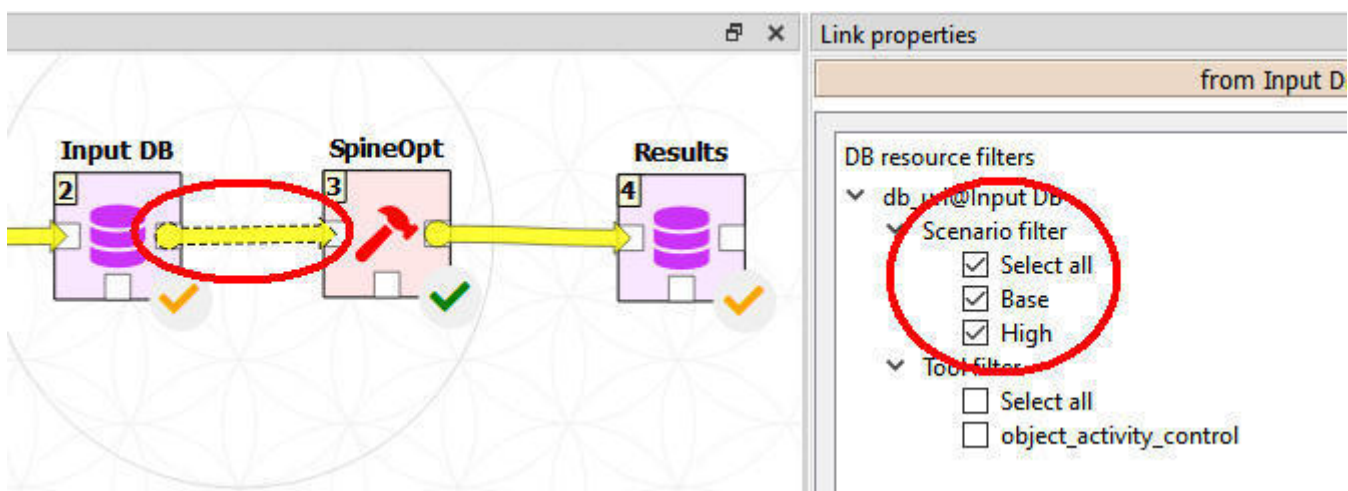


- Connect alternatives under a scenario. Toolbox modifies Base data with the data from the alternatives in the same scenario.
- Execute multiple scenarios in parallel. First run in a new Julia instance will need to compile SpineOpt taking some time.

Object parameter value					
object_class_name	object_name	parameter_name	alternative_name	value	database
model	instance	model_end	Base	2019-01-01 23:00:00	Input DB
model	instance	model_start	Base	2019-01-01 00:00:00	Input DB
node	node_elec	demand	Base	Time series	Input DB
node	node_gas	balance_type	Base	balance_type_none	Input DB
				None	Input DB

Relationship parameter value					
relationship_class_name	object_name_list	parameter_name	alternative_name	value	database
unit_from_node	gas_turbine   node_gas	operating_cost	Base	20.0	Input DB
unit_from_node	gas_turbine   node_gas	operating_cost	High	30.0	Input DB
unit_from_node	gas_turbine   node_gas	unit_capacity	Base	200.0	Input DB
unit_node_node	gas_turbine   node_gas   node_elec	fix_ratio_in_out_unit_flow	Base	0.5	Input DB
unit_to_node	gas_turbine   node_elec	unit_capacity	Base	100.0	Input DB
				None	Input DB



# Basics of the model structure

In *SpineOpt.jl*, the model structure is generated based on the input data, allowing it to be used for a multitude of different problems. Here, we aim to provide you with a basic understanding of the *SpineOpt.jl* model and data structure, while the [Object Classes](#), [Relationship Classes](#), [Parameters](#), and [Parameter Value Lists](#) sections provide more in-depth explanations of each concept.

## Introduction to object classes

Essentially, [Object Classes](#) represents different types of `objects` or *entities* that make up the model. For example, every power plant in the model is represented as an object of the object class `unit`, every power line as an object of the object class `connection`, and so forth. In order to add any new *entity* to a model, a new `object` has to be added to desired `object class` in the input data.

Each `object class` has a very specific purpose in *SpineOpt.jl*, so understanding their differences is key. The [Object Classes](#) can be roughly divided into three distinctive groups, namely [Systemic object classes](#), [Structural object classes](#), and [Meta object classes](#).

## Systemic object classes

As the name implies, *system [Object Classes](#)* are used to describe the system to be modelled. Essentially, they define *what* you want to model. These include:

- `commodity` represents different goods to be generated, consumed, transported, etc.
- `connection` handles the transfer of `commodities` between `nodes`.
- `node` ensures the balance of the `commodity` flows, and can be used to store `commodities` as well.
- `unit` handles the generation and consumption of `commodities`.

## Structural object classes

*Structural [Object Classes](#)* are used to define the temporal and stochastic structure of the modelled problem, as well as custom [Unit Constraints](#). Unlike the above *system [Object Classes](#)*, the *structural [Object Classes](#)* are more about *how* you want to model, instead of strictly *what* you want to model. These include:

- `stochastic_scenario` represents a different *forecast* or another type of an *alternative time period*.
- `stochastic_structure` acts as a handle for a group of `stochastic_scenarios` with set properties.
- `temporal_block` defines a period of *time* with the desired temporal [resolution](#).
- `unit_constraint` is an optional custom constraint generated based on the input data.

# Meta object classes

*Meta Object Classes* are used for defining things on the level of `models` or above, like `model output` and even multiple `models` for problem decompositions. These include:

- `model` represents an individual *model*, grouping together all the things relevant for itself.
- `output` defines which `Variables` are output from the `model`.
- `report` groups together multiple `output` objects.

# Introduction to relationship classes

While *Object Classes* define all the `objects` or *entities* that make up a `model`, *Relationship Classes* define how those *entities* are related to each other. Thus, *Relationship Classes* hold no meaning on their own, and always include at least one `object class`.

Similar to *Object Classes*, each `relationship class` has a very specific purpose in *SpineOpt.jl*, and understanding the purpose of each `relationship class` is paramount. The *Relationship Classes* can be roughly divided into *Systemic relationship classes*, *Structural relationship classes*, and *Meta relationship classes*, again similar to *Object Classes*.

## Systemic relationship classes

*Systemic Relationship Classes* define how *Systemic object classes* are related to each other, thus helping define the system to be modelled. Most of these relationships deal with *which* `units` and `connections` interact with *which* `nodes`, and *how* those interactions work. This essentially defines the possible `commodity` flows to be modelled. *Systemic Relationship Classes* include:

- `connection_from_node` defines which `node` the `connection` can transfer a `commodity` from.
- `connection_to_node` defines which `node` the `connection` can transfer a `commodity` to.
- `connection_node_node` holds `Parameters` for `connections` between two `nodes`.
- `node_commodity` defines which `node` holds which `commodity`.
- `node_node` holds parameters for direct `node-node` interactions, like diffusion of `commodities`.
- `unit_commodity` defines which `commodity` the `unit` handles.
- `unit_from_node` defines which `node` the `unit` can take an input `commodity` from.
- `unit_to_node` defines which `node` the `unit` can output a `commodity` to.
- `unit_node_node` holds parameters for `unit` interactions between two `nodes`.

## Structural relationship classes

*Structural Relationship Classes* primarily relate *Structural object classes* to *Systemic object classes*, defining what *structures* the individual parts of the *system* use. These are mostly used to determine the

temporal and stochastic structures to be used in different parts of the modelled *system*, or custom [Unit Constraints](#).

*SpineOpt.jl* has a very flexible temporal and stochastic structure, explained in detail in the [Temporal Framework](#) and [Stochastic Framework](#) sections of the documentation. Unfortunately, this flexibility requires quite a few different *structural Relationship Classes*, the most important of which are the following *basic structural Relationship Classes*:

- [node\\_stochastic\\_structure](#) defines the [stochastic\\_structure](#) used for the [node](#) balance.
- [node\\_temporal\\_block](#) defines the `temporal blocks` used for the [node](#) balance.
- [parent\\_stochastic\\_scenario\\_child\\_stochastic\\_scenario](#) defines the *stochastic directed acyclic graph (DAG)* of the [Stochastic Framework](#).
- [stochastic\\_structure\\_stochastic\\_scenario](#) holds parameters for `stochastic scenarios` in the [stochastic\\_structure](#).
- [units\\_on\\_stochastic\\_structure](#) defines the [stochastic\\_structure](#) used for the online variable of the [unit](#).
- [units\\_on\\_temporal\\_block](#) defines the `temporal blocks` used for the online variable of the [unit](#).

Furthermore, there are also a number of *advanced structural Relationship Classes*, which are only necessary when using some of the optional features of *SpineOpt.jl*. For [Investment Optimization](#), the following relationships control the stochastic and temporal structures of the investment [variables](#):

- [connection\\_investment\\_stochastic\\_structure](#) defines the [stochastic\\_structure](#) used for the investment [Variables](#) for the [connection](#).
- [connection\\_investment\\_temporal\\_block](#) defines the `temporal blocks` used for the investment [Variables](#) for the [connection.unit\\_constraint](#).
- [node\\_investment\\_stochastic\\_structure](#) defines the [stochastic\\_structure](#) used for the investment [Variables](#) for the [node](#).
- [node\\_investment\\_temporal\\_block](#) defines the [stochastic\\_structure](#) used for the investment [Variables](#) for the [node](#).
- [unit\\_investment\\_stochastic\\_structure](#) defines the [stochastic\\_structure](#) used for the investment [Variables](#) for the [unit](#).
- [unit\\_investment\\_temporal\\_block](#) defines the `temporal blocks` used for the investment [Variables](#) for the [unit](#).(@ref).

For [Unit Constraints](#), which are essentially generic data-driven custom constraints, the following relationships are used to control which [variables](#) are included and with what coefficients:

- [connection\\_from\\_node\\_unit\\_constraint](#) holds [Parameters](#) for the [connection\\_flow](#) variable *from* the [node](#) in question in the custom [unit\\_constraint](#).
- [connection\\_to\\_node\\_unit\\_constraint](#) holds [Parameters](#) for the [connection\\_flow](#) variable *to* the [node](#) in question in the custom [unit\\_constraint](#).
- [node\\_unit\\_constraint](#) holds [Parameters](#) for the [node\\_state](#) variable in the custom [unit\\_constraint](#).

- `unit_from_node_unit_constraint` holds [Parameters](#) for the `unit_flow` variable *from* the `node` in question in the custom `unit_constraint`.
- `unit_to_node_unit_constraint` holds [Parameters](#) for the `unit_flow` variable *to* the `node` in question in the custom `unit_constraint`.

## Meta relationship classes

*Meta Relationship Classes* are used for defining [model](#)-level settings, like which `temporal blocks` or `stochastic structures` are active, and what the `model output` is. These include:

- `model_default_investment_stochastic_structure` defines a default `stochastic_structure` to be used for investment [Variables](#) when no other definitions exist.
- `model_default_investment_temporal_block` defines a default `temporal_block` to be used for investment [Variables](#) when no other definitions exist.
- `model_default_stochastic_structure` defines a default `stochastic_structure` to be used for `nodes` and `units` when no other definitions exist.
- `model_default_temporal_block` defines a default `temporal_block` to be used for `nodes` and `units` when no other definitions exist.
- `model_report` connects each `report` to the desired `model`.
- `model_stochastic_structure` defines which `stochastic structures` are active in which `models`.
- `model_temporal_block` defines which `temporal blocks` are active in which `models`.
- `report_output` defines which `outputs` are part of which `report`.

## Introduction to parameters

While the primary function of [Object Classes](#) and [Relationship Classes](#) is to *define* the system to be modelled and its structure, [Parameters](#) exist to *constrain* them. Every `parameter` is attributed to at least one `object class` or `relationship class`, but some appear in many classes whenever they serve a similar purpose.

[Parameters](#) accept different types of values depending on their purpose, e.g. whether they act as a *flag* for some specific functionality or appear as a *coefficient* in [Constraints](#), so understanding each `parameter` is key. Most coefficient-type [Parameters](#) accept *constant*, *time series*, and even *stochastic time series* form input, but there are some exceptions. Most *flag-type* [Parameters](#), on the other hand, have a restricted list of acceptable values defined by their [Parameter Value Lists](#).

The existence of some [Constraints](#) is controlled based on if the relevant [Parameters](#) are defined. As a rule-of-thumb, a `constraint` only gets generated if at least one of the [Parameters](#) appearing in it is defined, but one should refer to the appropriate [Constraints](#) and [Parameters](#) sections when in doubt.

## Introduction to groups of objects

Groups of objects are used within SpineOpt for different purposes. To create a group of objects, simply right-click the corresponding [Object Class](#) in the *Spine Toolbox* database editor and select [Add object group](#). Groups are essentially special [objects](#), that act as a single handle for all of its members.

On the one hand, groups can be used in order to impose constraints on the aggregation of a variable, e.g. on the sum of multiple [unit\\_flow](#) variables. Constraints based on parameters associated with the [unit\\_node\\_node](#), [unit\\_to\\_node](#), [unit\\_from\\_node](#), [connection\\_node\\_node](#), [connection\\_to\\_node](#), [connection\\_from\\_node](#) can generally be used for this kind of flow aggregation by defining the parameters on groups of objects, typically node groups. (with the exception of variable fixing parameters, e.g. [fix\\_unit\\_flow](#), [fix\\_connection\\_flow](#) etc.). See for instance [constraint\\_unit\\_flow\\_capacity](#).

On the other hand, a node group can be used to for [PTDF based powerflows](#). Here a node group is used to enforce a nodal balance on system level, while suppressing the node balances at individual nodes. See also [balance\\_type](#) and [the node balance constraint](#).

# Object Classes

## commodity

A good or product that can be consumed, produced, traded. E.g., electricity, oil, gas, water...

**Related Parameters:** [commodity\\_lodf\\_tolerance](#), [commodity\\_physics](#), [commodity\\_ptdf\\_threshold](#) and [is\\_active](#)

**Related Relationship Classes:** [node\\_commodity](#) and [unit\\_commodity](#)

Commodities correspond to the type of energy traded. When associated with a [node](#) through the [node\\_commodity](#) relationship, a specific form of energy, i.e. commodity, can be associated with a specific location. Furthermore, by linking commodities with [units](#), it is possible to track the flows of a certain commodity and impose limitations on the use of a certain commodity (See also [max\\_cum\\_in\\_unit\\_flow\\_bound](#)). For the representation of specific commodity physics, related to e.g. the representation of the electric network, designated parameters can be defined to enforce commodity specific behaviour. (See also [commodity\\_physics](#))

## connection

A transfer of commodities between nodes. E.g. electricity line, gas pipeline...

**Related Parameters:** [candidate\\_connections](#), [connection\\_availability\\_factor](#), [connection\\_contingency](#), [connection\\_flow\\_cost](#), [connection\\_investment\\_cost](#), [connection\\_investment\\_lifetime](#), [connection\\_investment\\_variable\\_type](#), [connection\\_monitored](#), [connection\\_reactance\\_base](#), [connection\\_reactance](#), [connection\\_resistance](#), [connection\\_type](#), [fix\\_connections\\_invested\\_available](#), [fix\\_connections\\_invested](#), [has\\_binary\\_gas\\_flow](#) and [is\\_active](#)

**Related Relationship Classes:** [connection\\_from\\_node\\_unit\\_constraint](#), [connection\\_from\\_node](#), [connection\\_investment\\_stochastic\\_structure](#),



`connection_investment_temporal_block`, `connection_node_node`, `connection_to_node_unit_constraint` and `connection_to_node`

A `connection` represents a transfer of one `commodity` over space. For example, an electricity transmission line, a gas pipe, a river branch, can be modelled using a `connection`.

A `connection` always takes `commodities` from one or more `nodes`, and releases them to one or more (possibly the same) `nodes`. The former are specified through the `connection_from_node` relationship, and the latter through `connection_to_node`. Every `connection` inherits the temporal and stochastic structures from the associated nodes. The model will generate `connection_flow` variables for every combination of `connection`, `node`, *direction* (from node or to node), *time slice*, and *stochastic scenario*, according to the above relationships.

The operation of the `connection` is specified through a number of parameter values. For example, the capacity of the connection, as the maximum amount of energy that can enter or leave it, is given by `connection_capacity`. The conversion ratio of input to output can be specified using any of `fix_ratio_out_in_connection_flow`, `max_ratio_out_in_connection_flow`, and `min_ratio_out_in_connection_flow` parameters in the `connection_node_node` relationship. The delay on a connection, as the time it takes for the energy to go from one end to the other, is given by `connection_flow_delay`.

## model

An instance of `SpineOpt`, that specifies general parameters such as the temporal horizon.

**Related Parameters:** `big_m`, `duration_unit`, `is_active`, `max_gap`, `max_iterations`, `model_end`, `model_start`, `model_type`, `roll_forward`, `write_lodf_file`, `write_mps_file` and `write_ptdf_file`

**Related Relationship Classes:** `model_default_investment_stochastic_structure`, `model_default_investment_temporal_block`, `model_default_stochastic_structure`, `model_default_temporal_block`, `model_report`, `model_stochastic_structure` and `model_temporal_block`

The model object holds general information about the optimization problem at hand. Firstly, the modelling horizon is specified on the model object, i.e. the scope of the optimization model, and if applicable the duration of the rolling window (see also `model_start`, `model_end` and `roll_forward`). Secondly, the model works as an overarching assembler - only through linking `temporal_blocks` and `stochastic_structures` to a model object via relationships, they become part of the optimization problem, and respectively linked nodes, connections and units. If desired the user can also specify



defaults for temporals and stochastic via the designated default relationships (see e.g., [model\\_default\\_temporal\\_block](#)). In this case, the default temporal is populated for missing [node\\_temporal\\_block](#) relationships. Lastly, the model object contains information about the algorithm used for solving the problem (see [model\\_type](#)).

## node

A universal aggregator of commodity flows over units and connections, with storage capabilities.

**Related Parameters:** [balance\\_type](#), [candidate\\_storages](#), [demand](#), [downward\\_reserve](#), [fix\\_node\\_pressure](#), [fix\\_node\\_state](#), [fix\\_node\\_voltage\\_angle](#), [fix\\_storages\\_invested\\_available](#), [fix\\_storages\\_invested](#), [frac\\_state\\_loss](#), [fractional\\_demand](#), [has\\_pressure](#), [has\\_state](#), [has\\_voltage\\_angle](#), [is\\_active](#), [is\\_reserve\\_node](#), [max\\_node\\_pressure](#), [max\\_voltage\\_angle](#), [min\\_node\\_pressure](#), [min\\_voltage\\_angle](#), [minimum\\_reserve\\_activation\\_time](#), [nodal\\_balance\\_sense](#), [node\\_opf\\_type](#), [node\\_slack\\_penalty](#), [node\\_state\\_cap](#), [node\\_state\\_min](#), [state\\_coeff](#), [storage\\_investment\\_cost](#), [storage\\_investment\\_lifetime](#), [storage\\_investment\\_variable\\_type](#), [tax\\_in\\_unit\\_flow](#), [tax\\_net\\_unit\\_flow](#), [tax\\_out\\_unit\\_flow](#) and [upward\\_reserve](#)

**Related Relationship Classes:** [connection\\_from\\_node\\_unit\\_constraint](#), [connection\\_from\\_node](#), [connection\\_node\\_node](#), [connection\\_to\\_node\\_unit\\_constraint](#), [connection\\_to\\_node](#), [node\\_commodity](#), [node\\_investment\\_stochastic\\_structure](#), [node\\_investment\\_temporal\\_block](#), [node\\_node](#), [node\\_stochastic\\_structure](#), [node\\_temporal\\_block](#), [node\\_unit\\_constraint](#), [unit\\_from\\_node\\_unit\\_constraint](#), [unit\\_from\\_node](#), [unit\\_node\\_node](#), [unit\\_to\\_node\\_unit\\_constraint](#) and [unit\\_to\\_node](#)

The [node](#) is perhaps the most important [object class](#) out of the [Systemic object classes](#), as it is what connects the rest together via the [Systemic relationship classes](#). Essentially, [nodes](#) act as points in the modelled [commodity](#) network where [commodity](#) balance is enforced via the [node balance](#) and [node injection](#) constraints, tying together the inputs and outputs from [units](#) and [connections](#), as well as any external [demand](#). Furthermore, [nodes](#) play a crucial role for defining the temporal and stochastic structures of the [model](#) via the [node\\_temporal\\_block](#) and [node\\_stochastic\\_structure](#) relationships. For more details about the [Temporal Framework](#) and the [Stochastic Framework](#), please refer to the dedicated sections.

Since [nodes](#) act as the points where [commodity](#) balance is enforced, this also makes them a natural fit for implementing *storage*. The [has\\_state](#) parameter controls whether a [node](#) has a [node\\_state](#) variable, which essentially represents the [commodity](#) content of the [node](#). The [state\\_coeff](#) parameter

tells how the `node_state` variable relates to all the `commodity` flows. Storage losses are handled via the `frac_state_loss` parameter, and potential diffusion of `commodity` content to other `nodes` via the `diff_coeff` parameter for the `node__node` relationship.

## output

A variable name from SpineOpt that can be included in a report.

Related **Parameters**: `is_active`

Related **Relationship Classes**: `report__output`

An `output` is essentially a handle for a *SpineOpt* `variable` and `Objective function` to be included in a `report` and written into an output database. Typically, e.g. the `unit_flow` variables are desired as output from most `models`, so creating an `output` object called `unit_flow` allows one to designate it as something to be written in the desired `report`. Note that unless appropriate `model__report` and `report__output` relationships are defined, *SpineOpt* doesn't write any output!

## report

A results report from a particular SpineOpt run, including the value of specific variables.

Related **Parameters**: `is_active` and `output_db_url`

Related **Relationship Classes**: `model__report` and `report__output`

A `report` is essentially a group of `outputs` from a `model`, that gets written into the output database as a result of running *SpineOpt*. Note that unless appropriate `model__report` and `report__output` relationships are defined, *SpineOpt* doesn't write any output!

## stochastic\_scenario

A scenario for stochastic optimisation in SpineOpt.

Related [Parameters](#): [is\\_active](#)

Related [Relationship Classes](#): [parent\\_stochastic\\_scenario\\_\\_child\\_stochastic\\_scenario](#) and [stochastic\\_structure\\_\\_stochastic\\_scenario](#)

Essentially, a [stochastic\\_scenario](#) is a label for an alternative period of time, describing one possibility of what might come to pass. They are the basic building blocks of the scenario-based [Stochastic Framework](#) in *SpineOpt.jl*, but aren't really meaningful on their own. Only when combined into a [stochastic\\_structure](#) using the [stochastic\\_structure\\_\\_stochastic\\_scenario](#) and [parent\\_stochastic\\_scenario\\_\\_child\\_stochastic\\_scenario](#) relationships, along with [Parameters](#) like the [weight\\_relative\\_to\\_parents](#) and [stochastic\\_scenario\\_end](#), they become meaningful.

## stochastic\_structure

A group of stochastic scenarios that represent a structure.

Related [Parameters](#): [is\\_active](#)

Related [Relationship Classes](#): [connection\\_\\_investment\\_stochastic\\_structure](#), [model\\_default\\_investment\\_stochastic\\_structure](#), [model\\_default\\_stochastic\\_structure](#), [model\\_stochastic\\_structure](#), [node\\_investment\\_stochastic\\_structure](#), [node\\_stochastic\\_structure](#), [stochastic\\_structure\\_\\_stochastic\\_scenario](#), [unit\\_investment\\_stochastic\\_structure](#) and [units\\_on\\_\\_stochastic\\_structure](#)

The [stochastic\\_structure](#) is the key component of the scenario-based [Stochastic Framework](#) in *SpineOpt.jl*, and essentially represents a group of [stochastic\\_scenarios](#) with set [Parameters](#). The [stochastic\\_structure\\_\\_stochastic\\_scenario](#) relationship defines which [stochastic\\_scenarios](#) are included in which [stochastic\\_structures](#), and the [weight\\_relative\\_to\\_parents](#) and [stochastic\\_scenario\\_end](#) [Parameters](#) define the exact shape and impact of the [stochastic\\_structure](#), along with the [parent\\_stochastic\\_scenario\\_\\_child\\_stochastic\\_scenario](#) relationship.

The main reason as to why [stochastic\\_structures](#) are so important is, that they act as handles connecting the [Stochastic Framework](#) to the modelled system. This is handled using the [Structural relationship classes](#) e.g. [node\\_stochastic\\_structure](#), which define the [stochastic\\_structure](#) applied to each [object](#) describing the modelled system. Connecting each system [object](#) to the appropriate [stochastic\\_structure](#) individually can be a bit bothersome at times, so there are also a number of

convenience [Meta relationship classes](#) like the [model\\_default\\_stochastic\\_structure](#), which allow setting [model](#)-wide defaults to be used whenever specific definitions are missing.

## temporal\_block

A length of time with a particular resolution.

**Related Parameters:** [block\\_end](#), [block\\_start](#), [is\\_active](#), [representative\\_periods\\_mapping](#), [resolution](#) and [weight](#)

**Related Relationship Classes:** [connection\\_investment\\_temporal\\_block](#), [model\\_default\\_investment\\_temporal\\_block](#), [model\\_default\\_temporal\\_block](#), [model\\_temporal\\_block](#), [node\\_investment\\_temporal\\_block](#), [node\\_temporal\\_block](#), [unit\\_investment\\_temporal\\_block](#) and [units\\_on\\_temporal\\_block](#)

A temporal block defines the temporal properties of the optimization that is to be solved in the current window. It is the key building block of the [Temporal Framework](#). Most importantly, it holds the necessary information about the resolution and horizon of the optimization. A single model can have multiple temporal blocks, which is one of the main sources of temporal flexibility in Spine: by linking different parts of the model to different temporal blocks, a single model can contain aspects that are solved with different temporal resolutions or time horizons.

## unit

A conversion of one/many commodities between nodes.

**Related Parameters:** [candidate\\_units](#), [curtailment\\_cost](#), [fix\\_units\\_invested\\_available](#), [fix\\_units\\_invested](#), [fix\\_units\\_on](#), [fom\\_cost](#), [is\\_active](#), [min\\_down\\_time](#), [min\\_up\\_time](#), [number\\_of\\_units](#), [online\\_variable\\_type](#), [shut\\_down\\_cost](#), [start\\_up\\_cost](#), [unit\\_availability\\_factor](#), [unit\\_investment\\_cost](#), [unit\\_investment\\_lifetime](#) and [unit\\_investment\\_variable\\_type](#)

**Related Relationship Classes:** [unit\\_commodity](#), [unit\\_from\\_node\\_unit\\_constraint](#), [unit\\_from\\_node](#), [unit\\_investment\\_stochastic\\_structure](#), [unit\\_investment\\_temporal\\_block](#), [unit\\_node\\_node](#), [unit\\_to\\_node\\_unit\\_constraint](#), [unit\\_to\\_node](#), [unit\\_unit\\_constraint](#), [units\\_on\\_stochastic\\_structure](#) and [units\\_on\\_temporal\\_block](#)

A [unit](#) represents an energy conversion process, where energy of one [commodity](#) can be converted into energy of another [commodity](#). For example, a gas turbine, a power plant, or even a load, can be modelled using a [unit](#).

A [unit](#) always takes energy from one or more [nodes](#), and releases energy to one or more (possibly the same) [nodes](#). The former are specified through the [unit\\_from\\_node](#) relationship, and the latter through [unit\\_to\\_node](#). Every [unit](#) has a temporal and stochastic structures given by the [units\\_on\\_temporal\\_block](#) and [\[units\\_on\\_\\_stochastic\\_structure\]](#) relationships. The model will generate [unit\\_flow](#) variables for every combination of [unit](#), [node](#), *direction* (from node or to node), *time slice*, and *stochastic scenario*, according to the above relationships.

The operation of the [unit](#) is specified through a number of parameter values. For example, the capacity of the unit, as the maximum amount of energy that can enter or leave it, is given by [unit\\_capacity](#). The conversion ratio of input to output can be specified using any of [fix\\_ratio\\_out\\_in\\_unit\\_flow](#), [max\\_ratio\\_out\\_in\\_unit\\_flow](#), and [min\\_ratio\\_out\\_in\\_unit\\_flow](#). The variable operating cost is given by [vom\\_cost](#).

## unit\_constraint

A generic data-driven custom constraint.

**Related Parameters:** [constraint\\_sense](#), [is\\_active](#) and [right\\_hand\\_side](#)

**Related Relationship Classes:** [connection\\_from\\_node\\_unit\\_constraint](#), [connection\\_to\\_node\\_unit\\_constraint](#), [node\\_unit\\_constraint](#), [unit\\_from\\_node\\_unit\\_constraint](#), [unit\\_to\\_node\\_unit\\_constraint](#) and [unit\\_unit\\_constraint](#)

The [unit\\_constraint](#) is a generic data-driven [custom constraint](#), which allows for defining constraints involving multiple [units](#), [nodes](#), or [connections](#). The [constraint\\_sense](#) parameter changes the sense of the [unit\\_constraint](#), while the [right\\_hand\\_side](#) parameter allows for defining the constant terms of the constraint.

Coefficients for the different [variables](#) appearing in the [unit\\_constraint](#) are defined using relationships, like e.g. [unit\\_from\\_node\\_unit\\_constraint](#) and [connection\\_to\\_node\\_unit\\_constraint](#) for [unit\\_flow](#) and [connection\\_flow](#) variables, or [unit\\_unit\\_constraint](#) and [node\\_unit\\_constraint](#) for [units\\_on](#), [units\\_started\\_up](#), and [node\\_state](#) variables.

For more information, see the dedicated article on [Unit Constraints](#)



# Relationship Classes

## connection\_\_from\_node

Defines the `nodes` the `connection` can take input from, and holds most `connection_flow` variable specific parameters.

**Related Object Classes:** `connection` and `node`

**Related Parameters:** `connection_capacity`, `connection_conv_cap_to_flow`, `connection_emergency_capacity`, `fix_binary_gas_connection_flow`, `fix_connection_flow`, `fix_connection_intact_flow` and `graph_view_position`

int

`connection__from_node` is a two-dimensional relationship between a `connection` and a `node` and implies a `connection_flow` to the `connection` from the `node`. Specifying such a relationship will give rise to a `connection_flow_variable` with indices `connection=connection`, `node=node`, `direction=:from_node`. Relationships defined on this relationship will generally apply to this specific flow variable. For example, `connection_capacity` will apply only to this specific flow variable, unless the `connection` parameter `connection_type` is specified.

c\_structure

block

## connection\_\_from\_node\_\_unit\_constraint

when specified this relationship allows the relevant flow connection flow variable to be included in the specified user constraint

**Related Object Classes:** `connection`, `node` and `unit_constraint`

re

**Related Parameters:** `connection_flow_coefficient`

`connection__from_node__unit_constraint` is a three-dimensional relationship between a `connection`, a `node` and a `unit_constraint`. The relationship specifies that the `connection_flow` variable

to the specified `connection` from the specified `node` is involved in the specified `unit_constraint`. Parameters on this relationship generally apply to this specific `connection_flow` variable. For example the parameter `connection_flow_coefficient` defined on `connection__from_node__unit_constraint` represents the coefficient on the specific `connection_flow` variable in the specified `unit_constraint`

## connection\_\_investment\_stochastic\_structure

re

Defines the stochastic structure of the connections investments variable

**Related Object Classes:** `connection` and `stochastic_structure`

The `connection__investment_stochastic_structure` relationship defines the `stochastic_structure` of `connection`-related investment decisions. Essentially, it sets the `stochastic_structure` used by the `connections_invested_available` variable of the `connection`.

The `connection__investment_stochastic_structure` relationship uses the `model_default_investment_stochastic_structure` relationship if not defined.

int

## connection\_\_investment\_temporal\_block

ck

Defines the temporal resolution of the connections investments variable

**Related Object Classes:** `connection` and `temporal_block`

block

`connection__investment_temporal_block` is a two-dimensional relationship between a `connection` and a `temporal_block`. This relationship defines the temporal resolution and scope of a connection's investment decision. Note that in a decomposed investments problem with two model objects, one for the master problem model and another for the operations problem model, the link to the specific model is made indirectly through the `model__temporal_block` relationship. If a `model_default_investment_temporal_block` is specified and no `connection__investment_temporal_block` relationship is specified, the `model_default_investment_temporal_block` relationship will be used. Conversely if `connection__investment_temporal_block` is specified along with `model__temporal_block`, this will override `model_default_investment_temporal_block` for the specified `connection`.

See also [Investment Optimization](#)



stochastic\_scenario

## connection\_\_node\_\_node

scenario

Holds parameters spanning multiple `connection_flow` variables to and from multiple `nodes`.

**Related Object Classes:** `connection` and `node`

re

**Related Parameters:** `compression_factor`, `connection_flow_delay`, `connection_linepack_constant`, `fix_ratio_out_in_connection_flow`, `fixed_pressure_constant_0`, `fixed_pressure_constant_1`, `max_ratio_out_in_connection_flow` and `min_ratio_out_in_connection_flow`

`connection__node__node` is a three-dimensional relationship between a `connection`, a `node` (node 1) and another `node` (node 2). `connection__node__node` infers a conversion and a direction with respect to that conversion. Node 1 is assumed to be the input node and node 2 is assumed to be the output node. For example, the `fix_ratio_out_in_connection_flow` parameter defined on `connection__node__node` relates the output `connection_flow` to node 2 to the input `connection_flow` from node 1

structure

## connection\_\_to\_node

t

Defines the `nodes` the `connection` can output to, and holds most `connection_flow` variable specific parameters.

c\_structure

block

**Related Object Classes:** `connection` and `node`

**Related Parameters:** `connection_capacity`, `connection_conv_cap_to_flow`, `connection_emergency_capacity`, `fix_connection_flow`, `fix_connection_intact_flow` and `graph_view_position`

re

`connection__to_node` is a two-dimensional relationship between a `connection` and a `node` and implies a `connection_flow` from the `connection` to the `node`. Specifying such a relationship will give rise to a `connection_flow_variable` with indices `connection=connection`, `node=node`, `direction=:to_node`. Relationships defined on this relationship will generally apply to this specific flow variable. For example, `connection_capacity` will apply only to this specific flow variable, unless the connection parameter `connection_type` is specified.

stochastic\_scenario

## connection\_\_to\_node\_\_unit\_constraint

scenario

when specified this relationship allows the relevant flow connection flow variable to be included in the specified user constraint

re

Related **Object Classes**: [connection](#), [node](#) and [unit\\_constraint](#)

Related **Parameters**: [connection\\_flow\\_coefficient](#)

`connection__to_node__unit_constraint` is a three-dimensional relationship between a [connection](#), a [node](#) and a [unit\\_constraint](#). The relationship specifies that the `connection_flow` variable from the specified [connection](#) to the specified [node](#) is involved in the specified [unit\\_constraint](#). Parameters on this relationship generally apply to this specific `connection_flow` variable. For example the parameter [connection\\_flow\\_coefficient](#) defined on `connection__to_node__unit_constraint` represents the coefficient on the specific `connection_flow` variable in the specified [unit\\_constraint](#)

int

## model\_\_default\_investment\_stochastic\_structure

structure

ck

Defines the default stochastic structure used for investment variables, which will be replaced by more specific definitions

t

c\_structure

Related **Object Classes**: [model](#) and [stochastic\\_structure](#)

block

The [model\\_default\\_investment\\_stochastic\\_structure](#) relationship can be used to set [model](#)-wide default [unit\\_investment\\_stochastic\\_structure](#), [connection\\_investment\\_stochastic\\_structure](#), and [node\\_investment\\_stochastic\\_structure](#) relationships. Its main purpose is to allow users to avoid defining each relationship individually, and instead allow them to focus on defining only the exceptions. As such, any specific [unit\\_investment\\_stochastic\\_structure](#), [connection\\_investment\\_stochastic\\_structure](#), and [node\\_investment\\_stochastic\\_structure](#) relationships take priority over the [model\\_default\\_investment\\_stochastic\\_structure](#) relationship.

## model\_\_default\_investment\_temporal\_block

Defines the default temporal block used for investment variables, which will be replaced by more specific definitions

**Related Object Classes:** [model](#) and [temporal\\_block](#)

[model\\_\\_default\\_investment\\_temporal\\_block](#) is a two-dimensional relationship between a [model](#) and a [temporal\\_block](#). This relationship defines the default temporal resolution and scope for all investment decisions in the model ([units](#), [connections](#) and storages). Specifying [model\\_\\_default\\_investment\\_temporal\\_block](#) for a model avoids the need to specify individual [node\\_investment\\_temporal\\_block](#), [unit\\_investment\\_temporal\\_block](#) and [connection\\_investment\\_temporal\\_block](#) relationships. Conversely, if any of these individual relationships are defined (e.g. [connection\\_investment\\_temporal\\_block](#)) along with [model\\_temporal\\_block](#), these will override [model\\_\\_default\\_investment\\_temporal\\_block](#).

See also [Investment Optimization](#)

## model\_\_default\_stochastic\_structure

Defines the default stochastic structure used for model variables, which will be replaced by more specific definitions

int

structure

**Related Object Classes:** [model](#) and [stochastic\\_structure](#)

The [model\\_\\_default\\_stochastic\\_structure](#) relationship can be used to set a [model](#)-wide default for the [node\\_stochastic\\_structure](#) and [units\\_on\\_stochastic\\_structure](#) relationships. Its main purpose is to allow users to avoid defining each relationship individually, and instead allow them to focus on defining only the exceptions. As such, any specific [node\\_stochastic\\_structure](#) or [units\\_on\\_stochastic\\_structure](#) relationships take priority over the [model\\_\\_default\\_stochastic\\_structure](#) relationship.

## model\_\_default\_temporal\_block

Defines the default temporal block used for model variables, which will be replaced by more specific definitions

re

**Related Object Classes:** [model](#) and [temporal\\_block](#)

The [model\\_\\_default\\_temporal\\_block](#) relationship can be used to set a [model](#)-wide default for the [node\\_temporal\\_block](#) and [units\\_on\\_temporal\\_block](#) relationships. Its main purpose is to allow users to

stochastic\_scenario

avoid defining each relationship individually, and instead allow them to focus on defining only the exceptions. As such, any specific `node__temporal_block` or `units_on__temporal_block` relationships take priority over the `model__default_temporal_block` relationship.

## model\_\_report

re

Determines which reports are written for each model and in turn, which outputs are written for each model

Related **Object Classes**: `model` and `report`

The `model__report` relationship tells which `reports` are written by which `model`, where the contents of the `reports` are defined separately using the `report__output` relationship. Without appropriately defined `model__report` and `report__output` and relationships, *SpineOpt* doesn't write any output, so be sure to include at least one `report` connected to all the `output variables` of interest in the `model`!

## model\_\_stochastic\_structure

int

structure

ck Defines which `stochastic_structures` are included in which `models`.

Related **Object Classes**: `model` and `stochastic_structure`

t

c\_structure

The `[model__stochastic_structure]` relationship defines which `stochastic_structures` are active in which `models`. Essentially, this relationship allows for e.g. attributing multiple `node__stochastic_structure` relationships for a single `node`, and switching between them in different `models`. Any `stochastic_structure` in the `model__default_stochastic_structure` relationship is automatically assumed to be active in the connected `model`, so there's no need to include it in `[model__stochastic_structure]` separately.

## model\_\_temporal\_block

re

Defines which `temporal_blocks` are included in which `models`.

Related **Object Classes**: `model` and `temporal_block`

stochastic\_scenario

The `model__temporal_block` relationship is used to determine which [temporal\\_blocks](#) are included in a specific model. Note that defining this relationship does not yet imply that any element of the model will be governed by the specified `temporal_block`, for this to happen additional relationships have to be defined such as the [model\\_default\\_temporal\\_block](#) relationship.

## node\_\_commodity

re

Define a `commodity` for a `node`. Only a single `commodity` is permitted per `node`

Related **Object Classes**: [commodity](#) and [node](#)

`node__commodity` is a two-dimensional relationship between a [node](#) and a [commodity](#) and specifies the commodity that `flows` to or from the node. Generally, since flows are not dimensioned by [commodity](#), this has no meaning in terms of the variables and constraint equations. However, there are two specific uses for this relationship:

1. To specify that specific network physics should apply to the network formed by the member nodes for that commodity. See [powerflow](#)
2. Only connection flows that are between nodes of the same or no [commodity](#) are included in the `node_balance` constraint.

int

structure

ck

## node\_\_investment\_stochastic\_structure

t

`node__investment_stochastic_structure` defines the stochastic structure for node related investments, currently only storages

block

Related **Object Classes**: [node](#) and [stochastic\\_structure](#)

The [node\\_\\_investment\\_stochastic\\_structure](#) relationship defines the [stochastic\\_structure](#) of `node`-related investment decisions. Essentially, it sets the [stochastic\\_structure](#) used by the [storages\\_invested\\_available](#) variable of the `node`.

The [node\\_\\_investment\\_stochastic\\_structure](#) relationship uses the [model\\_default\\_investment\\_stochastic\\_structure](#) relationship if not defined.

## node\_\_investment\_temporal\_block

defines the temporal resolution for node related investments, currently only storages

stochastic\_scenario

Related **Object Classes**: [node](#) and [temporal\\_block](#)

`node__investment_temporal_block` is a two-dimensional relationship between a [node](#) and a [temporal\\_block](#). This relationship defines the temporal resolution and scope of a [node](#)'s investment decisions (currently only storage investments). Note that in a decomposed investments problem with two model objects, one for the master problem model and another for the operations problem model, the link to the specific model is made indirectly through the [model\\_\\_temporal\\_block](#) relationship. If a [model\\_\\_default\\_investment\\_temporal\\_block](#) is specified and no `node__investment_temporal_block` relationship is specified, the [model\\_\\_default\\_investment\\_temporal\\_block](#) relationship will be used. Conversely if `node__investment_temporal_block` is specified along with [model\\_\\_temporal\\_block](#), this will override [model\\_\\_default\\_investment\\_temporal\\_block](#) for the specified [node](#).

See also [Investment Optimization](#)

## node\_\_node

Holds parameters for direct interactions between two [nodes](#), e.g. [node\\_state](#) diffusion coefficients.

ck

Related **Object Classes**: [node](#)

Related **Parameters**: [diff\\_coeff](#)

The [node\\_\\_node](#) relationship is used for defining direct interactions between two [nodes](#), like diffusion of [commodity](#) content. Note that the [node\\_\\_node](#) relationship is assumed to be one-directional, meaning that

```
node__node(node1=n1, node2=n2) != node__node(node1=n2, node2=n1).
```

Thus, when one wants to define *symmetric relationships* between two [nodes](#), one needs to define both directions as separate relationships.

## node\_\_stochastic\_structure

stochastic\_scenario

Defines which specific `stochastic_structure` is used by the `node` and all `flow` variables associated with it. Only one `stochastic_structure` is permitted per `node`.

**Related Object Classes:** `node` and `stochastic_structure`

re

The `node__stochastic_structure` relationship defines which `stochastic_structure` the `node` uses. Essentially, it sets the `stochastic_structure` of all the `flow` variables connected to the `node`, as well as the potential `node_state` variable. Note that only one `stochastic_structure` can be defined per `node` per `model`, as interpreted based on the `node__stochastic_structure` and `model__stochastic_structure` relationships. Investment variables use dedicated relationships, as detailed in the [Investment Optimization](#) section.

The `node__stochastic_structure` relationship uses the `model__default_stochastic_structure` relationship if not specified.

## node\_\_temporal\_block

int

Defines the `temporal_blocks` used by the `node` and all the `flow` variables associated with it.

ck

**Related Object Classes:** `node` and `temporal_block`

t

**Related Parameters:** `cyclic_condition`

block

This relationship links a `node` to a `temporal_block` and as such it will determine which temporal block governs the temporal horizon and resolution of the variables associated with this node. Specifically, the `resolution` of the temporal block will directly imply the duration of the time slices for which both the regular and ramping `flow` variables and their associated constraints are created.

For a more detailed description of how the temporal structure in SpineOpt can be created, see [Temporal Framework](#).

re

## node\_\_unit\_constraint

specifying this relationship allows a node's demand or `node_state` to be included in the specified unit constraint

stochastic\_scenario

**Related Object Classes:** [node](#) and [unit\\_constraint](#)

scenario

**Related Parameters:** [demand\\_coefficient](#) and [node\\_state\\_coefficient](#)

`node__unit_constraint` is a two-dimensional relationship between a [node](#) and a [unit\\_constraint](#). The relationship specifies that a variable associated only with the node (currently only the `node_state`) is involved in the constraint. For example, the [node\\_state\\_coefficient](#) defined on `node__unit_constraint` specifies the coefficient of the [node](#)'s `node_state` variable in the specified [unit\\_constraint](#).

See also [unit\\_constraint](#)

## parent\_stochastic\_scenario\_\_child\_stochastic\_scenario

Defines the master stochastic direct acyclic graph, meaning how the `stochastic_scenarios` are related to each other.

int

structure

ck

**Related Object Classes:** [stochastic\\_scenario](#)

The [parent\\_stochastic\\_scenario\\_\\_child\\_stochastic\\_scenario](#) relationship defines how the individual [stochastic\\_scenarios](#) are related to each other, forming what is referred to as the *stochastic direct acyclic graph (DAG)* in the [Stochastic Framework](#) section. It acts as a sort of basis for the [stochastic\\_structures](#), but doesn't contain any [Parameters](#) necessary for describing how it relates to the [Temporal Framework](#) or the [Objective function](#).

The [parent\\_stochastic\\_scenario\\_\\_child\\_stochastic\\_scenario](#) relationship and the *stochastic DAG* it forms are crucial for [Constraint generation with stochastic path indexing](#). Every finite *stochastic DAG* has a limited number of unique ways of traversing it, called *full stochastic paths*, which are used when determining how many different constraints need to be generated over time periods where [stochastic\\_structures](#) branch or converge, or when generating constraints involving different [stochastic\\_structures](#). See the [Stochastic Framework](#) section for more information.

re

## report\_\_output

Output object related to a report object are returned to the output database (if they appear in the model as variables)



stochastic\_scenario

Related **Object Classes**: [output](#) and [report](#)

scenario

The [report\\_output](#) relationship tells which [output variables](#) to include in which [report](#) when writing *SpineOpt* output. Note that the [reports](#) also need to be connected to a [model](#) using the [model\\_report](#) relationship. Without appropriately defined [model\\_report](#) and [report\\_output](#) and relationships, *SpineOpt* doesn't write any output, so be sure to include at least one [report](#) connected to all the [output variables](#) of interest in the [model](#)!

## stochastic\_structure\_\_stochastic\_scenario

Defines which [stochastic\\_scenarios](#) are included in which [stochastic\\_structure](#), and holds the parameters required for realizing the structure in combination with the [temporal\\_blocks](#).

Related **Object Classes**: [stochastic\\_scenario](#) and [stochastic\\_structure](#)

int

Related **Parameters**: [stochastic\\_scenario\\_end](#) and [weight\\_relative\\_to\\_parents](#)

ck

The [stochastic\\_structure\\_\\_stochastic\\_scenario](#) relationship defines which [stochastic\\_scenarios](#) are included in which [stochastic\\_structure](#), as well as holds the [stochastic\\_scenario\\_end](#) and [weight\\_relative\\_to\\_parents](#) Parameters defining how the [stochastic\\_structure](#) interacts with the [Temporal Framework](#) and the [Objective function](#). Along with [parent\\_stochastic\\_scenario\\_\\_child\\_stochastic\\_scenario](#), this relationship is used to define the exact properties of each [stochastic\\_structure](#), which are then applied to the [objects](#) describing the modelled system according to the [Structural relationship classes](#), like the [node\\_\\_stochastic\\_structure](#) relationship.

## unit\_\_commodity

Holds parameters for [commodities](#) used by the [unit](#).

re

Related **Object Classes**: [commodity](#) and [unit](#)

Related **Parameters**: [max\\_cum\\_in\\_unit\\_flow\\_bound](#)

To impose a limit on the cumulative amount of commodity flows, the `max_cum_in_unit_flow_bound` can be imposed on a `unit_commodity` relationship. This can be very helpful, e.g. if a certain amount of emissions should not be surpassed throughout the optimization.

Note that, next to the `unit_commodity` relationship, also the nodes connected to the units need to be associated with their corresponding commodities, see `node_commodity`.

## unit\_from\_node

Defines the nodes the unit can take input from, and holds most `unit_flow` variable specific parameters.

**Related Object Classes:** `node` and `unit`

**Related Parameters:** `fix_nonspin_ramp_up_unit_flow`, `fix_nonspin_units_started_up`, `fix_ramp_up_unit_flow`, `fix_start_up_unit_flow`, `fix_unit_flow_op`, `fix_unit_flow`, `fuel_cost`, `graph_view_position`, `max_res_shutdown_ramp`, `max_res_startup_ramp`, `max_shutdown_ramp`, `max_startup_ramp`, `min_res_shutdown_ramp`, `min_res_startup_ramp`, `min_shutdown_ramp`, `min_startup_ramp`, `minimum_operating_point`, `operating_points`, `ramp_down_cost`, `ramp_down_limit`, `ramp_up_cost`, `ramp_up_limit`, `reserve_procurement_cost`, `unit_capacity`, `unit_conv_cap_to_flow` and `vom_cost`

The `unit_to_node` and `unit_from_node` unit relationships are core elements of SpineOpt. For each `unit_to_node` or `unit_from_node`, a `unit_flow` variable is automatically added to the model, i.e. a commodity flow of a unit *to* or *from* a specific node, respectively.

Various parameters can be defined on the `unit_from_node` relationship, in order to constrain the associated unit flows. In most cases a `unit_capacity` will be defined for an upper bound on the commodity flows. Apart from that, ramping abilities of a unit can be defined. For further details on ramps see [Ramping and Reserves](#).

To associate costs with a certain commodity flows, cost terms, such as `fuel_costs` and `vom_costs`, can be included for the `unit_from_node` relationship.

It is important to note, that the parameters associated with the `unit_from_node` can be defined either for a specific `node`, or for a group of nodes. Grouping nodes for the described parameters will result in an aggregation of the unit flows for the triggered constraint, e.g. the definition of the `unit_capacity` on a group of nodes will result in an upper bound on the sum of all individual `unit_flows`.

stochastic\_scenario

## unit\_\_from\_node\_\_unit\_constraint

scenario

Defines which input `unit_flows` are included in the `unit_constraint`, and holds their parameters.

re

Related **Object Classes**: [node](#), [unit\\_constraint](#) and [unit](#)

Related **Parameters**: [graph\\_view\\_position](#) and [unit\\_flow\\_coefficient](#)

`unit__from_node__unit_constraint` is a three-dimensional relationship between a [unit](#), a [node](#) and a [unit\\_constraint](#). The relationship specifies that the `unit_flow` variable to the specified [unit](#) from the specified [node](#) is involved in the specified [unit\\_constraint](#). Parameters on this relationship generally apply to this specific `unit_flow` variable. For example the parameter [unit\\_flow\\_coefficient](#) defined on `unit__from_node__unit_constraint` represents the coefficient on the specific `unit_flow` variable in the specified [unit\\_constraint](#)

int

## unit\_\_investment\_stochastic\_structure

structure

ck

Sets the stochastic structure for investment decisions - overrides `model__default_investment_stochastic_structure`.

t

c\_structure

Related **Object Classes**: [stochastic\\_structure](#) and [unit](#)

block

The [unit\\_investment\\_stochastic\\_structure](#) relationship defines the [stochastic\\_structure](#) of [unit](#)-related investment decisions. Essentially, it sets the [stochastic\\_structure](#) used by the [units\\_invested\\_available](#) variable of the [unit](#).

The [unit\\_investment\\_stochastic\\_structure](#) relationship uses the [model\\_default\\_investment\\_stochastic\\_structure](#) relationship if not defined.

re

## unit\_\_investment\_temporal\_block

Sets the temporal resolution of investment decisions - overrides `model__default_investment_temporal_block`

stochastic\_scenario

**Related Object Classes:** [temporal\\_block](#) and [unit](#)

scenario

`unit__investment_temporal_block` is a two-dimensional relationship between a [unit](#) and a [temporal\\_block](#). This relationship defines the temporal resolution and scope of a unit's investment decision. Note that in a decomposed investments problem with two model objects, one for the master problem model and another for the operations problem model, the link to the specific model is made indirectly through the [model\\_\\_temporal\\_block](#) relationship. If a [model\\_\\_default\\_investment\\_temporal\\_block](#) is specified and no `unit__investment_temporal_block` relationship is specified, the [model\\_\\_default\\_investment\\_temporal\\_block](#) relationship will be used. Conversely if `unit__investment_temporal_block` is specified along with [model\\_\\_temporal\\_block](#), this will override [model\\_\\_default\\_investment\\_temporal\\_block](#) for the specified [unit](#).

See also [Investment Optimization](#)

## unit\_\_node\_\_node

Holds parameters spanning multiple `unit_flow` variables to and from multiple `nodes`.

int

**Related Object Classes:** [node](#) and [unit](#)

block

**Related Parameters:** [fix\\_ratio\\_in\\_in\\_unit\\_flow](#), [fix\\_ratio\\_in\\_out\\_unit\\_flow](#), [fix\\_ratio\\_out\\_in\\_unit\\_flow](#), [fix\\_ratio\\_out\\_out\\_unit\\_flow](#), [fix\\_units\\_on\\_coefficient\\_in\\_in](#), [fix\\_units\\_on\\_coefficient\\_in\\_out](#), [fix\\_units\\_on\\_coefficient\\_out\\_in](#), [fix\\_units\\_on\\_coefficient\\_out\\_out](#), [max\\_ratio\\_in\\_in\\_unit\\_flow](#), [max\\_ratio\\_in\\_out\\_unit\\_flow](#), [max\\_ratio\\_out\\_in\\_unit\\_flow](#), [max\\_ratio\\_out\\_out\\_unit\\_flow](#), [max\\_units\\_on\\_coefficient\\_in\\_in](#), [max\\_units\\_on\\_coefficient\\_in\\_out](#), [max\\_units\\_on\\_coefficient\\_out\\_in](#), [max\\_units\\_on\\_coefficient\\_out\\_out](#), [min\\_ratio\\_in\\_in\\_unit\\_flow](#), [min\\_ratio\\_in\\_out\\_unit\\_flow](#), [min\\_ratio\\_out\\_in\\_unit\\_flow](#), [min\\_ratio\\_out\\_out\\_unit\\_flow](#), [min\\_units\\_on\\_coefficient\\_in\\_in](#), [min\\_units\\_on\\_coefficient\\_in\\_out](#), [min\\_units\\_on\\_coefficient\\_out\\_in](#), [min\\_units\\_on\\_coefficient\\_out\\_out](#), [unit\\_idle\\_heat\\_rate](#), [unit\\_incremental\\_heat\\_rate](#) and [unit\\_start\\_flow](#)

While the relationships [unit\\_\\_to\\_node](#) and [unit\\_\\_to\\_node](#) take care of the automatic generation of the `unit_flow` variables, the [unit\\_\\_node\\_\\_node](#) relationships hold the information how the different commodity flows of a unit interact. Only through this relationship and the associated parameters, the topology of a unit, i.e. which intakes lead to which products etc., becomes unambiguous.

In almost all cases, at least one of the `..._ratio_...` parameters will be defined, e.g. to set a fixed ratio between outgoing and incoming commodity flows of unit (see also e.g. [fix\\_ratio\\_out\\_in\\_unit\\_flow](#)). Note that the parameters can also be defined on a relationship between groups of objects, e.g. to force a fixed

stochastic scenario

ratio between a group of nodes. In the triggered constraints, this will lead to an aggregation of the individual unit flows.

scenario

## unit\_\_to\_node

re

Defines the nodes the unit can output to, and holds most unit\_flow variable specific parameters.

Related **Object Classes**: [node](#) and [unit](#)

int

structure

ck

The [unit\\_\\_to\\_node](#) and [unit\\_\\_from\\_node](#) unit relationships are core elements of SpineOpt. For each [unit\\_\\_to\\_node](#) or [unit\\_\\_from\\_node](#), a [unit\\_flow](#) variable is automatically added to the model, i.e. a commodity flow of a unit *to* or *from* a specific node, respectively.

t

Various parameters can be defined on the [unit\\_\\_to\\_node](#) relationship, in order to constrain the associated unit flows. In most cases a [unit\\_capacity](#) will be defined for an upper bound on the commodity flows. Apart from that, ramping abilities of a unit can be defined. For further details on ramps see [Ramping and Reserves](#).

block

To associate costs with a certain commodity flow, cost terms, such as [fuel\\_costs](#) and [vom\\_costs](#), can be included for the [unit\\_\\_to\\_node](#) relationship.

It is important to note, that the parameters associated with the [unit\\_\\_to\\_node](#) can be defined either for a specific [node](#), or for a group of nodes. Grouping nodes for the described parameters will result in an aggregation of the unit flows for the triggered constraint, e.g. the definition of the [unit\\_capacity](#) on a group of nodes will result in an upper bound on the sum of all individual [unit\\_flows](#).

## unit\_\_to\_node\_\_unit\_constraint

stochastic\_scenario

Defines which output `unit_flows` are included in the `unit_constraint`, and holds their parameters.

scenario

**Related Object Classes:** [node](#), [unit\\_constraint](#) and [unit](#)

re

**Related Parameters:** [graph\\_view\\_position](#) and [unit\\_flow\\_coefficient](#)

`unit__to_node__unit_constraint` is a three-dimensional relationship between a [unit](#), a [node](#) and a [unit\\_constraint](#). The relationship specifies that the `unit_flow` variable from the specified [unit](#) to the specified [node](#) is involved in the specified [unit\\_constraint](#). Parameters on this relationship generally apply to this specific `unit_flow` variable. For example the parameter [unit\\_flow\\_coefficient](#) defined on `unit__to_node__unit_constraint` represents the coefficient on the specific `unit_flow` variable in the specified [unit\\_constraint](#)

## unit\_\_unit\_constraint

int

Defines which `units_on` variables are included in the `unit_constraint`, and holds their parameters.

structure

**Related Object Classes:** [unit\\_constraint](#) and [unit](#)

t

c\_structure

block

**Related Parameters:** [units\\_on\\_coefficient](#) and [units\\_started\\_up\\_coefficient](#)

`unit__unit_constraint` is a two-dimensional relationship between a [unit](#) and a [unit\\_constraint](#). The relationship specifies that a variable or variable(s) associated only with the unit (not a `unit_flow` for example) are involved in the constraint. For example, the [units\\_on\\_coefficient](#) defined on `unit__unit_constraint` specifies the coefficient of the [unit](#)'s `units_on` variable in the specified [unit\\_constraint](#).

re

See also [unit\\_constraint](#)

## units\_on\_\_stochastic\_structure

stochastic\_scenario

Defines which specific `stochastic_structure` is used for the `units_on` variable of the `unit`.

Only one `stochastic_structure` is permitted per `unit`.

**Related Object Classes:** [stochastic\\_structure](#) and [unit](#)

re

The [units\\_on\\_stochastic\\_structure](#) relationship defines the [stochastic\\_structure](#) used by the `units_on` variable. Essentially, this relationship permits defining a different [stochastic\\_structure](#) for the online decisions regarding the `units_on` variable, than what is used for the production [unit\\_flow](#) variables. A common use-case is e.g. using only one `units_on` variable across multiple [stochastic\\_scenarios](#) for the [unit\\_flow](#) variables. Note that only one [units\\_on\\_stochastic\\_structure](#) relationship can be defined per `unit` per `model`, as interpreted by the [units\\_on\\_stochastic\\_structure](#) and [model\\_stochastic\\_structure](#) relationships.

The [units\\_on\\_stochastic\\_structure](#) relationship uses the [model\\_default\\_stochastic\\_structure](#) relationship if not specified.

## units\_on\_\_temporal\_block

int

structure Defines which specific `temporal_blocks` are used by the `units_on` variable of the `unit`.

ck

**Related Object Classes:** [temporal\\_block](#) and [unit](#)

t

c\_structure

[units\\_on\\_\\_temporal\\_block](#) is a relationship linking the `units_on` variable of a unit to a specific [temporal\\_block](#) object. As such, this relationship will determine which temporal block governs the on- and offline status of the unit. The temporal block holds information on the temporal scope and resolution for which the variable should be optimized.

« [Object Classes](#)

[Parameters](#) »

re

# Parameters

## balance\_type

A selector for how the `:nodal_balance` constraint should be handled.

**Default value:** `balance_type_node`

Uses **Parameter Value Lists:** `balance_type_list`

**Related Object Classes:** `node`

The `balance_type` parameter determines whether or not a `node` needs to be balanced, in the classical sense that the sum of flows entering the `node` is equal to the sum of flows leaving it.

The values `balance_type_node` (the default) and `balance_type_group` mean that the `node` is always balanced. The only exception is if the `node` belongs in a group that has itself `balance_type` equal to `balance_type_group`. The value `balance_type_none` means that the `node` doesn't need to be balanced.

## big\_m

Sufficiently large number used for linearization bilinear terms, e.g. to enforce bidirectional flow for gas pipelines

**Default value:** `1000000`

**Related Object Classes:** `model`



The `big_m` parameter is a property of the `model` object. The `bigM` method is commonly used for the purpose of recasting non-linear constraints into a mixed-integer reformulation. In SpineOpt, the `bigM` formulation is used to describe the sign of gas flow through a connection (if a pressure driven gas transfer model is used). The `big_m` parameter in combination with the binary variable `binary_gas_connection_flow` is used in the constraints on `the gas flow capacity` and `the fixed node pressure points` and ensures that the average flow through a pipeline is only in one direction and is constraint by the fixed pressure points from the outer approximation of the Weymouth equation. See [Schwele - Coordination of Power and Natural Gas Systems: Convexification Approaches for Linepack Modeling](#) for reference.

## block\_end

The end time for the `temporal_block`. Can be given either as a `DateTime` for a static end point, or as a `Duration` for an end point relative to the start of the current optimization.

Related [Object Classes](#): `temporal_block`

Indicates the end of this temporal block. The default value is equal to a duration of 0. It is useful to distinguish here between two cases: a single solve, or a rolling window optimization.

**single solve** When a Date time value is chosen, this is directly the end of the optimization for this temporal block. In a single solve optimization, a combination of `block_start` and `block_end` can easily be used to run optimizations that cover only part of the model horizon. Multiple `temporal_block` objects can then be used to create optimizations for disconnected time periods, which is commonly used in the method of representative days. The default value coincides with the `model_end`.

**rolling window optimization** To create a temporal block that is rolling along with the optimization window, a rolling temporal block, a duration value should be chosen. The `block_end` parameter will in this case determine the size of the optimization window, with respect to the start of each optimization window. If multiple temporal blocks with different `block_end` parameters exist, the maximum value will determine the size of the optimization window. Note, this is different from the `roll_forward` parameter, which determines how much the window moves for after each optimization. For more info, see [One single temporal\\_block](#). The default value is equal to the `roll_forward` parameter.

## block\_start

The start time for the `temporal_block`. Can be given either as a `DateTime` for a static start point, or as a `Duration` for an start point relative to the start of the current optimization.

Related [Object Classes](#): [temporal\\_block](#)

Indicates the start of this temporal block. The main use of this parameter is to create an offset from the model start. The default value is equal to a duration of 0. It is useful to distinguish here between two cases: a single solve, or a rolling window optimization.

**single solve** When a Date time value is chosen, this is directly the start of the optimization for this temporal block. When a duration is chosen, it is added to the [model\\_start](#) to obtain the start of this [temporal\\_block](#). In the case of a duration, the chosen value directly marks the offset of the optimization with respect to the [model\\_start](#). The default value for this parameter is the [model\\_start](#).

**rolling window optimization** To create a temporal block that is rolling along with the optimization window, a rolling temporal block, a duration value should be chosen. The temporal [block\\_start](#) will again mark the offset of the optimization start but now with respect to the start of each optimization window.

## candidate\_connections

The number of connections that may be invested in

Related [Object Classes](#): [connection](#)

The [candidate\\_connections](#) parameter denotes the possibility of investing on a certain [connection](#).

The default value of `nothing` means that the [connection](#) can't be invested in, because it's already in operation. An integer value represents the maximum investment possible at any point in time, as a factor of the [connection\\_capacity](#).

In other words, [candidate\\_connections](#) is the upper bound of the [connections\\_invested\\_available](#) variable.

## candidate\_storages

Determines the maximum number of new storages which may be invested in

Related [Object Classes](#): [node](#)

Within an investments problem `candidate_storages` determines the upper bound on the storages investment decision variable in constraint `storages_invested_available`. In constraint

`node_state_cap` the maximum node state will be the product of the `storage_investment_variable_type` and `node_state_cap`. Thus, the interpretation of `candidate_storages` depends on `storage_investment_variable_type` which determines the investment decision variable type. If `storage_investment_variable_type` is integer or binary, then `candidate_storages` represents the maximum number of discrete storages of size `node_state_cap` that may be invested in at the corresponding node. If `storage_investment_variable_type` is continuous, `candidate_storages` is more analogous to a maximum storage capacity with `node_state_cap` being analogous to a scaling parameter.

Note that `candidate_storages` is the main investment switch and setting a value other than none/nothing triggers the creation of the investment variable for storages at the corresponding node. Note that a value of zero will still trigger the variable creation but its value will be fixed to zero. This can be useful if an inspection of the related dual variables will yield the value of this resource.

See also [Investment Optimization](#) and [storage\\_investment\\_variable\\_type](#)

## candidate\_units

Number of units which may be additionally constructed

Related [Object Classes](#): [unit](#)

Within an investments problem `candidate_units` determines the upper bound on the `unit` investment decision variable in constraint `units_invested_available`. In constraint `unit_flow_capacity` the maximum `unit_flow` will be the product of the `units_invested_available` and the corresponding `unit_capacity`. Thus, the interpretation of `candidate_units` depends on `unit_investment_variable_type` which determines the `unit` investment decision variable type. If `unit_investment_variable_type` is integer or binary, then `candidate_units` represents the maximum number of discrete units that may be invested in. If `unit_investment_variable_type` is continuous, `candidate_units` is more analogous to a maximum storage capacity.

Note that `candidate_units` is the main investment switch and setting a value other than none/nothing triggers the creation of the investment variable for the `unit`. Note that a value of zero will still trigger the variable creation but its value will be fixed to zero. This can be useful if an inspection of the related dual variables will yield the value of this resource.

See also [Investment Optimization](#) and [unit\\_investment\\_variable\\_type](#)

## commodity\_lodf\_tolerance

The minimum absolute value of the line outage distribution factor (LODF) that is considered meaningful.

**Default value:** 0.1

**Related Object Classes:** [commodity](#)

Given two [connections](#), the line outage distribution factor (LODF) is the fraction of the pre-contingency flow on the first one, that will flow on the second after the contingency. [commodity\\_lodf\\_tolerance](#) is the minimum absolute value of the LODF that is considered meaningful. Any value below this tolerance (in absolute value) will be treated as zero.

The LODFs are used to model contingencies on some [connections](#) and their impact on some other [connections](#). To model contingencies on a [connection](#), set [connection\\_contingency](#) to `true`; to study the impact of such contingencies on another [connection](#), set [connection\\_monitored](#) to `true`.

In addition, define a [commodity](#) with [commodity\\_physics](#) set to [commodity\\_physics\\_lodf](#), and associate that [commodity](#) (via [node\\_commodity](#)) to both [connections](#)' [nodes](#) (given by [connection\\_to\\_node](#) and [connection\\_from\\_node](#)).

## commodity\_physics

Defines if the [commodity](#) follows lodf or ptdf physics.

**Default value:** `commodity_physics_none`

**Uses Parameter Value Lists:** [commodity\\_physics\\_list](#)

**Related Object Classes:** [commodity](#)

This parameter determines the specific formulation used to carry out dc load flow within a model. To enable power transfer distribution factor (ptdf) based load flow for a network of [nodes](#) and [connections](#), all [nodes](#) must be related to a commodity with [commodity\\_physics](#) set to [commodity\\_physics\\_ptdf](#). To enable security constraint unit commitment based on ptdfs and line outage

distribution factors (lodf) all [nodes](#) must be related to a commodity with [commodity\\_physics](#) set to [commodity\\_physics\\_lodf](#).

See also [powerflow](#)

## commodity\_ptdf\_threshold

The minimum absolute value of the power transfer distribution factor (PTDF) that is considered meaningful.

**Default value:** 0.0001

**Related Object Classes:** [commodity](#)

Given a [connection](#) and a [node](#), the power transfer distribution factor (PTDF) is the fraction of the flow injected into the [node](#) that will flow on the [connection](#). [commodity\\_ptdf\\_threshold](#) is the minimum absolute value of the PTDF that is considered meaningful. Any value below this threshold (in absolute value) will be treated as zero.

The PTDFs are used to model DC power flow on certain [connections](#). To model DC power flow on a [connection](#), set [connection\\_monitored](#) to `true`.

In addition, define a [commodity](#) with [commodity\\_physics](#) set to either [commodity\\_physics\\_ptdf](#), or [commodity\\_physics\\_lodf](#). and associate that [commodity](#) (via [node\\_commodity](#)) to both [connections'](#) [nodes](#) (given by [connection\\_to\\_node](#) and [connection\\_from\\_node](#)).

## compression\_factor

The compression factor establishes a compression from an origin node to a receiving node, which are connected through a connection. The first node corresponds to the origin node, the second to the (compressed) destination node. Typically the value is  $\geq 1$ .

**Related Relationship Classes:** [connection\\_node\\_node](#)

This parameter is specific to the use of pressure driven gas transfer. To represent a compression between two nodes in the gas network, the [compression\\_factor](#) can be defined. This factor ensures that the pressure of a node is equal to (or lower than) the pressure at the sending node times the

compression\_factor. The relationship [connection\\_\\_node\\_\\_node](#) that hosts this parameter should be defined in a way that the first node represents the origin node and the second node represents the compressed node.

## connection\_availability\_factor

Availability of the `connection`, acting as a multiplier on its `connection_capacity`. Typically between 0-1.

Default value: 1.0

Related [Object Classes](#): [connection](#)

To indicate that a connection is only available to a certain extent or at certain times of the optimization, the [connection\\_availability\\_factor](#) can be used. A typical use case could be an availability timeseries for connection with expected outage times. By default the availability factor is set to 1. The availability is, among others, used in the [constraint\\_connection\\_flow\\_capacity](#).

## connection\_capacity

Limits the `connection_flow` variable to the `to_node`. `to_node` can be a group of nodes, in which case the sum of the `connection_flow` is constrained.

Related [Relationship Classes](#): [connection\\_\\_from\\_node](#) and [connection\\_\\_to\\_node](#)

Defines the upper bound on the corresponding `connection_flow` variable. If the connection is a candidate connection, the effective `connection_flow` upper bound is the product of the investment variable, `connections_invested_available` and `connection_capacity`. If ptdf based dc load flow is enabled, `connection_capacity` represents the normal rating of a `connection` (line) while [connection\\_emergency\\_capacity](#) represents the maximum post contingency flow.

## connection\_contingency

A boolean flag for defining a contingency `connection`.

Uses **Parameter Value Lists**: [boolean\\_value\\_list](#)

Related **Object Classes**: [connection](#)

Specifies that the connection in question is to be included as a contingency when security constrained unit commitment is enabled. When using security constrained unit commitment by setting [commodity\\_physics](#) to [commodity\\_physics\\_lodf](#), an N-1 security constraint is created for each monitored line (`connection_monitored = true`) for each specified contingency (`connection_contingency = true`).

See also [powerflow](#)

## connection\_conv\_cap\_to\_flow

Optional coefficient for `connection_capacity` unit conversions in the case the `connection_capacity` value is incompatible with the desired `connection_flow` units.

Default value: 1.0

Related **Relationship Classes**: [connection\\_from\\_node](#) and [connection\\_to\\_node](#)

The [connection\\_conv\\_cap\\_to\\_flow](#) can be used to perform the conversion between the measurement unit of the [connection\\_capacity](#) to the measurement unit of the [connection\\_flow](#) variable. The default of this parameter is 1, i.e. assuming that both are given in the same measurement unit.

## connection\_emergency\_capacity

The maximum post-contingency flow on a monitored `connection`.

Related **Relationship Classes**: [connection\\_from\\_node](#) and [connection\\_to\\_node](#)

The [connection\\_emergency\\_capacity](#) parameter represents the maximum post-contingency flow on a *monitored connection* if ptdf and lodf based security constrained unit commitment is enabled ([commodity\\_physics](#) is set to [`commodity_physics_lodf`]).

If you set this value, make sure that you also set [connection\\_monitored](#) to `true` for the involved [connection](#).

## connection\_flow\_coefficient

defines the unit constraint coefficient on the connection flow variable in the to direction

**Default value:** 0.0

**Related Relationship Classes:** [connection\\_from\\_node\\_unit\\_constraint](#) and [connection\\_to\\_node\\_unit\\_constraint](#)

The [connection\\_flow\\_coefficient](#) is an optional parameter that can be used to include the [connection\\_flow](#) variable from or to a [node](#) in a [unit\\_constraint](#) via the [connection\\_from\\_node\\_unit\\_constraint](#) and [connection\\_to\\_node\\_unit\\_constraint](#) relationships. Essentially, [connection\\_flow\\_coefficient](#) appears as a coefficient for the [connection\\_flow](#) variable from or to the [node](#) in the [unit\\_constraint](#).

## connection\_flow\_cost

Variable costs of a flow through a [connection](#). E.g. EUR/MWh of energy throughput.

**Related Object Classes:** [connection](#)

By defining the [connection\\_flow\\_cost](#) parameter for a specific [connection](#), a cost term will be added to the objective function that values all [connection\\_flow](#) variables associated with that connection during the current optimization window.

## connection\_flow\_delay

Delays the [connection\\_flows](#) associated with the latter [node](#) in respect to the [connection\\_flows](#) associated with the first [node](#).



**Default value:** Dict{String, Any}{"data" => "0h", "type" => "duration"}

**Related Relationship Classes:** [connection\\_\\_node\\_\\_node](#)

The [connection\\_flow\\_delay](#) parameter denotes the amount of time that it takes for the flow to go through a [connection](#). In other words, the flow that enters the [connection](#) is only seen at the other side after [connection\\_flow\\_delay](#) units of time.

## connection\_investment\_cost

The per unit investment cost for the connection over the [connection\\_investment\\_lifetime](#)

**Related Object Classes:** [connection](#)

By defining the [connection\\_investment\\_cost](#) parameter for a specific [connection](#), a cost term will be added to the objective function whenever a connection investment is made during the current optimization window.

## connection\_investment\_lifetime

Determines the minimum investment lifetime of a connection. Once invested, it remains in service for this long

**Related Object Classes:** [connection](#)

[connection\\_investment\\_lifetime](#) is the minimum amount of time that a [connection](#) has to stay in operation once it's invested-in. Only after that time, the [connection](#) can be decomissioned. Note that [connection\\_investment\\_lifetime](#) is a dynamic parameter that will impact the amount of solution history that must remain available to the optimisation in each step - this may impact performance.

## connection\_investment\_variable\_type

Determines whether the investment variable is integer `variable_type_integer` or continuous `variable_type_continuous`

**Default value:** `variable_type_integer`

**Uses [Parameter Value Lists](#):** `variable_type_list`

**Related [Object Classes](#):** `connection`

The `connection_investment_variable_type` parameter represents the *type* of the `connections_invested_available` decision variable.

The default value, `variable_type_integer`, means that only integer factors of the `connection_capacity` can be invested in. The value `variable_type_continuous` means that any fractional factor can also be invested in. The value `variable_type_binary` means that only a factor of 1 or zero are possible.

## connection\_linepack\_constant

The linepack constant is a property of gas pipelines and relates the linepack to the pressure of the adjacent nodes.

**Related [Relationship Classes](#):** `connection__node__node`

The linepack constant is a physical property of a connection representing a pipeline and holds information on how the linepack flexibility relates to pressures of the adjacent nodes. If, and only if, this parameter is defined, the linepack flexibility of a pipeline can be modelled. The existence of the parameter triggers the generation of the [constraint on line pack storage](#). The `connection_linepack_constant` should always be defined on the tuple (connection pipeline, linepack storage node, node group (containing both pressure nodes, i.e. start and end of the pipeline)). [See also](#).

## connection\_monitored

A boolean flag for defining a contingency `connection`.

Default value: false

Uses [Parameter Value Lists](#): [boolean\\_value\\_list](#)

Related [Object Classes](#): [connection](#)

When using ptdf-based load flow by setting [commodity\\_physics](#) to either [commodity\\_physics\\_ptdf](#) or [commodity\\_physics\\_ptdf](#), a constraint is created for each connection for which `connection_monitored = true`. Thus, to monitor the ptdf-based flow on a particular connection `connection_monitored` must be set to `true`.

See also [powerflow](#)

## connection\_reactance

The per unit reactance of a `connection`.

Related [Object Classes](#): [connection](#)

The per unit reactance of a transmission line. Used in ptdf based dc load flow where the relative reactances of lines determine the ptdfs of the network and in lossless dc powerflow where the flow on a line is given by  $\text{flow} = 1/x(\theta_{\text{to}} - \theta_{\text{from}})$  where  $x$  is the reactance of the line,  $\theta_{\text{to}}$  is the voltage angle of the remote node and  $\theta_{\text{from}}$  is the voltage angle of the sending node.

## connection\_reactance\_base

If the reactance is given for a p.u. (e.g. p.u. = 100MW), the `connection_reactance_base` can be set to perform this conversion (e.g. \*100).

Default value: 1

Related [Object Classes](#): [connection](#)

As the [connection\\_reactance](#) is often given on a per unit basis, often different than the units used elsewhere, the [connection\\_reactance\\_base](#) parameter serves as a conversion factor, scaling the [connection\\_reactance](#) with its p.u..

## connection\_resistance

The per unit resistance of a [connection](#).

Related [Object Classes](#): [connection](#)

The per unit resistance of a transmission line. **Currently unimplemented!**

## connection\_type

A selector between a normal and a lossless bidirectional [connection](#).

**Default value:** `connection_type_normal`

Uses [Parameter Value Lists](#): [connection\\_type\\_list](#)

Related [Object Classes](#): [connection](#)

Used to control specific pre-processing actions on connections. Currently, the primary purpose of [connection\\_type](#) is to simplify the data that is required to define a simple bi-directional, lossless line. If [connection\\_type](#)=`:connection_type_lossless_bidirectional`, it is only necessary to specify the following minimum data:

- relationship: [connection\\_\\_from\\_node](#)
- relationship: [connection\\_\\_to\\_node](#)
- parameter: [connection\\_capacity](#) (defined on [connection\\_\\_from\\_node](#) and/or [connection\\_\\_to\\_node](#))

If [connection\\_type](#)=`:connection_type_lossless_bidirectional` the following pre-processing actions are taken:

- reciprocal [connection\\_from\\_node](#) and [connection\\_to\\_node](#) relationships are created if they don't exist
- a new [connection\\_node\\_node](#) relationship is created if none exists already
- [fix\\_ratio\\_out\\_in\\_connection\\_flow](#) parameter is created with the value of 1 if no existing parameter found (therefore this value can be overridden)
- The first [connection\\_capacity](#) parameter found is copied to [connection\\_from\\_nodes](#) and [connection\\_to\\_nodes](#) without a defined [connection\\_capacity](#).

## constraint\_sense

A selector for the sense of the [unit\\_constraint](#).

Default value: ==

Uses [Parameter Value Lists](#): [constraint\\_sense\\_list](#)

Related [Object Classes](#): [unit\\_constraint](#)

The [constraint\\_sense](#) parameter determines the *sense* of a custom user constraint.

See [User constraints](#) for details.

## curtailment\_cost

Costs for curtailing generation. Essentially, accrues costs whenever [unit\\_flow](#) not operating at its maximum available capacity. E.g. EUR/MWh

Related [Object Classes](#): [unit](#)

By defining the [curtailment\\_cost](#) parameter for a specific [unit](#), a cost term will be added to the objective function whenever this unit's available capacity exceeds its activity (i.e., the [unit\\_flow](#) variable) over the course of the operational dispatch during the current optimization window.

## cyclic\_condition

If the cyclic condition is set to true for a storage node, the `node_state` at the end of the optimization window has to be larger than or equal to the initial storage state.

**Default value:** false

**Uses [Parameter Value Lists](#):** `boolean_value_list`

**Related [Relationship Classes](#):** `node__temporal_block`

The `cyclic_condition` parameter is used to enforce that the storage level at the end of the optimization window is higher or equal to the storage level at the beginning optimization. If the `cyclic_condition` parameter is set to `true` for a `node__temporal_block` relationship, and the `has_state` parameter of the corresponding `node` is set to `true`, the `constraint_cyclic_node_state` will be triggered.

## demand

Demand for the `commodity` of a `node`. Energy gains can be represented using negative `demand`.

**Default value:** 0.0

**Related [Object Classes](#):** `node`

The `demand` parameter represents a "demand" or a "load" of a `commodity` on a `node`. It appears in the `node injection constraint`, with positive values interpreted as "demand" or "load" for the modelled system, while negative values provide the system with "influx" or "gain". When the node is part of a group, the `fractional_demand` parameter can be used to split `demand` into fractions, when desired. See also: [Introduction to groups of objects](#)

The `demand` parameter can also be included in custom `unit_constraints` using the `demand_coefficient` parameter for the `node__unit_constraint` relationship.

## demand\_coefficient

coefficient of the specified node's demand in the specified unit constraint

Default value: 0.0

Related [Relationship Classes](#): [node\\_\\_unit\\_constraint](#)

The [demand\\_coefficient](#) is an optional parameter that can be used to include the [demand](#) of the a [node](#) in a [unit\\_constraint](#) via the [node\\_\\_unit\\_constraint](#) relationship. Essentially, [demand\\_coefficient](#) appears as a coefficient for the [demand](#) parameter of the connected [node](#) in the [unit constraint](#).

## diff\_coeff

Commodity diffusion coefficient between two [nodes](#). Effectively, denotes the *diffusion power per unit of state* from the first [node](#) to the second.

Default value: 0.0

Related [Relationship Classes](#): [node\\_\\_node](#)

The [diff\\_coeff](#) parameter represents diffusion of a [commodity](#) between the two [nodes](#) in the [node\\_\\_node](#) relationship. It appears as a coefficient on the [node\\_state](#) variable in the [node injection](#) constraint, essentially representing *diffusion power per unit of state*. Note that the [diff\\_coeff](#) is interpreted as *one-directional*, meaning that if one defines

```
diff_coeff(node1=n1, node2=n2),
```

there will only be diffusion from [n1](#) to [n2](#), but not vice versa. *Symmetric diffusion* is likely used in most cases, requiring defining the [diff\\_coeff](#) both ways

```
diff_coeff(node1=n1, node2=n2) == diff_coeff(node1=n2, node2=n1).
```

## downward\_reserve

Identifier for nodes providing downward reserves

**Default value:** false

**Related Object Classes:** [node](#)

If a [node](#) has a `true` [is\\_reserve\\_node](#) parameter, it will be treated as a reserve node in the model. To define whether the node corresponds to an upward or downward reserve commodity, the [upward\\_reserve](#) or the [downward\\_reserve](#) parameter needs to be set to true, respectively.

## duration\_unit

Defines the base temporal unit of the `model`. Currently supported values are either an `hour` or a `minute`.

**Default value:** minute

**Uses Parameter Value Lists:** [duration\\_unit\\_list](#)

**Related Object Classes:** [model](#)

The [duration\\_unit](#) parameter specifies the base unit of time in a [model](#). Two values are currently supported, `hour` and the default `minute`. E.g. if the [duration\\_unit](#) is set to `hour`, a `Duration` of one `minute` gets converted into `1/60 hours` for the calculations.

## fix\_binary\_gas\_connection\_flow

Fix the value of the `connection_flow_binary` variable, and hence pre-determine the direction of flow in the connection.

**Related Relationship Classes:** [connection\\_\\_from\\_node](#)



The binary flow of a gas pipelines for pressure driven gas transfer is enabled through the binary variable [binary\\_gas\\_connection\\_flow](#) and the [big\\_m](#) constant. To fix this binary variable, i.e. pre-define the direction of gas through the pipelines, the [fix\\_binary\\_gas\\_connection\\_flow](#) parameter can be used.

## fix\_connection\_flow

Fix the value of the [connection\\_flow](#) variable.

Related [Relationship Classes](#): [connection\\_\\_from\\_node](#) and [connection\\_\\_to\\_node](#)

The [fix\\_connection\\_flow](#) parameter fixes the value of the [connection\\_flow](#) variable.

## fix\_connection\_intact\_flow

Fix the value of the [connection\\_intact\\_flow](#) variable.

Related [Relationship Classes](#): [connection\\_\\_from\\_node](#) and [connection\\_\\_to\\_node](#)

The [fix\\_connection\\_intact\\_flow](#) parameter can be used to fix the values of the [connection\\_intact\\_flow](#) variable to preset values. If set to a [Scalar](#) type value, the [connection\\_intact\\_flow](#) variable is fixed to that value for all time steps and [stochastic\\_scenarios](#). Values for individual time steps can be fixed using [TimeSeries](#) type values.

## fix\_connections\_invested

Setting a value fixes the [connections\\_invested](#) variable accordingly

Related [Object Classes](#): [connection](#)

The [fix\\_connections\\_invested](#) parameter can be used to fix the values of the [connections\\_invested](#) variable to preset values. If set to a [Scalar](#) type value, the [connections\\_invested](#) variable is fixed to that value for all time steps and [stochastic\\_scenarios](#). Values for individual time steps can be fixed using [TimeSeries](#) type values.

See [Investment Optimization](#) for more information about the investment framework in *SpineOpt.jl*.

## fix\_connections\_invested\_available

Setting a value fixes the connections *invested* available variable accordingly

Related [Object Classes](#): [connection](#)

The [fix\\_connections\\_invested\\_available](#) parameter represents a *forced* connection investment.

In other words, it is the fix value of the [connections\\_invested\\_available](#) variable.

## fix\_node\_pressure

Fixes the corresponding `node_pressure` variable to the provided value

Related [Object Classes](#): [node](#)

In a pressure driven gas model, gas network nodes are associated with the [node\\_pressure](#) variable. In order to fix the pressure at a certain node or to give initial conditions the [fix\\_node\\_pressure](#) parameter can be used.

## fix\_node\_state

Fixes the corresponding `node_state` variable to the provided value. Can be used for e.g. fixing boundary conditions.

Related [Object Classes](#): [node](#)

The [fix\\_node\\_state](#) parameter simply fixes the value of the `node_state` variable to the provided value, if one is found. Common uses for the parameter include e.g. providing initial values for `node_state` variables, by fixing the value on the first modelled time step (*or the value before the first modelled time step*) using a `TimeSeries` type parameter value with an appropriate timestamp. Due to the way *SpineOpt* handles `TimeSeries` data, the `node_state` variables are only fixed for time steps with defined [fix\\_node\\_state](#) parameter values.

# fix\_node\_voltage\_angle

Fixes the corresponding `node_voltage_angle` variable to the provided value

Related **Object Classes**: [node](#)

For a lossless nodal DC power flow network, each node is associated with a [node\\_voltage\\_angle](#) variable. In order to fix the voltage angle at a certain node or to give initial conditions the [fix\\_node\\_voltage\\_angle](#) parameter can be used.

# fix\_nonspin\_ramp\_down\_unit\_flow

Fix the `nonspin_ramp_down_unit_flow` variable.

Related **Relationship Classes**: [unit\\_to\\_node](#)

The `fix_nonspin_ramp_down_unit_flow` parameter simply fixes the value of the [nonspin\\_ramp\\_down\\_unit\\_flow](#) variable to the provided value. As such, it determines directly how much non-spinning downward reserve commodity flows the relevant [unit](#) is providing to the [node](#) to which it is linked by the [unit\\_to\\_node](#) relationship.

When a single value is selected, this value is kept constant throughout the model. It is also possible to provide a timeseries of values, which can be used for example to impose initial conditions by providing a value only for the first timestep included in the model.

# fix\_nonspin\_ramp\_up\_unit\_flow

Fix the `nonspin_ramp_up_unit_flow` variable.

Related **Relationship Classes**: [unit\\_from\\_node](#) and [unit\\_to\\_node](#)

The `fix_nonspin_ramp_up_unit_flow` parameter simply fixes the value of the [nonspin\\_ramp\\_up\\_unit\\_flow](#) variable to the provided value. As such, it determines directly how much non-spinning upward reserve commodity flows the relevant [unit](#) is providing to the [node](#) to which it is linked by the [unit\\_to\\_node](#) relationship.

When a single value is selected, this value is kept constant throughout the model. It is also possible to provide a timeseries of values, which can be used for example to impose initial conditions by providing a value only for the first timestep included in the model.

## fix\_nonspin\_units\_shut\_down

Fix the `nonspin_units_shut_down` variable.

Related [Relationship Classes](#): [unit\\_to\\_node](#)

The `fix_nonspin_units_shut_down` parameter simply fixes the value of the [nonspin\\_units\\_shut\\_down](#) variable to the provided value. As such, it determines directly how many member [units](#) are involved in providing downward reserve commodity flows to the [node](#) to which it is linked by the [unit\\_to\\_node](#) relationship.

When a single value is selected, this value is kept constant throughout the model. It is also possible to provide a timeseries of values, which can be used for example to impose initial conditions by providing a value only for the first timestep included in the model.

## fix\_nonspin\_units\_started\_up

Fix the `nonspin_units_started_up` variable.

Related [Relationship Classes](#): [unit\\_from\\_node](#) and [unit\\_to\\_node](#)

The `fix_nonspin_units_started_up` parameter simply fixes the value of the [nonspin\\_units\\_started\\_up](#) variable to the provided value. As such, it determines directly how many member [units](#) are involved in providing upward reserve commodity flows to the [node](#) to which it is linked by the [unit\\_to\\_node](#) relationship.

When a single value is selected, this value is kept constant throughout the model. It is also possible to provide a timeseries of values, which can be used for example to impose initial conditions by providing a value only for the first timestep included in the model.

## fix\_ramp\_down\_unit\_flow

Fix the `ramp_down_unit_flow` variable.

Related [Relationship Classes](#): [unit\\_to\\_node](#)

The `fix_ramp_down_unit_flow` parameter simply fixes the value of the [ramp\\_down\\_unit\\_flow](#) variable to the provided value. It is possible to provide an incomplete timeseries of values, which can be used for example to impose initial conditions by providing a value only for the first timestep included in the model.

## `fix_ramp_up_unit_flow`

Fix the `ramp_up_unit_flow` variable.

Related [Relationship Classes](#): [unit\\_from\\_node](#) and [unit\\_to\\_node](#)

The `fix_ramp_up_unit_flow` parameter simply fixes the value of the [ramp\\_up\\_unit\\_flow](#) variable to the provided value. It is possible to provide an incomplete timeseries of values, which can be used for example to impose initial conditions by providing a value only for the first timestep included in the model.

## `fix_ratio_in_in_unit_flow`

Fix the ratio between two `unit_flows` coming into the `unit` from the two `nodes`.

Related [Relationship Classes](#): [unit\\_node\\_node](#)

The definition of the [fix\\_ratio\\_in\\_in\\_unit\\_flow](#) parameter triggers the generation of the [constraint\\_fix\\_ratio\\_in\\_in\\_unit\\_flow](#) and fixes the ratio between incoming flows of a unit. The parameter is defined on the relationship class [unit\\_node\\_node](#), where both nodes (or group of nodes) in this relationship represent `from_nodes`, i.e. the incoming flows to the unit. The ratio parameter is interpreted such that it constrains the ratio of `in1` over `in2`, where `in1` is the [unit\\_flow](#) variable from the first [node](#) in the [unit\\_node\\_node](#) relationship in a left-to-right order. This parameter can be useful, for instance if a unit requires a specific commodity mix as a fuel supply.

To enforce e.g. for a unit `u` a fixed share of `0.8` of its incoming flow from the node `supply_fuel_1` compared to its incoming flow from the node group `supply_fuel_2` (consisting of the two nodes

supply\_fuel\_2\_component\_a and supply\_fuel\_2\_component\_b) the [fix\\_ratio\\_in\\_in\\_unit\\_flow](#) parameter would be set to 0.8 for the relationship u\_\_supply\_fuel\_1\_\_supply\_fuel\_2.

## fix\_ratio\_in\_out\_unit\_flow

Fix the ratio between an incoming [unit\\_flow](#) from the first [node](#) and an outgoing [unit\\_flow](#) to the second [node](#).

Related [Relationship Classes](#): [unit\\_node\\_node](#)

The definition of the [fix\\_ratio\\_in\\_out\\_unit\\_flow](#) parameter triggers the generation of the [constraint\\_fix\\_ratio\\_in\\_out\\_unit\\_flow](#) and fixes the ratio between incoming and outgoing flows of a unit. The parameter is defined on the relationship class [unit\\_node\\_node](#), where the first node (or group of nodes) in this relationship represents the [from\\_node](#), i.e. the incoming flows to the unit, and the second node (or group of nodes), represents the [to\\_node](#) i.e. the outgoing flow from the unit. The ratio parameter is interpreted such that it constrains the ratio of [in](#) over [out](#), where [in](#) is the [unit\\_flow](#) variable from the first [node](#) in the [unit\\_node\\_node](#) relationship in a left-to-right order.

To enforce e.g. a fixed ratio of 1.4 for a unit [u](#) between its incoming gas flow from the node [ng](#) and its outgoing flows to the node group [e1\\_heat](#) (consisting of the two nodes [e1](#) and [heat](#)), the [fix\\_ratio\\_in\\_out\\_unit\\_flow](#) parameter would be set to 1.4 for the relationship u\_\_ng\_\_e1\_heat.

## fix\_ratio\_out\_in\_connection\_flow

Fix the ratio between the [connection\\_flow](#) from the first [node](#) and the [connection\\_flow](#) to the second [node](#).

Related [Relationship Classes](#): [connection\\_node\\_node](#)

The definition of the [fix\\_ratio\\_out\\_in\\_connection\\_flow](#) parameter triggers the generation of the [constraint\\_fix\\_ratio\\_out\\_in\\_connection\\_flow](#) and fixes the ratio between outgoing and incoming flows of a connection. The parameter is defined on the relationship class [connection\\_node\\_node](#), where the first node (or group of nodes) in this relationship represents the [to\\_node](#), i.e. the outgoing flow from the [connection](#), and the second node (or group of nodes), represents the [from\\_node](#), i.e. the incoming flows to the [connection](#). In most cases the [fix\\_ratio\\_out\\_in\\_connection\\_flow](#) parameter is set to equal or lower than 1, linking the flows entering to the flows leaving the connection. The ratio parameter is interpreted such that it constrains the ratio of [out](#) over [in](#), where [out](#) is the [connection\\_flow](#) variable

from the first `node` in the `connection__node__node` relationship in a left-to-right order. The parameter can be used to e.g. account for losses over a connection in a certain direction.

To enforce e.g. a fixed ratio of `0.8` for a connection `conn` between its outgoing electricity flow to `node` `e11` and its incoming flows from the node `node` `e12`, the `fix_ratio_out_in_connection_flow` parameter would be set to `0.8` for the relationship `u__e11__e12`.

## `fix_ratio_out_in_unit_flow`

Fix the ratio between an outgoing `unit_flow` to the first `node` and an incoming `unit_flow` from the second `node`.

Related [Relationship Classes](#): `unit__node__node`

The definition of the `fix_ratio_out_in_unit_flow` parameter triggers the generation of the `constraint_fix_ratio_out_in_unit_flow` and fixes the ratio between out and incoming flows of a `unit`. The parameter is defined on the relationship class `unit__node__node`, where the first node (or group of nodes) in this relationship represents the `to_node`, i.e. the outgoing flow from the unit, and the second node (or group of nodes), represents the `from_node`, i.e. the incoming flows to the unit. The ratio parameter is interpreted such that it constrains the ratio of `out` over `in`, where `out` is the `unit_flow` variable from the first `node` in the `unit__node__node` relationship in a left-to-right order.

To enforce e.g. a fixed ratio of `0.8` for a unit `u` between its outgoing flows to the node group `e1_heat` (consisting of the two nodes `e1` and `heat`) and its incoming gas flow from `ng` the `fix_ratio_out_in_unit_flow` parameter would be set to `0.8` for the relationship `u__e1_heat__ng`.

## `fix_ratio_out_out_unit_flow`

Fix the ratio between two `unit_flows` going from the `unit` into the two `nodes`.

Related [Relationship Classes](#): `unit__node__node`

The definition of the `fix_ratio_out_out_unit_flow` parameter triggers the generation of the `constraint_fix_ratio_out_out_unit_flow` and fixes the ratio between outgoing flows of a unit. The parameter is defined on the relationship class `unit__node__node`, where the nodes (or group of nodes) in this relationship represent the `to_node`'s, i.e. outgoing flow from the unit. The ratio parameter is interpreted such that it constrains the ratio of `out1` over `out2`, where `out1` is the `unit_flow` variable from the first `node` in the `unit__node__node` relationship in a left-to-right reading order.

To enforce a fixed ratio between two products of a unit `u`, e.g. fixing the share of produced electricity flowing to node `e1` to `0.4` of the production of heat flowing to node `heat`, the `fix_ratio_out_out_unit_flow` parameter would be set to `0.4` for the relationship `u__e1__heat`.

## fix\_shut\_down\_unit\_flow

Fix the `shut_down_unit_flow` variable.

Related [Relationship Classes](#): `unit__to_node`

The `fix_shut_down_unit_flow` parameter fixes the value of the `shut_down_unit_flow` to the provided value, if the parameter is defined.

Common uses for the parameter include e.g. providing initial values for the `shut_down_unit_flow`, by fixing the value on the first modelled time step (*or the value before the first modelled time step*) using a `TimeSeries` type parameter value with an appropriate timestamp. Due to the way *SpineOpt* handles `TimeSeries` data, the `shut_down_unit_flow` variable is only fixed for time steps with defined `fix_shut_down_unit_flow` parameter values.

Other uses can include e.g. a constant or time-varying **exogenous** commodity flow from or to a unit.

Note that the mentioned `shut_down_unit_flow` variable is only included if the parameter `max_startup_ramp` exist for the correspond `unit__to_node` or `unit__from_node` relationship. The usage of ramps is described in [Ramping and Reserves](#).

## fix\_start\_up\_unit\_flow

Fix the `start_up_unit_flow` variable.

Related [Relationship Classes](#): `unit__from_node` and `unit__to_node`

The `fix_start_up_unit_flow` parameter fixes the value of the `start_up_unit_flow` to the provided value, if the parameter is defined.

Common uses for the parameter include e.g. providing initial values for the `start_up_unit_flow`, by fixing the value on the first modelled time step (*or the value before the first modelled time step*) using a `TimeSeries` type parameter value with an appropriate timestamp. Due to the way *SpineOpt* handles `TimeSeries` data, the `start_up_unit_flow` variable is only fixed for time steps with defined `fix_start_up_unit_flow` parameter values.



Other uses can include e.g. a constant or time-varying **exogenous** commodity flow from or to a unit.

Note that the mentioned [start\\_up\\_unit\\_flow](#) variable is only included if the parameter [max\\_startup\\_ramp](#) exist for the correspond [unit\\_to\\_node](#) or [unit\\_from\\_node](#) relationship. The usage of ramps is described in [Ramping and Reserves](#).

## fix\_storages\_invested

Used to fix the value of the storages\_invested variable

Related [Object Classes](#): [node](#)

Used primarily to fix the value of the [storages\\_invested](#) variable which represents the point-in-time storage investment decision variable at a [node](#) and how many candidate storages are invested-in in a particular timeslice at the corresponding [node](#).

See also [Investment Optimization](#), [candidate\\_storages](#) and [storage\\_investment\\_variable\\_type](#)

## fix\_storages\_invested\_available

Used to fix the value of the storages\_invested\_available variable

Related [Object Classes](#): [node](#)

Used primarily to fix the value of the [storages\\_invested\\_available](#) variable which represents the storages investment decision variable and how many candidate storages are available at the corresponding node, time step and stochastic scenario. Used also in the decomposition framework to communicate the value of the master problem solution variables to the operational sub-problem.

See also [candidate\\_storages](#) and [Investment Optimization](#)

## fix\_unit\_flow

Fix the [unit\\_flow](#) variable.

Related [Relationship Classes](#): [unit\\_from\\_node](#) and [unit\\_to\\_node](#)

The `fix_unit_flow` parameter fixes the value of the `unit_flow` variable to the provided value, if the parameter is defined.

Common uses for the parameter include e.g. providing initial values for the `unit_flow` variable, by fixing the value on the first modelled time step (*or the value before the first modelled time step*) using a `TimeSeries` type parameter value with an appropriate timestamp. Due to the way *SpineOpt* handles `TimeSeries` data, the `unit_flow` variable is only fixed for time steps with defined `fix_unit_flow` parameter values.

Other uses can include e.g. a constant or time-varying **exogenous** commodity flow from or to a unit.

## fix\_unit\_flow\_op

Fix the `unit_flow_op` variable.

Related **Relationship Classes**: `unit_from_node` and `unit_to_node`

If `operating_points` is defined on a certain `unit_to_node` or `unit_from_node` flow, the corresponding `unit_flow` flow variable is decomposed into a number of sub-variables, `unit_flow_op` one for each operating point, with an additional index, `i` to reference the specific operating point. `fix_unit_flow_op` can thus be used to fix the value of one or more of the variables as desired.

## fix\_units\_invested

Fix the value of the `units_invested` variable.

Related **Object Classes**: `unit`

Used primarily to fix the value of the `units_invested` variable which represents the point-in-time `unit` investment decision variable and how many candidate units are invested-in in a particular timeslice.

See also [Investment Optimization](#), `candidate_units` and `unit_investment_variable_type`

## fix\_units\_invested\_available

Fix the value of the `units_invested_available` variable

Related [Object Classes](#): [unit](#)

Used primarily to fix the value of the `units_invested_available` variable which represents the unit investment decision variable and how many candidate units are invested-in and available at the corresponding node, time step and stochastic scenario. Used also in the decomposition framework to communicate the value of the master problem solution variables to the operational sub-problem.

See also [Investment Optimization](#), [candidate\\_units](#) and [unit\\_investment\\_variable\\_type](#)

## fix\_units\_on

Fix the value of the `units_on` variable.

Related [Object Classes](#): [unit](#)

The `fix_units_on` parameter simply fixes the value of the [units\\_on](#) variable to the provided value. As such, it determines directly how many members of the specific [unit](#) will be online throughout the model when a single value is selected. It is also possible to provide a timeseries of values, which can be used for example to impose initial conditions by providing a value only for the first timestep included in the model.

## fix\_units\_on\_coefficient\_in\_in

Optional coefficient for the `units_on` variable impacting the `fix_ratio_in_in_unit_flow` constraint.

**Default value:** 0.0

Related [Relationship Classes](#): [unit\\_\\_node\\_\\_node](#)

The `fix_units_on_coefficient_in_in` parameter is an optional coefficient in the [unit input-input ratio constraint](#) controlled by the `fix_ratio_in_in_unit_flow` parameter. Essentially, it acts as a coefficient for the [units\\_on](#) variable in the constraint, allowing for fixing the conversion ratio depending on the amount of online capacity.

Note that there are different parameters depending on the directions of the [unit\\_flow](#) variables being constrained: [fix\\_units\\_on\\_coefficient\\_in\\_out](#), [fix\\_units\\_on\\_coefficient\\_out\\_in](#), and [fix\\_units\\_on\\_coefficient\\_out\\_out](#), all of which apply to their respective [constraints](#). Similarly, there are different parameters for setting minimum or maximum conversion rates, e.g. [min\\_units\\_on\\_coefficient\\_in\\_in](#) and [max\\_units\\_on\\_coefficient\\_in\\_in](#).

## fix\_units\_on\_coefficient\_in\_out

Optional coefficient for the `units_on` variable impacting the `fix_ratio_in_out_unit_flow` constraint.

Default value: 0.0

Related [Relationship Classes](#): [unit\\_\\_node\\_\\_node](#)

The [fix\\_units\\_on\\_coefficient\\_in\\_out](#) parameter is an optional coefficient in the [unit input-output ratio constraint](#) controlled by the [fix\\_ratio\\_in\\_out\\_unit\\_flow](#) parameter. Essentially, it acts as a coefficient for the [units\\_on](#) variable in the constraint, allowing for fixing the conversion ratio depending on the amount of online capacity.

Note that there are different parameters depending on the directions of the [unit\\_flow](#) variables being constrained: [fix\\_units\\_on\\_coefficient\\_in\\_in](#), [fix\\_units\\_on\\_coefficient\\_out\\_in](#), and [fix\\_units\\_on\\_coefficient\\_out\\_out](#), all of which apply to their respective [constraints](#). Similarly, there are different parameters for setting minimum or maximum conversion rates, e.g. [min\\_units\\_on\\_coefficient\\_in\\_out](#) and [max\\_units\\_on\\_coefficient\\_in\\_out](#).

## fix\_units\_on\_coefficient\_out\_in

Optional coefficient for the `units_on` variable impacting the `fix_ratio_out_in_unit_flow` constraint.

Default value: 0.0

Related [Relationship Classes](#): [unit\\_\\_node\\_\\_node](#)

The `fix_units_on_coefficient_out_in` parameter is an optional coefficient in the [unit output-input ratio constraint](#) controlled by the `fix_ratio_out_in_unit_flow` parameter. Essentially, it acts as a coefficient for the `units_on` variable in the constraint, allowing for fixing the conversion ratio depending on the amount of online capacity.

Note that there are different parameters depending on the directions of the `unit_flow` variables being constrained: `fix_units_on_coefficient_in_in`, `fix_units_on_coefficient_in_out`, and `fix_units_on_coefficient_out_out`, all of which apply to their respective [constraints](#). Similarly, there are different parameters for setting minimum or maximum conversion rates, e.g. `min_units_on_coefficient_out_in` and `max_units_on_coefficient_out_in`.

## fix\_units\_on\_coefficient\_out\_out

Optional coefficient for the `units_on` variable impacting the `fix_ratio_out_out_unit_flow` constraint.

Default value: 0.0

Related [Relationship Classes](#): `unit__node__node`

The `fix_units_on_coefficient_out_out` parameter is an optional coefficient in the [unit output-output ratio constraint](#) controlled by the `fix_ratio_out_out_unit_flow` parameter. Essentially, it acts as a coefficient for the `units_on` variable in the constraint, allowing for fixing the conversion ratio depending on the amount of online capacity.

Note that there are different parameters depending on the directions of the `unit_flow` variables being constrained: `fix_units_on_coefficient_in_in`, `fix_units_on_coefficient_in_out`, and `fix_units_on_coefficient_out_in`, all of which apply to their respective [constraints](#). Similarly, there are different parameters for setting minimum or maximum conversion rates, e.g. `min_units_on_coefficient_out_out` and `max_units_on_coefficient_out_out`.

## fixed\_pressure\_constant\_0

Fixed pressure points for pipelines for the outer approximation of the Weymouth approximation. The direction of flow is the first node in the relationship to the second node in the relationship.

Related [Relationship Classes](#): [connection\\_\\_node\\_\\_node](#)

For the MILP representation of pressure driven gas transfer, we use an outer approximation approach as described by [Schwele et al.](#). The Weymouth equation is approximated around fixed pressure points, as described by the constraint on [fixed node pressure points](#), constraining the average flow in each direction dependent on the adjacent node pressures. The second fixed pressure constant, which will be multiplied with the pressure of the destination node, is represented by an Array value of the [fixed\\_pressure\\_constant\\_0](#). The first pressure constant corresponds to the related parameter [fixed\\_pressure\\_constant\\_1](#). Note that the [fixed\\_pressure\\_constant\\_0](#) parameter should be defined on a [connection\\_\\_node\\_\\_node](#) relationship, for which the first node corresponds to the origin node, while the second node corresponds to the destination node. For a typical gas pipeline, the will be a [fixed\\_pressure\\_constant\\_1](#) for both directions of flow.

## fixed\_pressure\_constant\_1

Fixed pressure points for pipelines for the outer approximation of the Weymouth approximation. The direction of flow is the first node in the relationship to the second node in the relationship.

Related [Relationship Classes](#): [connection\\_\\_node\\_\\_node](#)

For the MILP representation of pressure driven gas transfer, we use an outer approximation approach as described by [Schwele et al.](#). The Weymouth equation is approximated around fixed pressure points, as described by the constraint on [fixed node pressure points](#), constraining the average flow in each direction dependent on the adjacent node pressures. The first fixed pressure constant, which will be multiplied with the pressure of the origin node, is represented by an Array value of the [fixed\\_pressure\\_constant\\_1](#). The second pressure constant corresponds to the related parameter [fixed\\_pressure\\_constant\\_0](#). Note that the [fixed\\_pressure\\_constant\\_1](#) parameter should be defined on a [connection\\_\\_node\\_\\_node](#) relationship, for which the first node corresponds to the origin node, while the second node corresponds to the destination node. For a typical gas pipeline, the will be a [fixed\\_pressure\\_constant\\_1](#) for both directions of flow.

## fom\_cost

Fixed operation and maintenance costs of a [unit](#). Essentially, a cost coefficient on the [number\\_of\\_units](#) and [unit\\_capacity](#) parameters. E.g. EUR/MWh

Related [Object Classes](#): [unit](#)

By defining the `fom_cost` parameter for a specific `unit`, a cost term will be added to the objective function to account for the fixed operation and maintenance costs associated with that unit during the current optimization window.

## frac\_state\_loss

Self-discharge coefficient for `node_state` variables. Effectively, represents the *loss power per unit of state*.

Default value: 0.0

Related [Object Classes](#): [node](#)

The `frac_state_loss` parameter allows setting self-discharge losses for `nodes` with the `node_state` variables enabled using the `has_state` variable. Effectively, the `frac_state_loss` parameter acts as a coefficient on the `node_state` variable in the [node injection constraint](#), imposing losses for the `node`. In simple cases, storage losses are typically fractional, e.g. a `frac_state_loss` parameter value of 0.01 would represent 1% of `node_state` lost per unit of time. However, a more general definition of what the `frac_state_loss` parameter represents in *SpineOpt* would be *loss power per unit of node\_state*.

## fractional\_demand

The fraction of a `node` group's `demand` applied for the `node` in question.

Default value: 0.0

Related [Object Classes](#): [node](#)

Whenever a `node` is a member of a group, the `fractional_demand` parameter represents its share of the group's `demand`.

## fuel\_cost

Variable fuel costs than can be attributed to a `unit_flow`. E.g. EUR/MWh

Related [Relationship Classes](#): `unit_from_node` and `unit_to_node`

By defining the `fuel_cost` parameter for a specific `unit`, `node`, and `direction`, a cost term will be added to the objective function to account for costs associated with the unit's fuel usage over the course of its operational dispatch during the current optimization window.

## graph\_view\_position

An optional setting for tweaking the position of the different elements when drawing them via Spine Toolbox Graph View.

Related [Relationship Classes](#): `connection_from_node`, `connection_to_node`, `unit_from_node_unit_constraint`, `unit_from_node`, `unit_to_node_unit_constraint` and `unit_to_node`

The `graph_view_position` parameter can be used to fix the positions of various objects and relationships when plotted using the *Spine Toolbox Graph View*. If not defined, *Spine Toolbox* simply plots the element in question wherever it sees fit in the graph.

## has\_binary\_gas\_flow

This parameter needs to be set to `true` in order to represent bidirectional pressure drive gas transfer.

Default value: `false`

Uses [Parameter Value Lists](#): `boolean_value_list`



Related [Object Classes](#): [connection](#)

This parameter is necessary for the use of pressure driven gas transfer, for which the direction of flow is not known a priori. The parameter [has\\_binary\\_gas\\_flow](#) is a boolean method parameter, which - when set to [true](#) - triggers the generation of the binary variables [binary\\_gas\\_connection\\_flow](#), which (together with the [big\\_m](#) parameter) forces the average flow through a pipeline to be unidirectional.

## has\_pressure

A boolean flag for whether a [node](#) has a [node\\_pressure](#) variable.

**Default value:** false

Uses [Parameter Value Lists](#): [boolean\\_value\\_list](#)

Related [Object Classes](#): [node](#)

If a node is to represent a node in a pressure driven gas network, the boolean parameter [has\\_pressure](#) should be set true, in order to trigger the generation of the [node\\_pressure](#) variable. The pressure at a certain node can also be constrained through the parameters [max\\_node\\_pressure](#) and [min\\_node\\_pressure](#). More details on the use of pressure driven gas transfer are described [here](#)

## has\_state

A boolean flag for whether a [node](#) has a [node\\_state](#) variable.

**Default value:** false

Uses [Parameter Value Lists](#): [boolean\\_value\\_list](#)

Related [Object Classes](#): [node](#)

The `has_state` parameter is simply a `Bool` flag for whether a `node` has a `node_state` variable. By default, it is set to `false`, so the `nodes` enforce instantaneous `commodity` balance according to the `nodal balance` and `node injection` constraints. If set to `true`, the `node` will have a `node_state` variable generated for it, allowing for `commodity` storage at the `node`. Note that you'll also have to specify a value for the `state_coeff` parameter, as otherwise the `node_state` variable has zero `commodity` capacity.

## has\_voltage\_angle

A boolean flag for whether a `node` has a `node_voltage_angle` variable.

**Default value:** `false`

**Uses Parameter Value Lists:** `boolean_value_list`

**Related Object Classes:** `node`

For the use of node-based lossless DC powerflow, each node will be associated with a `node_voltage_angle` variable. To enable the generation of the variable in the optimization model, the boolean parameter `has_voltage_angle` should be set `true`. The voltage angle at a certain node can also be constrained through the parameters `max_voltage_angle` and `min_voltage_angle`. More details on the use of lossless nodal DC power flows are described [here](#)

## is\_active

If `false`, the object is excluded from the model if the tool filter object activity control is specified

**Default value:** `true`

**Uses Parameter Value Lists:** `boolean_value_list`

Related **Object Classes**: [commodity](#), [connection](#), [model](#), [node](#), [output](#), [report](#), [stochastic\\_scenario](#), [stochastic\\_structure](#), [temporal\\_block](#), [unit\\_constraint](#) and [unit](#)

`is_active` is a universal, utility parameter that is defined for every object class. When used in conjunction with the `activity_control` feature, the `is_active` parameter allows one to control whether or not a specific object is active within a model or not.

## is\_reserve\_node

A boolean flag for whether a `node` is acting as a `reserve_node`

**Default value:** false

Uses **Parameter Value Lists**: [boolean\\_value\\_list](#)

Related **Object Classes**: [node](#)

By setting the parameter `is_reserve_node` to `true`, a node is treated as a reserve [node](#) in the model. Units that are linked through a `unit_to_node` relationship will be able to provide balancing services to the reserve node, but within their technical feasibility. The mathematical formulation holds a chapter on [Ramping and reserve constraints](#) and the general concept of setting up a model with reserves is described in [Ramping and Reserves](#).

## max\_cum\_in\_unit\_flow\_bound

Set a maximum cumulative upper bound for a `unit_flow`

Related **Relationship Classes**: [unit\\_commodity](#)

To impose a limit on the cumulative in flows to a unit for the entire modelling horizon, e.g. to enforce limits on emissions, the `max_cum_in_unit_flow_bound` parameter can be used. Defining this parameter triggers the generation of the `constraint_max_cum_in_unit_flow_bound`.

Assuming for instance that the total intake of a unit `u_A` should not exceed `10MWh` for the entire modelling horizon, then the `max_cum_in_unit_flow_bound` would need to take the value `10`. (Assuming here that the `unit_flow` variable is in `MW`, and the model `duration_unit` is `hours`)

## max\_gap

Specifies the maximum optimality gap for the model. Currently only used for the master problem within a decomposed structure

Default value: 0.05

Related [Object Classes](#): `model`

This determines the optimality convergence criterion and is the benders gap tolerance for the master problem in a decomposed investments model. The benders gap is the relative difference between the current objective function upper bound (*zupper*) and lower bound (*zlower*) and is defined as  $2 \cdot (zupper - zlower) / (zupper + zlower)$ . When this value is lower than `max_gap` the benders algorithm will terminate having achieved satisfactory optimality.

## max\_iterations

Specifies the maximum number of iterations for the model. Currently only used for the master problem within a decomposed structure

Default value: 10.0

Related [Object Classes](#): `model`

When the `model` in question is of type `:spineopt_master`, this determines the maximum number of Benders iterations.

## max\_node\_pressure

Maximum allowed gas pressure at `node`.

Related **Object Classes**: `node`

If a node has a `node_pressure` variable (see also the parameter `has_pressure` and this [chapter](#)), an upper bound on the pressure can be introduced through the `max_node_pressure` parameter, which triggers the generation of the `maximum node pressure` constraint.

## `max_ratio_in_in_unit_flow`

Maximum ratio between two `unit_flows` coming into the `unit` from the two `nodes`.

Related **Relationship Classes**: `unit__node__node`

The definition of the `max_ratio_in_in_unit_flow` parameter triggers the generation of the `constraint_max_ratio_in_in_unit_flow` and enforces an upper bound on the ratio between incoming flows of a unit. The parameter is defined on the relationship class `unit__node__node`, where both nodes (or group of nodes) in this relationship represent `from_nodes`, i.e. the incoming flows to the unit. The ratio parameter is interpreted such that it constrains the ratio of `in1` over `in2`, where `in1` is the `unit_flow` variable from the first `node` in the `unit__node__node` relationship in a left-to-right reading order. This parameter can be useful, for instance if a unit requires a specific commodity mix as a fuel supply.

To enforce e.g. for a unit `u` a maximum share of `0.8` of its incoming flow from the node `supply_fuel_1` compared to its incoming flow from the node group `supply_fuel_2` (consisting of the two nodes `supply_fuel_2_component_a` and `supply_fuel_2_component_b`) the `max_ratio_in_in_unit_flow` parameter would be set to `0.8` for the relationship `u__supply_fuel_1__supply_fuel_2`.

## `max_ratio_in_out_unit_flow`

Maximum ratio between an incoming `unit_flow` from the first `node` and an outgoing `unit_flow` to the second `node`.

Related **Relationship Classes**: `unit__node__node`

The definition of the `max_ratio_in_out_unit_flow` parameter triggers the generation of the `constraint_max_ratio_in_out_unit_flow` and sets an upper bound on the ratio between incoming and outgoing flows of a unit. The parameter is defined on the relationship class `unit_node_node`, where the first node (or group of nodes) in this relationship represents the `from_node`, i.e. the incoming flows to the unit, and the second node (or group of nodes), represents the `to_node` i.e. the outgoing flow from the unit. The ratio parameter is interpreted such that it constrains the ratio of `in` over `out`, where `in` is the `unit_flow` variable from the first `node` in the `unit_node_node` relationship in a left-to-right reading order.

To enforce e.g. a maximum ratio of `1.4` for a unit `u` between its incoming gas flow from the node `ng` and its outgoing flow to the node group `e1_heat` (consisting of the two nodes `e1` and `heat`), the `max_ratio_in_out_unit_flow` parameter would be set to `1.4` for the relationship `u__ng__e1_heat`.

## max\_ratio\_out\_in\_connection\_flow

Maximum ratio between the `connection_flow` from the first `node` and the `connection_flow` to the second `node`.

Related [Relationship Classes](#): `connection_node_node`

The definition of the `max_ratio_out_in_connection_flow` parameter triggers the generation of the `constraint_max_ratio_out_in_connection_flow` and sets an upper bound on the ratio between outgoing and incoming flows of a connection. The parameter is defined on the relationship class `connection_node_node`, where the first node (or group of nodes) in this relationship represents the `to_node`, i.e. the outgoing flow from the connection, and the second node (or group of nodes), represents the `from_node`, i.e. the incoming flows to the connection. The ratio parameter is interpreted such that it constrains the ratio of `out` over `in`, where `out` is the `connection_flow` variable from the first `node` in the `connection_node_node` relationship in a left-to-right reading order.

To enforce e.g. a maximum ratio of `0.8` for a connection `conn` between its outgoing electricity flow to `node commodity1` and its incoming flows from the node `node commodity2`, the `max_ratio_out_in_connection_flow` parameter would be set to `0.8` for the relationship `conn__commodity1__commodity2`.

Note that the ratio can also be defined for `connection_node_node` relationships where one or both of the nodes correspond to node groups in order to impose a ratio on aggregated connection flows.

## max\_ratio\_out\_in\_unit\_flow

Maximum ratio between an outgoing `unit_flow` to the first `node` and an incoming `unit_flow` from the second `node`.

Related [Relationship Classes](#): `unit_node_node`

The definition of the `max_ratio_out_in_unit_flow` parameter triggers the generation of the `constraint_max_ratio_out_in_unit_flow` and enforces an upper bound on the ratio between outgoing and incoming flows of a unit. The parameter is defined on the relationship class `unit_node_node`, where the first node (or group of nodes) in this relationship represents the `to_node`, i.e. the outgoing flow from the unit, and the second node (or group of nodes), represents the `from_node`, i.e. the incoming flows to the unit. The ratio parameter is interpreted such that it constrains the ratio of `out` over `in`, where `out` is the `unit_flow` variable from the first `node` in the `unit_node_node` relationship in a left-to-right reading order.

To enforce e.g. a maximum ratio of `0.8` for a unit `u` between its outgoing flows to the node group `e1_heat` (consisting of the two nodes `e1` and `heat`) and its incoming gas flow from `ng` the `max_ratio_out_in_unit_flow` parameter would be set to `0.8` for the relationship `u__e1_heat__ng`.

## max\_ratio\_out\_out\_unit\_flow

Maximum ratio between two `unit_flows` going from the `unit` into the two `nodes`.

Related [Relationship Classes](#): `unit_node_node`

The definition of the `max_ratio_out_out_unit_flow` parameter triggers the generation of the `constraint_max_ratio_out_out_unit_flow` and sets an upper bound on the ratio between outgoing flows of a unit. The parameter is defined on the relationship class `unit_node_node`, where the nodes (or group of nodes) in this relationship represent the `to_node`'s, i.e. outgoing flow from the unit. The ratio parameter is interpreted such that it constrains the ratio of `out1` over `out2`, where `out1` is the `unit_flow` variable from the first `node` in the `unit_node_node` relationship in a left-to-right reading order.

To enforce a maximum ratio between two products of a unit `u`, e.g. setting the maximum share of produced electricity flowing to node `e1` to `0.4` of the production of heat flowing to node `heat`, the `fix_ratio_out_out_unit_flow` parameter would be set to `0.4` for the relationship `u__e1__heat`.

## max\_res\_shutdown\_ramp

Maximum non-spinning reserve ramp-down for online units providing reserves during shut-downs

**Related Relationship Classes:** [unit\\_from\\_node](#) and [unit\\_to\\_node](#)

A unit can provide spinning and nonspinning reserves to a [reserve node](#). These reserves can be either [upward\\_reserve](#) or [downward\\_reserve](#). Nonspinning downward reserves are provided to a [downward\\_reserve](#) node by contracted units holding available to shutdown. To include the provision of nonspinning downward reserves, the parameter [max\\_res\\_shutdown\\_ramp](#) needs to be defined on the corresponding [unit\\_to\\_node](#) relationship. This will trigger the generation of the variables [nonspin\\_units\\_shut\\_down](#) and [nonspin\\_ramp\\_down\\_unit\\_flow](#) and the constraint [on maximum downward nonspinning reserve provision](#). Note that [max\\_res\\_shutdown\\_ramp](#) is given as a fraction of the [unit\\_capacity](#).

A detailed description of the usage of ramps and reserves is given in the chapter [Ramping and Reserves](#). The chapter [Ramping and reserve constraints](#) in the Mathematical Formulation presents the equations related to ramps and reserves.

## max\_res\_startup\_ramp

Maximum non-spinning reserve ramp-up for offline units scheduled for reserve provision

**Related Relationship Classes:** [unit\\_from\\_node](#) and [unit\\_to\\_node](#)

A unit can provide spinning and nonspinning reserves to a [reserve node](#). These reserves can be either [upward\\_reserve](#) or [downward\\_reserve](#). Nonspinning upward reserves are provided to a [upward\\_reserve](#) node by contracted offline units holding available to startup. To include the provision of nonspinning upward reserves, the parameter [max\\_res\\_startup\\_ramp](#) needs to be defined on the corresponding [unit\\_to\\_node](#) relationship. This will trigger the generation of the variables [nonspin\\_units\\_started\\_up](#) and [nonspin\\_ramp\\_up\\_unit\\_flow](#) and the constraint [on maximum upward nonspinning reserve provision](#). Note that [max\\_res\\_startup\\_ramp](#) is given as a fraction of the [unit\\_capacity](#).

A detailed description of the usage of ramps and reserves is given in the chapter [Ramping and Reserves](#). The chapter [Ramping and reserve constraints](#) in the Mathematical Formulation presents the equations related to ramps and reserves.

## max\_shutdown\_ramp



Maximum ramp-down during shutdowns

Related [Relationship Classes](#): [unit\\_from\\_node](#) and [unit\\_to\\_node](#)

The definition of the `max_shutdown_ramp` parameter will trigger the creation of the [constraint on maximum shutdown ramp](#). It sets an upper bound on the [unit\\_flow](#) variable for the timestep right before a shutdown.

It can be defined for [unit\\_to\\_node](#) or [unit\\_from\\_node](#) relationships, as well as their counterparts for node groups. It will then impose restrictions on the `unit_flow` variables that indicate flows between the two members of the relationship for which the parameter is defined. The parameter is given as a fraction of the [unit\\_capacity](#) parameter. When the parameter is not included, the aforementioned constraint will not be created, which is equivalent to choosing a value of 1.

## max\_startup\_ramp

Maximum ramp-up during startups

Related [Relationship Classes](#): [unit\\_from\\_node](#) and [unit\\_to\\_node](#)

The definition of the `max_startup_ramp` parameter will trigger the creation of the [Constraint on upward start up ramp\\_up](#). It sets an upper bound on the [unit\\_flow](#) variable for the timestep right after a startup.

It can be defined for [unit\\_to\\_node](#) or [unit\\_from\\_node](#) relationships, as well as their counterparts for node groups. It will then impose restrictions on the `unit_flow` variables that indicate flows between the two members of the relationship for which the parameter is defined. The parameter is given as a fraction of the [unit\\_capacity](#) parameter. When the parameter is not included, the aforementioned constraint will not be created, which is equivalent to choosing a value of 1.

## max\_units\_on\_coefficient\_in\_in

Optional coefficient for the `units_on` variable impacting the `max_ratio_in_in_unit_flow` constraint.

Default value: 0.0

Related [Relationship Classes](#): [unit\\_\\_node\\_\\_node](#)

The [max\\_units\\_on\\_coefficient\\_in\\_in](#) parameter is an optional coefficient in the [unit input-input ratio constraint](#) controlled by the [max\\_ratio\\_in\\_in\\_unit\\_flow](#) parameter. Essentially, it acts as a coefficient for the [units\\_on](#) variable in the constraint, allowing for making the maximum conversion ratio dependent on the amount of online capacity.

Note that there are different parameters depending on the directions of the [unit\\_flow](#) variables being constrained: [max\\_units\\_on\\_coefficient\\_in\\_out](#), [max\\_units\\_on\\_coefficient\\_out\\_in](#), and [max\\_units\\_on\\_coefficient\\_out\\_out](#), all of which apply to their respective [constraints](#). Similarly, there are different parameters for setting minimum or fixed conversion rates, e.g. [min\\_units\\_on\\_coefficient\\_in\\_in](#) and [fix\\_units\\_on\\_coefficient\\_in\\_in](#).

## max\_units\_on\_coefficient\_in\_out

Optional coefficient for the [units\\_on](#) variable impacting the [max\\_ratio\\_in\\_out\\_unit\\_flow](#) constraint.

Default value: 0.0

Related [Relationship Classes](#): [unit\\_\\_node\\_\\_node](#)

The [max\\_units\\_on\\_coefficient\\_in\\_out](#) parameter is an optional coefficient in the [unit input-output ratio constraint](#) controlled by the [max\\_ratio\\_in\\_out\\_unit\\_flow](#) parameter. Essentially, it acts as a coefficient for the [units\\_on](#) variable in the constraint, allowing for making the maximum conversion ratio dependent on the amount of online capacity.

Note that there are different parameters depending on the directions of the [unit\\_flow](#) variables being constrained: [max\\_units\\_on\\_coefficient\\_in\\_in](#), [max\\_units\\_on\\_coefficient\\_out\\_in](#), and [max\\_units\\_on\\_coefficient\\_out\\_out](#), all of which apply to their respective [constraints](#). Similarly, there are different parameters for setting minimum or fixed conversion rates, e.g. [min\\_units\\_on\\_coefficient\\_in\\_out](#) and [fix\\_units\\_on\\_coefficient\\_in\\_out](#).

## max\_units\_on\_coefficient\_out\_in

Optional coefficient for the `units_on` variable impacting the `max_ratio_out_in_unit_flow` constraint.

**Default value:** 0.0

**Related Relationship Classes:** `unit_node_node`

The `max_units_on_coefficient_out_in` parameter is an optional coefficient in the `unit output-input ratio constraint` controlled by the `max_ratio_out_in_unit_flow` parameter. Essentially, it acts as a coefficient for the `units_on` in the constraint, allowing for making the maximum conversion ratio dependent on the amount of online capacity.

Note that there are different parameters depending on the directions of the `unit_flow` being constrained: `max_units_on_coefficient_in_in`, `max_units_on_coefficient_in_out`, and `max_units_on_coefficient_out_out`, all of which apply to their respective `constraints`. Similarly, there are different parameters for setting minimum or fixed conversion rates, e.g. `min_units_on_coefficient_out_in` and `fix_units_on_coefficient_out_in`.

## max\_units\_on\_coefficient\_out\_out

Optional coefficient for the `units_on` variable impacting the `max_ratio_out_out_unit_flow` constraint.

**Default value:** 0.0

**Related Relationship Classes:** `unit_node_node`

The `max_units_on_coefficient_out_out` parameter is an optional coefficient in the `unit output-output ratio constraint` controlled by the `max_ratio_out_out_unit_flow` parameter. Essentially, it acts as a coefficient for the `units_on` variable in the constraint, allowing for making the maximum conversion ratio dependent on the amount of online capacity.

Note that there are different parameters depending on the directions of the `unit_flow` variables being constrained: `max_units_on_coefficient_in_in`, `max_units_on_coefficient_out_in`, and `max_units_on_coefficient_in_out`, all of which apply to their respective `constraints`. Similarly, there are

different parameters for setting minimum or fixed conversion rates, e.g. [min\\_units\\_on\\_coefficient\\_out\\_out](#) and [fix\\_units\\_on\\_coefficient\\_out\\_out](#).

## max\_voltage\_angle

Maximum allowed voltage angle at [node](#).

Related [Object Classes](#): [node](#)

If a node has a [node\\_voltage\\_angle](#) variable (see also the parameter [has\\_voltage\\_angle](#) and this [chapter](#)), an upper bound on the voltage angle can be introduced through the [max\\_voltage\\_angle](#) parameter, which triggers the generation of the [maximum node voltage angle](#) constraint.

## min\_down\_time

Minimum downtime of a [unit](#) after it shuts down.

Related [Object Classes](#): [unit](#)

The definition of the [min\\_down\\_time](#) parameter will trigger the creation of the [Constraint on minimum down time](#). It sets a lower bound on the period that a unit has to stay offline after a shutdown.

It can be defined for a [unit](#) and will then impose restrictions on the [units\\_on](#) variables that represent the on- or offline status of the unit. The parameter is given as a duration value. When the parameter is not included, the aforementioned constraint will not be created, which is equivalent to choosing a value of 0.

For a more complete description of unit commitment restrictions, see [Unit commitment](#).

## min\_node\_pressure

Minimum allowed gas pressure at [node](#).

Related [Object Classes](#): [node](#)

If a node has a [node\\_pressure](#) variable (see also the parameter [has\\_pressure](#) and this [chapter](#)), a lower bound on the pressure can be introduced through the [min\\_node\\_pressure](#) parameter, which triggers the generation of the [minimum node pressure](#) constraint.

## min\_ratio\_in\_in\_unit\_flow

Minimum ratio between two [unit\\_flows](#) coming into the [unit](#) from the two [nodes](#).

Related [Relationship Classes](#): [unit\\_\\_node\\_\\_node](#)

The definition of the [min\\_ratio\\_in\\_in\\_unit\\_flow](#) parameter triggers the generation of the [constraint\\_min\\_ratio\\_in\\_in\\_unit\\_flow](#) and sets a lower bound for the ratio between incoming flows of a unit. The parameter is defined on the relationship class [unit\\_\\_node\\_\\_node](#), where both nodes (or group of nodes) in this relationship represent [from\\_nodes](#), i.e. the incoming flows to the unit. The ratio parameter is interpreted such that it constrains the ratio of [in1](#) over [in2](#), where [in1](#) is the [unit\\_flow](#) variable from the first [node](#) in the [unit\\_\\_node\\_\\_node](#) relationship in a left-to-right reading order. This parameter can be useful, for instance if a unit requires a specific commodity mix as a fuel supply.

To enforce e.g. for a unit [u](#) a minimum share of [0.2](#) of its incoming flow from the node [supply\\_fuel\\_1](#) compared to its incoming flow from the node group [supply\\_fuel\\_2](#) (consisting of the two nodes [supply\\_fuel\\_2\\_component\\_a](#) and [supply\\_fuel\\_2\\_component\\_b](#)) the [min\\_ratio\\_in\\_in\\_unit\\_flow](#) parameter would be set to [0.2](#) for the relationship [u\\_\\_supply\\_fuel\\_1\\_\\_supply\\_fuel\\_2](#).

## min\_ratio\_in\_out\_unit\_flow

Minimum ratio between an incoming [unit\\_flow](#) from the first [node](#) and an outgoing [unit\\_flow](#) to the second [node](#).

Related [Relationship Classes](#): [unit\\_\\_node\\_\\_node](#)

The definition of the [min\\_ratio\\_in\\_out\\_unit\\_flow](#) parameter triggers the generation of the [constraint\\_min\\_ratio\\_in\\_out\\_unit\\_flow](#) and enforces a lower bound on the ratio between incoming and outgoing flows of a unit. The parameter is defined on the relationship class [unit\\_\\_node\\_\\_node](#), where the first node (or group of nodes, see) in this relationship represents the [from\\_node](#), i.e. the incoming flow to the unit, and the second node (or group of nodes) represents the [to\\_node](#) i.e. the outgoing flow from the unit. The ratio parameter is interpreted such that it constrains the ratio of [in](#) over [out](#), where [in](#) is the [unit\\_flow](#) variable from the first [node](#) in the [unit\\_\\_node\\_\\_node](#) relationship in a left-to-right reading order.

To enforce e.g. a minimum ratio of 1.4 for a unit `u` between its incoming gas flow from the node `ng` and its outgoing flow to the node group `e1_heat` (consisting of the two nodes `e1` and `heat`), the `fix_ratio_in_out_unit_flow` parameter would be set to 1.4 for the relationship `u__ng__e1_heat`.

## min\_ratio\_out\_in\_connection\_flow

Minimum ratio between the `connection_flow` from the first `node` and the `connection_flow` to the second `node`.

Related **Relationship Classes**: `connection__node__node`

The definition of the `min_ratio_out_in_connection_flow` parameter triggers the generation of the `constraint_min_ratio_out_in_connection_flow` and sets a lower bound on the ratio between outgoing and incoming flows of a connection. The parameter is defined on the relationship class `connection__node__node`, where the first node (or group of nodes) in this relationship represents the `to_node`, i.e. the outgoing flow from the connection, and the second node (or group of nodes), represents the `from_node`, i.e. the incoming flows to the connection. The ratio parameter is interpreted such that it constrains the ratio of `out` over `in`, where `out` is the `connection_flow` variable from the first `node` in the `connection__node__node` relationship in a left-to-right reading order.

Note that the ratio can also be defined for `connection__node__node` relationships, where one or both of the nodes correspond to node groups in order to impose a ratio on aggregated connection flows.

To enforce e.g. a minimum ratio of 0.2 for a connection `conn` between its outgoing electricity flow to `node commodity1` and its incoming flows from the node `node commodity2`, the `min_ratio_out_in_connection_flow` parameter would be set to 0.8 for the relationship `conn__commodity1__commodity2`.

## min\_ratio\_out\_in\_unit\_flow

Minimum ratio between an outgoing `unit_flow` to the first `node` and an incoming `unit_flow` from the second `node`.

Related **Relationship Classes**: `unit__node__node`

The definition of the `[min_ratio_out_in_unit_flow]` parameter triggers the generation of the `constraint_min_ratio_out_in_unit_flow` and corresponds to a lower bound of the ratio between out and incoming flows of a unit. The parameter is defined on the relationship class `unit__node__node`, where

the first node (or group of nodes) in this relationship represents the `to_node`, i.e. the outgoing flow from the unit, and the second node (or group of nodes), represents the `from_node`, i.e. the incoming flows to the unit. The ratio parameter is interpreted such that it constrains the ratio of `out` over `in`, where `out` is the `unit_flow` variable from the first `node` in the `unit_node_node` relationship in a left-to-right reading order.

To enforce e.g. a minimum ratio of 0.8 for a unit `u` between its outgoing flows to the node group `e1_heat` (consisting of the two nodes `e1` and `heat`) and its incoming gas flow from `ng` the `min_ratio_out_in_unit_flow` parameter would be set to 0.8 for the relationship `u__e1_heat__ng`.

## min\_ratio\_out\_out\_unit\_flow

Minimum ratio between two `unit_flows` going from the `unit` into the two nodes.

Related [Relationship Classes](#): `unit_node_node`

The definition of the `min_ratio_out_out_unit_flow` parameter triggers the generation of the `constraint_min_ratio_out_out_unit_flow` and enforces a lower bound on the ratio between outgoing flows of a unit. The parameter is defined on the relationship class `unit_node_node`, where the nodes (or group of nodes) in this relationship represent the `to_node`'s', i.e. outgoing flow from the unit. The ratio parameter is interpreted such that it constrains the ratio of `out1` over `out2`, where `out1` is the `unit_flow` variable from the first `node` in the `unit_node_node` relationship in a left-to-right reading order.

To enforce a minimum ratio between two products of a unit `u`, e.g. setting the minimum share of produced electricity flowing to node `e1` to 0.4 of the production of heat flowing to node `heat`, the `fix_ratio_out_out_unit_flow` parameter would be set to 0.4 for the relationship `u__e1__heat`.

## min\_res\_shutdown\_ramp

Minimum non-spinning reserve ramp-down for online units providing reserves during shutdowns

Related [Relationship Classes](#): `unit_from_node` and `unit_to_node`

A unit can provide spinning and nonspinning reserves to a `reserve node`. These reserves can be either `upward_reserve` or `downward_reserve`. Nonspinning downward reserves are provided by contracted units holding available to shutdown to a `downward_reserve` node. If a unit is scheduled to provide

nonspinning reserve, a limit on the minimum amount of reserves provided can be imposed by defining the parameter [min\\_res\\_shutdown\\_ramp](#) on a [unit\\_to\\_node](#) relationship, which triggers the constraint [on minimum downward nonspinning reserve provision](#). The parameter [min\\_res\\_shutdown\\_ramp](#) is given as a fraction of the [unit\\_capacity](#) of the corresponding [unit\\_to\\_node](#) relationship.

Note that to include the provision of nonspinning downward reserves, the parameter [max\\_res\\_shutdown\\_ramp](#) needs to be defined on the corresponding [unit\\_to\\_node](#) relationship, which triggers the generation of the variables [nonspin\\_units\\_shut\\_down](#) and [nonspin\\_ramp\\_down\\_unit\\_flow](#).

A detailed description of the usage of ramps and reserves is given in the chapter [Ramping and Reserves](#). The chapter [Ramping and reserve constraints](#) in the Mathematical Formulation presents the equations related to ramps and reserves.

## min\_res\_startup\_ramp

Minimum non-spinning reserve ramp-up for offline units scheduled for reserve provision

Related [Relationship Classes](#): [unit\\_from\\_node](#) and [unit\\_to\\_node](#)

A unit can provide spinning and nonspinning reserves to a [reserve node](#). These reserves can be either [upward\\_reserve](#) or [downward\\_reserve](#). Nonspinning upward reserves are provided to an [upward\\_reserve](#) node by contracted offline units holding available to startup. If a unit is scheduled to provide nonspinning reserve, a limit on the minimum amount of reserves provided can be imposed by defining the parameter [min\\_res\\_startup\\_ramp](#) on a [unit\\_to\\_node](#) relationship, which triggers the constraint [on minimum upward nonspinning reserve provision](#). The parameter [min\\_res\\_startup\\_ramp](#) is given as a fraction of the [unit\\_capacity](#) of the corresponding [unit\\_to\\_node](#) relationship.

Note that to include the provision of nonspinning upward reserves, the parameter [max\\_res\\_startup\\_ramp](#) needs to be defined on the corresponding [unit\\_to\\_node](#) relationship, which triggers the generation of the variables [nonspin\\_units\\_started\\_up](#) and [nonspin\\_ramp\\_up\\_unit\\_flow](#).

A detailed description of the usage of ramps and reserves is given in the chapter [Ramping and Reserves](#). The chapter [Ramping and reserve constraints](#) in the Mathematical Formulation presents the equations related to ramps and reserves.

## min\_shutdown\_ramp

Minimum ramp-up during startups



Related [Relationship Classes](#): [unit\\_from\\_node](#) and [unit\\_to\\_node](#)

The definition of the `min_shutdown_ramp` parameter will trigger the creation of the [constraint on minimum shutdown ramp](#). It sets a lower bound on the [unit\\_flow](#) variable for the timestep right before a shutdown.

It can be defined for [unit\\_to\\_node](#) or [unit\\_from\\_node](#) relationships, as well as their counterparts for node groups. It will then impose restrictions on the `unit_flow` variables that indicate flows between the two members of the relationship for which the parameter is defined. The parameter is given as a fraction of the [unit\\_capacity](#) parameter. When the parameter is not included, the aforementioned constraint will not be created, which is equivalent to choosing a value of 0.

## min\_startup\_ramp

Minimum ramp-up during startups

Related [Relationship Classes](#): [unit\\_from\\_node](#) and [unit\\_to\\_node](#)

The definition of the `min_startup_ramp` parameter will trigger the creation of the [constraint on minimum startup ramp](#). It sets a lower bound on the [unit\\_flow](#) variable for the timestep right after a startup.

It can be defined for [unit\\_to\\_node](#) or [unit\\_from\\_node](#) relationships, as well as their counterparts for node groups. It will then impose restrictions on the `unit_flow` variables that indicate flows between the two members of the relationship for which the parameter is defined. The parameter is given as a fraction of the [unit\\_capacity](#) parameter. When the parameter is not included, the aforementioned constraint will not be created, which is equivalent to choosing a value of 0.

## min\_units\_on\_coefficient\_in\_in

Optional coefficient for the `units_on` variable impacting the `min_ratio_in_in_unit_flow` constraint.

Default value: 0.0

Related [Relationship Classes](#): [unit\\_node\\_node](#)

The `min_units_on_coefficient_in_in` parameter is an optional coefficient in the [unit input-input ratio constraint](#) controlled by the `min_ratio_in_in_unit_flow` parameter. Essentially, it acts as a coefficient for the `units_on` variable in the constraint, allowing for making the minimum conversion ratio dependent on the amount of online capacity.

Note that there are different parameters depending on the directions of the `unit_flow` variables being constrained: `min_units_on_coefficient_in_out`, `min_units_on_coefficient_out_in`, and `min_units_on_coefficient_out_out`, all of which apply to their respective [constraints](#). Similarly, there are different parameters for setting maximum or fixed conversion rates, e.g. `max_units_on_coefficient_in_in` and `fix_units_on_coefficient_in_in`.

## min\_units\_on\_coefficient\_in\_out

Optional coefficient for the `units_on` variable impacting the `min_ratio_in_out_unit_flow` constraint.

Default value: 0.0

Related [Relationship Classes](#): `unit__node__node`

The `min_units_on_coefficient_in_out` parameter is an optional coefficient in the [unit input-output ratio constraint](#) controlled by the `min_ratio_in_out_unit_flow` parameter. Essentially, it acts as a coefficient for the `units_on` variable in the constraint, allowing for making the minimum conversion ratio dependent on the amount of online capacity.

Note that there are different parameters depending on the directions of the `unit_flow` variables being constrained: `min_units_on_coefficient_in_in`, `min_units_on_coefficient_out_in`, and `min_units_on_coefficient_out_out`, all of which apply to their respective [constraints](#). Similarly, there are different parameters for setting maximum or fixed conversion rates, e.g. `max_units_on_coefficient_in_out` and `fix_units_on_coefficient_in_out`.

## min\_units\_on\_coefficient\_out\_in

Optional coefficient for the `units_on` variable impacting the `min_ratio_out_in_unit_flow` constraint.

Default value: 0.0

Related [Relationship Classes](#): [unit\\_\\_node\\_\\_node](#)

The [min\\_units\\_on\\_coefficient\\_out\\_in](#) parameter is an optional coefficient in the [unit output-input ratio constraint](#) controlled by the [min\\_ratio\\_out\\_in\\_unit\\_flow](#) parameter. Essentially, it acts as a coefficient for the [units\\_on](#) variable in the constraint, allowing for making the minimum conversion ratio dependent on the amount of online capacity.

Note that there are different parameters depending on the directions of the [unit\\_flow](#) variables being constrained: [min\\_units\\_on\\_coefficient\\_in\\_in](#), [min\\_units\\_on\\_coefficient\\_in\\_out](#), and [min\\_units\\_on\\_coefficient\\_out\\_out](#), all of which apply to their respective [constraints](#). Similarly, there are different parameters for setting maximum or fixed conversion rates, e.g. [max\\_units\\_on\\_coefficient\\_out\\_in](#) and [fix\\_units\\_on\\_coefficient\\_out\\_in](#).

## min\_units\_on\_coefficient\_out\_out

Optional coefficient for the [units\\_on](#) variable impacting the [min\\_ratio\\_out\\_out\\_unit\\_flow](#) constraint.

Default value: 0.0

Related [Relationship Classes](#): [unit\\_\\_node\\_\\_node](#)

The [min\\_units\\_on\\_coefficient\\_out\\_out](#) parameter is an optional coefficient in the [unit output-output ratio constraint](#) controlled by the [min\\_ratio\\_out\\_out\\_unit\\_flow](#) parameter. Essentially, it acts as a coefficient for the [units\\_on](#) variable in the constraint, allowing for making the minimum conversion ratio dependent on the amount of online capacity.

Note that there are different parameters depending on the directions of the [unit\\_flow](#) variables being constrained: [min\\_units\\_on\\_coefficient\\_in\\_in](#), [min\\_units\\_on\\_coefficient\\_in\\_out](#), and [min\\_units\\_on\\_coefficient\\_out\\_in](#), all of which apply to their respective [constraints](#). Similarly, there are different parameters for setting maximum or fixed conversion rates, e.g. [max\\_units\\_on\\_coefficient\\_out\\_out](#) and [fix\\_units\\_on\\_coefficient\\_out\\_out](#).

## min\_up\_time

Minimum uptime of a `unit` after it starts up.

Related [Object Classes](#): `unit`

The definition of the `min_up_time` parameter will trigger the creation of the [Constraint on minimum up time](#). It sets a lower bound on the period that a unit has to stay online after a startup.

It can be defined for a `unit` and will then impose restrictions on the `units_on` variables that represent the on- or offline status of the unit. The parameter is given as a duration value. When the parameter is not included, the aforementioned constraint will not be created, which is equivalent to choosing a value of 0.

For a more complete description of unit commitment restrictions, see [Unit commitment](#).

## min\_voltage\_angle

Minimum allowed voltage angle at `node`.

Related [Object Classes](#): `node`

If a node has a `node_voltage_angle` variable (see also the parameter `has_voltage_angle` and this [chapter](#)), a lower bound on the pressure can be introduced through the `min_voltage_angle` parameter, which triggers the generation of the [minimum node voltage angle](#) constraint.

## minimum\_operating\_point

Minimum level for the `unit_flow` relative to the `units_on` online capacity.

Related [Relationship Classes](#): `unit_from_node` and `unit_to_node`

The definition of the `minimum_operating_point` parameter will trigger the creation of the [Constraint on minimum operating point](#). It sets a lower bound on the value of the `unit_flow` variable for a unit that is online.

It can be defined for `unit_to_node` or `unit_from_node` relationships, as well as their counterparts for node groups. It will then impose restrictions on the `unit_flow` variables that indicate flows between

the two members of the relationship for which the parameter is defined. The parameter is given as a fraction of the [unit\\_capacity](#) parameter. When the parameter is not included, the aforementioned constraint will not be created, which is equivalent to choosing a value of 0.

## minimum\_reserve\_activation\_time

Duration a certain reserve product needs to be online/available

Related [Object Classes](#): [node](#)

The parameter [minimum\\_reserve\\_activation\\_time](#) is the duration a reserve product needs to be online, before it can be replaced by another (slower) reserve product.

In SpineOpt, the parameter is used to model reserve provision through storages. If a storage provides reserves to a reserve [node](#) (see also [is\\_reserve\\_node](#)) one needs to ensure that the node state is sufficiently high to provide these scheduled reserves as least for the duration of the [minimum\\_reserve\\_activation\\_time](#). The [constraint on the minimum node state with reserve provision](#) is triggered by the existence of the [minimum\\_reserve\\_activation\\_time](#). See also [Ramping and Reserves](#)

## model\_end

Defines the last timestamp to be modelled. Rolling optimization terminates after passing this point.

**Default value:** Dict{String, Any}{"data" => "2000-01-02T00:00:00", "type" => "date\_time"}

Related [Object Classes](#): [model](#)

Together with the [model\\_start](#) parameter, it is used to define the temporal horizon of the model. In case of a single solve optimization, the parameter marks the end of the last timestep that is possibly part of the optimization. Note that it poses an upper bound, and that the optimization does not necessarily include this timestamp when the [block\\_end](#) parameters are more stringent.

In case of a rolling horizon optimization, it will tell to the model to stop rolling forward once an optimization has been performed for which the result of the indicated timestamp has been kept in the final results. For example, assume that a [model\\_end](#) value of `2030-01-01T05:00:00` has been chosen, a [block\\_end](#) of `3h`, and a [roll\\_forward](#) of `2h`. The [roll\\_forward](#) parameter indicates here that the results of

the first two hours of each optimization window are kept as final, therefore the last optimization window will span the timeframe [ 2030-01-01T04:00:00 - 2030-01-01T06:00:00 ].

A DateTime value should be chosen for this parameter.

## model\_start

Defines the first timestamp to be modelled. Relative `temporal_blocks` refer to this value for their start and end.

**Default value:** Dict{String, Any}{"data" => "2000-01-01T00:00:00", "type" => "date\_time"}

Related **Object Classes:** [model](#)

Together with the [model\\_end](#) parameter, it is used to define the temporal horizon of the model. For a single solve optimization, it marks the timestamp from which the relative offset in a [temporal\\_block](#) is defined by the [block\\_start](#) parameter. In the rolling optimization framework, it does this for the first optimization window.

A DateTime value should be chosen for this parameter.

## model\_type

Used to identify model objects as relating to the master problem or operational sub problems (default)

**Default value:** spineopt\_operations

Uses **Parameter Value Lists:** [model\\_type\\_list](#)

Related **Object Classes:** [model](#)

This parameter is used, generally, to control [model](#) dependent functionality and specify [model](#)-level parameters for different [models](#). Currently, the main use is to identify the [model](#) objects that represent

the master and operational sub problems within a decomposed investment problem structure. To trigger the decomposed structure, a [model](#) object with `model_type=:spineopt_master` must exist and another with `model_type=:spineopt_operations` must also be present. To deactivate the decomposition functionality, the `model_type` of the master problem can be set to `:spineopt_other`.

See also [Decomposition](#).

## nodal\_balance\_sense

A selector for `nodal_balance` constraint sense.

Default value: `==`

Uses [Parameter Value Lists](#): `constraint_sense_list`

Related [Object Classes](#): `node`

`nodal_balance_sense` determines whether or not a [node](#) is able to naturally consume or produce energy. The default value, `==`, means that the [node](#) is unable to do any of that, and thus it needs to be perfectly balanced. The value `>=` means that the [node](#) is a *sink*, that is, it can *consume* any amounts of energy. The value `<=` means that the [node](#) is a *source*, that is, it can *produce* any amounts of energy.

## node\_opf\_type

A selector for the reference `node` (slack bus) when PTDF-based DC load-flow is enabled.

Default value: `node_opf_type_normal`

Uses [Parameter Value Lists](#): `node_opf_type_list`

Related [Object Classes](#): `node`

Used to identify the reference node (or slack bus) when ptdf based dc load flow is enabled ([commodity\\_physics](#) set to [commodity\\_physics\\_ptdf](#) or [commodity\\_physics\\_lodf](#). To identify the reference node, set `node_opf_type = :node_opf_type_reference`

See also [powerflow](#).

## node\_slack\_penalty

A penalty cost for `node_slack_pos` and `node_slack_neg` variables. The slack variables won't be included in the model unless there's a cost defined for them.

Related [Object Classes](#): [node](#)

`node_slack_penalty` triggers the creation of node slack variables, `node_slack_pos` and `node_slack_neg`. This allows the model to violate the [node\\_balance](#) constraint with these violations penalised in the objective function with a coefficient equal to `node_slack_penalty`. If `node_slack_penalty = 0` the slack variables are created and violations are unpenalised. If set to none or undefined, the variables are not created and violation of the [node\\_balance](#) constraint is not possible.

## node\_state\_cap

The maximum permitted value for a `node_state` variable.

Related [Object Classes](#): [node](#)

The [node\\_state\\_cap](#) parameter represents the maximum allowed value for the `node_state` variable. Note that in order for a [node](#) to have a `node_state` variable in the first place, the [has\\_state](#) parameter must be set to `true`. However, if the [node](#) has storage investments enabled using the [candidate\\_storages](#) parameter, the [node\\_state\\_cap](#) parameter acts as a coefficient for the `storages_invested_available` variable. Essentially, with investments, the [node\\_state\\_cap](#) parameter represents *storage capacity per storage investment*.

## node\_state\_coefficient

Coefficient of the specified node's state variable in the specified unit constraint.



Default value: 0.0

Related [Relationship Classes](#): [node\\_\\_unit\\_constraint](#)

The [node\\_state\\_coefficient](#) is an optional parameter that can be used to include the [node\\_state](#) variable of a [node](#) in a [unit\\_constraint](#) via the [node\\_\\_unit\\_constraint](#) relationship. Essentially, [node\\_state\\_coefficient](#) appears as a coefficient for the [node\\_state](#) variable of the [node](#) in the [unit\\_constraint](#).

## node\_state\_min

The minimum permitted value for a [node\\_state](#) variable.

Default value: 0.0

Related [Object Classes](#): [node](#)

The [node\\_state\\_min](#) parameter sets the lower bound for the [node\\_state](#) variable, if one has been enabled by the [has\\_state](#) parameter. For reserve [nodes](#) with [minimum\\_reserve\\_activation\\_time](#), the [node\\_state\\_min](#) is considered also via a special constraint.

## number\_of\_units

Denotes the number of 'sub units' aggregated to form the modelled [unit](#).

Default value: 1.0

Related [Object Classes](#): [unit](#)

Defines how many members a certain [unit](#) object represents. Typically this parameter takes a binary (UC) or integer (clustered UC) value. Together with the [unit\\_availability\\_factor](#), this will determine the maximum number of members that can be online at any given time. (Thus restricting the [units\\_on](#)

variable). It is possible to allow the model to increase the `number_of_units` itself, through [Investment Optimization](#)

The default value for this parameter is 1.

## online\_variable\_type

A selector for how the `units_on` variable is represented within the model.

**Default value:** `unit_online_variable_type_linear`

**Uses [Parameter Value Lists](#):** `unit_online_variable_type_list`

**Related [Object Classes](#):** `unit`

`online_variable_type` is a method parameter closely related to the [number\\_of\\_units](#) and can take the values "unit\_online\_variable\_type\_binary", "unit\_online\_variable\_type\_integer", "unit\_online\_variable\_type\_linear". If the binary value is chosen, the units status is modelled as a binary (classic UC). For clustered unit commitment units, the integer type is applicable. Note that if the parameter is not defined, the default will be linear. If the units status is not crucial, this can reduce the computational burden.

## operating\_points

Decomposes the flow variable into a number of separate operating segment variables. Used to in conjunction with `unit_incremental_heat_rate` and/or `unit_constraints`

**Related [Relationship Classes](#):** `unit__from_node` and `unit__to_node`

If `operating_points` is defined as an array type on a certain `unit__to_node` or `unit__from_node` flow, the corresponding `unit_flow` flow variable is decomposed into a number of sub operating segment variables, `unit_flow_op` one for each operating segment, with an additional index, `i` to reference the specific operating segment. Each value in the array represents the upper bound of the operating segment, normalized on `unit_capacity` for the corresponding `unit__to_node` or `unit__from_node` flow. `operating_points` is used in conjunction with [unit\\_incremental\\_heat\\_rate](#)

where the array dimension must match and is used to define the normalized operating point bounds for the corresponding incremental heat rate. `operating_points` is also used in conjunction with `unit_constraint` where the array dimension must match any corresponding piecewise linear `unit_flow_coefficient`. Here `operating_points` is used also to define the normalized operating point bounds for the corresponding `unit_flow_coefficients`.

Note that `operating_points` is defined on a capacity-normalized basis and the values represent the upper bound of the corresponding operating segment variable. So if `operating_points` is specified as `[0.5, 1]`, this creates two operating segments, one from zero to 50% of the corresponding `unit_capacity` and a second from 50% to 100% of the corresponding `unit_capacity`.

## output\_db\_url

Database url for SpineOpt output.

Related [Object Classes](#): `report`

The `output_db_url` parameter is the url of the database to write the results of the model run. It overrides the value of the second argument passed to `run_spineopt`.

## ramp\_down\_cost

Costs of ramping down

Related [Relationship Classes](#): `unit_from_node` and `unit_to_node`

By defining the `ramp_down_cost` parameter for a specific `unit_to_node` or `unit_from_node` relationship, a cost term will be added to the objective function whenever the unit ramps down its activity (i.e., when the `ramp_down_unit_flow` is nonzero) over the course of its operational dispatch during the current optimization window.

## ramp\_down\_limit

Limit the maximum ramp-down rate of an online unit, given as a fraction of the unit *capacity*.  
*[rampdown\_limit] = %/t*, e.g. 0.2/h

Related [Relationship Classes](#): [unit\\_from\\_node](#) and [unit\\_to\\_node](#)

The definition of the `ramp_down_limit` parameter will trigger the creation of the [Constraint on spinning downward ramps](#). It will limit the maximum decrease in the [unit\\_flow](#) variable between two consecutive timesteps for which the unit is online.

It can be defined for [unit\\_to\\_node](#) or [unit\\_from\\_node](#) relationships, as well as their counterparts for node groups. It will then impose restrictions on the `unit_flow` variables that indicate flows between the two members of the relationship for which the parameter is defined. The parameter is given as a fraction of the [unit\\_capacity](#) parameter. When the parameter is not included, the aforementioned constraint will not be created, which is equivalent to choosing a value of 1.

## ramp\_up\_cost

Costs of ramping up

Related [Relationship Classes](#): [unit\\_from\\_node](#) and [unit\\_to\\_node](#)

By defining the `ramp_up_cost` parameter for a specific [unit\\_to\\_node](#) or [unit\\_from\\_node](#) relationship, a cost term will be added to the objective function whenever the unit ramps up its activity (i.e., when the [ramp\\_up\\_unit\\_flow](#) is nonzero) over the course of its operational dispatch during the current optimization window.

## ramp\_up\_limit

Limit the maximum ramp-up rate of an online unit, given as a fraction of the `unit_capacity`.  
*[rampup\_limit] = %/t, e.g. 0.2/h*

Related [Relationship Classes](#): [unit\\_from\\_node](#) and [unit\\_to\\_node](#)

The definition of the `ramp_up_limit` parameter will trigger the creation of the [Constraint on spinning upwards ramp\\_up](#). It limits the maximum increase in the [unit\\_flow](#) variable between two consecutive timesteps for which the unit is online.

It can be defined for [unit\\_to\\_node](#) or [unit\\_from\\_node](#) relationships, as well as their counterparts for node groups. It will then impose restrictions on the `unit_flow` variables that indicate flows between the two members of the relationship for which the parameter is defined. The parameter is given as a

fraction of the `unit_capacity` parameter. When the parameter is not included, the aforementioned constraint will not be created, which is equivalent to choosing a value of 1.

For a more complete description of how ramping restrictions can be implemented, see [Ramping and Reserves](#).

## representative\_periods\_mapping

Mapping from real timelices to their corresponding representative days

Related [Object Classes](#): [temporal\\_block](#)

For representative periods with seasonal storages, `SpineOpt.jl` can be interlinked with the package `SpinePeriods.jl`. `SpinePeriods.jl` provides the [representative\\_periods\\_mapping](#) parameter, which maps each non-representative period of the whole optimization window to its representative [temporal\\_block](#). The map is organized as timeseries (indicating the start of each the non-representative period) with the names of the representative `temporal_blocks` as entries.

## reserve\_procurement\_cost

Procurement cost for reserves

Related [Relationship Classes](#): [unit\\_from\\_node](#) and [unit\\_to\\_node](#)

By defining the [reserve\\_procurement\\_cost](#) parameter for a specific [unit\\_to\\_node](#) or [unit\\_from\\_node](#) relationship, a cost term will be added to the objective function whenever that unit is used over the course of the operational dispatch during the current optimization window.

## resolution

Temporal resolution of the `temporal_block`. Essentially, divides the period between `block_start` and `block_end` into `TimeSlices` with the input `resolution`.

**Default value:** `Dict{String, Any}{"data" => "1h", "type" => "duration"}`

Related [Object Classes](#): [temporal\\_block](#)

This parameter specifies the resolution of the temporal block, or in other words: the length of the timesteps used in the optimization run. Generally speaking, variables and constraints are generated for each timestep of an optimization. For example, the nodal balance constraint must hold for each timestep.

An array of duration values can be used to have a resolution that varies with time itself. It can for example be used when uncertainty in one of the inputs rises as the optimization moves away from the model start. Think of a forecast of for instance wind power generation, which might be available in quarter hourly detail for one day in the future, and in hourly detail for the next two days. It is possible to take a quarter hourly resolution for the full horizon of three days. However, by lowering the temporal resolution after the first day, the computational burden is lowered substantially.

## right\_hand\_side

The right-hand side, constant term in a [unit\\_constraint](#). Can be time-dependent and used e.g. for complicated efficiency approximations.

**Default value:** 0.0

Related [Object Classes](#): [unit\\_constraint](#)

Used to specify the right-hand-side, constant term in a [unit\\_constraint](#). See also [unit\\_constraint](#).

## roll\_forward

Defines how much the model moves ahead in time between solves in a rolling optimization. Without this parameter, everything is solved in as a single optimization.

Related [Object Classes](#): [model](#)

This parameter defines how much the optimization window rolls forward in a rolling horizon optimization and should be expressed as a duration. In a rolling horizon optimization, a (small) part of the model is optimized at each iteration, after which the window rolls forward to optimize a different

part. Overlap between consecutive optimization windows is possible. In the practical approaches presented in [Temporal Framework](#), the rolling window optimization will be explained in more detail. The default value of this parameter is the entire model time horizon, which leads to a single optimization for the entire time horizon.

## shut\_down\_cost

Costs of shutting down a 'sub unit', e.g. EUR/shutdown.

Related [Object Classes](#): [unit](#)

By defining the shut\_down\_cost parameter for a specific [unit](#), a cost term will be added to the objective function whenever this unit shuts down over the course of its operational dispatch during the current optimization window.

## start\_up\_cost

Costs of starting up a 'sub unit', e.g. EUR/startup.

Related [Object Classes](#): [unit](#)

By defining the start\_up\_cost parameter for a specific [unit](#), a cost term will be added to the objective function whenever this unit starts up over the course of its operational dispatch during the current optimization window.

## state\_coeff

Represents the [commodity](#) content of a [node\\_state](#) variable in respect to the [unit\\_flow](#) and [connection\\_flow](#) variables. Essentially, acts as a coefficient on the [node\\_state](#) variable in the [:node\\_injection](#) constraint.

Default value: 0.0

Related [Object Classes](#): [node](#)

The [state\\_coeff](#) parameter acts as a coefficient for the [node\\_state](#) variable in the [node injection constraint](#). Essentially, it tells how the [node\\_state](#) variable should be treated in relation to the [commodity](#) flows and [demand](#), and can be used for e.g. scaling or unit conversions. For most use-cases a [state\\_coeff](#) parameter value of `1.0` should suffice, e.g. having a MWh storage connected to MW flows in a model with hour as the basic unit of time.

Note that in order for the [state\\_coeff](#) parameter to have an impact, the [node](#) must first have a [node\\_state](#) variable to begin with, defined using the [has\\_state](#) parameter. By default, the [state\\_coeff](#) is set to zero as a precaution, so that the user always has to set its value explicitly for it to have an impact on the model.

## stochastic\_scenario\_end

A [Duration](#) for when a [stochastic\\_scenario](#) ends and its [child\\_stochastic\\_scenarios](#) start. Values are interpreted relative to the start of the current solve, and if no value is given, the [stochastic\\_scenario](#) is assumed to continue indefinitely.

Related [Relationship Classes](#): [stochastic\\_structure\\_\\_stochastic\\_scenario](#)

The [stochastic\\_scenario\\_end](#) is a [Duration](#)-type parameter, defining when a [stochastic\\_scenario](#) ends relative to the start of the current optimization. As it is a parameter for the [stochastic\\_structure\\_\\_stochastic\\_scenario](#) relationship, different [stochastic\\_structures](#) can have different values for the same [stochastic\\_scenario](#), making it possible to define slightly different [stochastic\\_structures](#) using the same [stochastic\\_scenarios](#). See the [Stochastic Framework](#) section for more information about how different [stochastic\\_structures](#) interact in *SpineOpt.jl*.

When a [stochastic\\_scenario](#) ends at the point in time defined by the [stochastic\\_scenario\\_end](#) parameter, it spawns its children according to the [parent\\_stochastic\\_scenario\\_\\_child\\_stochastic\\_scenario](#) relationship. Note that the children will be inherently assumed to belong to the same [stochastic\\_structure](#) their parent belonged to, even without explicit [stochastic\\_structure\\_\\_stochastic\\_scenario](#) relationships! Thus, you might need to define the [weight\\_relative\\_to\\_parents](#) parameter for the children.

If no [stochastic\\_scenario\\_end](#) is defined, the [stochastic\\_scenario](#) is assumed to go on indefinitely.

## storage\_investment\_cost



Determines the investment cost per unit state\_cap over the investment life of a storage

Related [Object Classes](#): [node](#)

By defining the storage\_investment\_cost parameter for a specific node, a cost term will be added to the objective function whenever a storage investment is made during the current optimization window.

## storage\_investment\_lifetime

Minimum lifetime for storage investment decisions.

Related [Object Classes](#): [node](#)

Duration parameter that determines the minimum duration of storage investment decisions. Once a storage has been invested-in, it must remain invested-in for storage\_investment\_lifetime. Note that storage\_investment\_lifetime is a dynamic parameter that will impact the amount of solution history that must remain available to the optimisation in each step - this may impact performance.

See also [Investment Optimization](#) and [candidate\\_storages](#)

## storage\_investment\_variable\_type

Determines whether the storage investment variable is continuous (usually representing capacity) or integer (representing discrete units invested)

**Default value:** variable\_type\_integer

Uses [Parameter Value Lists](#): [variable\\_type\\_list](#)

Related [Object Classes](#): [node](#)

Within an investments problem `storage_investment_variable_type` determines the storage investment decision variable type. Since a `node`'s `node_state` will be limited to the product of the investment variable and the corresponding `node_state_cap` and since `candidate_storages` represents the upper bound of the storage investment decision variable, `storage_investment_variable_type` thus determines what the investment decision represents. If `storage_investment_variable_type` is integer or binary, then `candidate_storages` represents the maximum number of discrete storages that may be invested-in. If `storage_investment_variable_type` is continuous, `candidate_storages` is more analagous to a capacity with `node_state_cap` being analagous to a scaling parameter. For example, if `storage_investment_variable_type = integer`, `candidate_storages = 4` and `node_state_cap = 1000 MWh`, then the investment decision is how many 1000h MW storages to build. If `storage_investment_variable_type = continuous`, `candidate_storages = 1000` and `node_state_cap = 1 MWh`, then the investment decision is how much storage capacity to build. Finally, if `storage_investment_variable_type = integer`, `candidate_storages = 10` and `node_state_cap = 100 MWh`, then the investment decision is how many 100MWh storage blocks to build.

See also [Investment Optimization](#) and [candidate\\_storages](#).

## tax\_in\_unit\_flow

Tax costs for incoming `unit_flows` on this `node`. E.g. EUR/MWh.

Default value: 0.0

Related [Object Classes](#): [node](#)

By defining the `tax_in_unit_flow` parameter for a specific `node`, a cost term will be added to the objective function to account the taxes associated with all `unit_flow` variables with direction `to_node` over the course of the operational dispatch during the current optimization window.

## tax\_net\_unit\_flow

Tax costs for net incoming and outgoing `unit_flows` on this `node`. Incoming flows accrue positive net taxes, and outgoing flows accrue negative net taxes.

Default value: 0.0

Related [Object Classes](#): [node](#)

By defining the [tax\\_net\\_unit\\_flow](#) parameter for a specific [node](#), a cost term will be added to the objective function to account the taxes associated with the net total of all [unit\\_flow](#) variables with the direction [to\\_node](#) for this specific node minus all [unit\\_flow](#) variables with direction [from\\_node](#).

## tax\_out\_unit\_flow

Tax costs for outgoing [unit\\_flows](#) from this [node](#). E.g. EUR/MWh.

Default value: 0.0

Related [Object Classes](#): [node](#)

By defining the [tax\\_out\\_unit\\_flow](#) parameter for a specific [node](#), a cost term will be added to the objective function to account the taxes associated with all [unit\\_flow](#) variables with direction [from\\_node](#) over the course of the operational dispatch during the current optimization window.

## unit\_availability\_factor

Availability of the [unit](#), acting as a multiplier on its [unit\\_capacity](#). Typically between 0-1.

Default value: 1.0

Related [Object Classes](#): [unit](#)

To indicate that a unit is only available to a certain extent or at certain times of the optimization, the [unit\\_availability\\_factor](#) can be used. A typical use case could be an availability timeseries for a variable renewable energy source. By default the availability factor is set to [1](#). The availability is, among others, used in the [constraint\\_units\\_available](#).

## unit\_capacity

Maximum `unit_flow` capacity of a single 'sub\_unit' of the `unit`.

Related [Relationship Classes](#): `unit_from_node` and `unit_to_node`

To set an upper bound on the commodity flow of a unit in a certain direction, the `unit_capacity` constraint needs to be defined on a `unit_to_node` or `unit_from_node` relationship. By defining the parameter, the `unit_flow` variables to or from a `node` or a group of nodes will be constrained by the `capacity constraint`.

Note that if the `unit_capacity` parameter is defined on a node group, the sum of all `unit_flows` within the specified node group will be constrained by the `unit_capacity`.

## unit\_conv\_cap\_to\_flow

Optional coefficient for `unit_capacity` unit conversions in the case the `unit_capacity` value is incompatible with the desired `unit_flow` units.

Default value: 1.0

Related [Relationship Classes](#): `unit_from_node` and `unit_to_node`

The `unit_conv_cap_to_flow`, as defined for a `unit_to_node` or `unit_from_node`, allows the user to align between `unit_flow` variables and the `unit_capacity` parameter, which may be expressed in different units. An example would be when the `unit_capacity` is expressed in GWh, while the demand on the node is expressed in MWh. In that case, a `unit_conv_cap_to_flow` parameter of 1000 would be applicable.

## unit\_flow\_coefficient

Coefficient of a `unit_flow` variable for a custom `unit_constraint`.

Default value: 0.0

Related [Relationship Classes](#): [unit\\_from\\_node\\_unit\\_constraint](#) and [unit\\_to\\_node\\_unit\\_constraint](#)

The [unit\\_flow\\_coefficient](#) is an optional parameter that can be used to include the [unit\\_flow](#) or [unit\\_flow\\_op](#) variables from or to a [node](#) in a [unit\\_constraint](#) via the [unit\\_from\\_node\\_unit\\_constraint](#) and [unit\\_to\\_node\\_unit\\_constraint](#) relationships. Essentially, [unit\\_flow\\_coefficient](#) appears as a coefficient for the [unit\\_flow](#) and [unit\\_flow\\_op](#) variables from or to the [node](#) in the [unit constraint](#).

Note that the [unit\\_flow\\_op](#) variables are a bit of a special case, defined using the [operating\\_points](#) parameter.

## unit\_idle\_heat\_rate

Flow from node1 per unit time and per [units\\_on](#) that results in no additional flow to node2

Default value: 0.0

Related [Relationship Classes](#): [unit\\_node\\_node](#)

Used to implement the no-load or idle heat rate of a unit. This is the y-axis offset of the heat rate function and is the fuel consumed per unit time when a unit is online and that results in no additional output. This is defined on the [unit\\_\\_node\\_\\_node](#) relationship and it is assumed that the input flow from node 1 represents fuel consumption and the output flow to node 2 is the electrical output. While the units depend on the data, [unit\\_idle\\_heat\\_rate](#) is generally expressed in GJ/hr. Used in conjunction with [unit\\_incremental\\_heat\\_rate](#). [unit\\_idle\\_heat\\_rate](#) is only currently considered if [unit\\_incremental\\_heat\\_rate](#) is specified. A trivial [unit\\_incremental\\_heat\\_rate](#) of zero can be defined if there is no incremental heat rate.

## unit\_incremental\_heat\_rate

Standard piecewise incremental heat rate where node1 is assumed to be the fuel and node2 is assumed to be electricity. Assumed monotonically increasing. Array type or single coefficient where the number of coefficients must match the dimensions of [unit\\_operating\\_points](#)

Related [Relationship Classes](#): [unit\\_node\\_node](#)

Used to implement simple or piecewise linear incremental heat rate functions. Used in the constraint `unit_pw_heat_rate` - the input fuel flow at node 1 is the sum of the electrical MW output at node 2 times the incremental heat rate over all heat rate segments, plus the `unit_idle_heat_rate`. The units are determined by the data, but generally, incremental heat rates are given in GJ/MWh. Note that the formulation assumes a convex, monotonically increasing heat rate function. The formulation relies on optimality to load the heat rate segments in the correct order and no additional integer variables are created to enforce the correct loading order. The heat rate segment MW operating points are defined by `operating_points`.

To implement a simple incremental heat rate function, `unit_incremental_heat_rate` should be given as a simple scalar representing the incremental heat rate over the entire operating range of the unit. To implement a piecewise linear heat rate function, `unit_incremental_heat_rate` should be specified as an array type. It is then used in conjunction with the `unit` parameter `operating_points` which should also be defined as an array type of equal dimension. When defined as an array type `unit_incremental_heat_rate[i]` is the effective incremental heat rate between `operating_points[i-1]` (or zero if  $i=1$ ) and `operating_points[i]`. Note that `operating_points` is defined on a capacity-normalized basis so if `operating_points` is specified as `[0.5, 1]`, this creates two operating segments, one from zero to 50% of the corresponding `unit_capacity` and a second from 50% to 100% of the corresponding `unit_capacity`.

## unit\_investment\_cost

Investment cost per 'sub unit' built.

Related [Object Classes](#): `unit`

By defining the `unit_investment_cost` parameter for a specific `unit`, a cost term will be added to the objective function whenever a unit investment is made during the current optimization window.

## unit\_investment\_lifetime

Minimum lifetime for unit investment decisions.

Related [Object Classes](#): `unit`

Duration parameter that determines the minimum duration of `unit` investment decisions. Once a `unit` has been invested-in, it must remain invested-in for `unit_investment_lifetime`. Note that

`unit_investment_lifetime` is a dynamic parameter that will impact the amount of solution history that must remain available to the optimisation in each step - this may impact performance.

See also [Investment Optimization](#) and [candidate\\_units](#)

## unit\_investment\_variable\_type

Determines whether investment variable is integer or continuous.

**Default value:** `unit_investment_variable_type_continuous`

**Uses Parameter Value Lists:** [unit\\_investment\\_variable\\_type\\_list](#)

**Related Object Classes:** [unit](#)

Within an investments problem `unit_investment_variable_type` determines the [unit](#) investment decision variable type. Since the `unit_flows` will be limited to the product of the investment variable and the corresponding [unit\\_capacity](#) for each `unit_flow` and since [candidate\\_units](#) represents the upper bound of the investment decision variable, `unit_investment_variable_type` thus determines what the investment decision represents. If [unit\\_investment\\_variable\\_type](#) is integer or binary, then [candidate\\_units](#) represents the maximum number of discrete units that may be invested. If [unit\\_investment\\_variable\\_type](#) is continuous, `candidate_units` is more analogous to a capacity with [unit\\_capacity](#) being analogous to a scaling parameter. For example, if `unit_investment_variable_type = integer`, `candidate_units = 4` and `unit_capacity` for a particular `unit_flow = 400` MW, then the investment decision is how many 400 MW units to build. If `unit_investment_variable_type = continuous`, `candidate_units = 400` and `unit_capacity` for a particular `unit_flow = 1` MW, then the investment decision is how much capacity if this particular unit to build. Finally, if `unit_investment_variable_type = integer`, `candidate_units = 10` and `unit_capacity` for a particular `unit_flow = 50` MW, then the investment decision is many 50MW blocks of capacity of this particular unit to build.

See also [Investment Optimization](#) and [candidate\\_units](#)

## unit\_start\_flow

Flow from node1 that is incurred when a unit is started up.

Default value: 0.0

Related [Relationship Classes](#): [unit\\_\\_node\\_\\_node](#)

Used to implement unit startup fuel consumption where node 1 is assumed to be input fuel and node 2 is assumed to be output electrical energy. This is a flow from node 1 that is incurred when the value of the variable `unitsstartedup` is 1 in the corresponding time period. This flow does not result in additional output flow at node 2. Used in conjunction with [unit\\_incremental\\_heat\\_rate](#). `unit_start_flow` is only currently considered if [unit\\_incremental\\_heat\\_rate](#) is specified. A trivial [unit\\_incremental\\_heat\\_rate](#) of zero can be defined if there is no incremental heat rate.

## units\_on\_coefficient

Coefficient of a `units_on` variable for a custom `unit_constraint`.

Default value: 0.0

Related [Relationship Classes](#): [unit\\_\\_unit\\_constraint](#)

The [units\\_on\\_coefficient](#) is an optional parameter that can be used to include the `units_on` variable of a `unit` in a `unit_constraint` via the `unit__unit_constraint` relationship. Essentially, [units\\_on\\_coefficient](#) appears as a coefficient for the `units_on` variable of the `unit` in the `unit constraint`.

## units\_started\_up\_coefficient

Coefficient of a `units_started_up` variable for a custom `unit_constraint`.

Default value: 0.0

Related [Relationship Classes](#): [unit\\_\\_unit\\_constraint](#)



The `units_started_up_coefficient` is an optional parameter that can be used to include the `units_started_up` variable of a `unit` in a `unit_constraint` via the `unit_unit_constraint` relationship. Essentially, `units_started_up_coefficient` appears as a coefficient for the `units_started_up` variable of the `unit` in the `unit constraint`.

## upward\_reserve

Identifier for `nodes` providing upward reserves

**Default value:** false

**Related Object Classes:** `node`

If a `node` has a `true` `is_reserve_node` parameter, it will be treated as a reserve node in the model. To define whether the node corresponds to an upward or downward reserve commodity, the `upward_reserve` or the `downward_reserve` parameter needs to be set to `true`, respectively.

## vom\_cost

Variable operating costs of a `unit_flow` variable. E.g. EUR/MWh.

**Related Relationship Classes:** `unit_from_node` and `unit_to_node`

By defining the `vom_cost` parameter for a specific `unit`, `node`, and `direction`, a cost term will be added to the objective function to account for the variable operation and maintenance costs associated with that unit over the course of its operational dispatch during the current optimization window.

## weight

Weighting factor of the temporal block associated with the objective function

**Default value:** 1.0

Related [Object Classes](#): [temporal\\_block](#)

The `weight` variable, defined for a [temporal\\_block](#) object can be used to assign different weights to different temporal periods that are modeled. It basically determines how important a certain temporal period is in the total cost, as it enters the [Objective](#) function. The main use of this parameter is for representative periods, where each representative period represents a specific fraction of a year or so.

## weight\_relative\_to\_parents

The weight of the `stochastic_scenario` in the objective function relative to its parents.

Default value: 1.0

Related [Relationship Classes](#): [stochastic\\_structure\\_\\_stochastic\\_scenario](#)

The [weight\\_relative\\_to\\_parents](#) parameter defines how much weight the [stochastic\\_scenario](#) gets in the [Objective function](#). As a [stochastic\\_structure\\_\\_stochastic\\_scenario](#) relationship parameter, different [stochastic\\_structures](#) can use different weights for the same [stochastic\\_scenario](#). Note that every [stochastic\\_scenario](#) that appears in the [model](#) must have a [weight\\_relative\\_to\\_parents](#) defined for it related to the used [stochastic\\_structure](#)! See the [Stochastic Framework](#) section for more information about how different [stochastic\\_structures](#) interact in *SpineOpt.jl*.)

Since the [Stochastic Framework](#) in *SpineOpt.jl* supports *stochastic directed acyclic graphs* instead of simple *stochastic trees*, it is possible to define [stochastic\\_structures](#) with converging [stochastic\\_scenarios](#). In these cases, the child [stochastic\\_scenarios](#) inherit the weight of all of their parents, and the final weight that will appear in the [Objective function](#) is calculated as shown below:

```
# For root `stochastic_scenarios` (meaning no parents)

weight(scenario) = weight_relative_to_parents(scenario)

# If not a root `stochastic_scenario`

weight(scenario) = sum([weight(parent) * weight_relative_to_parents(scenario)] for p
```

The above calculation is performed starting from the roots, generation by generation, until the leaves of the *stochastic DAG*. Thus, the final weight of each [stochastic\\_scenario](#) is dependent on the [weight\\_relative\\_to\\_parents Parameters](#) of all its ancestors.

## write\_lodf\_file

A boolean flag for whether the LODF values should be written to a results file.

**Default value:** false

**Uses [Parameter Value Lists:](#)** [boolean\\_value\\_list](#)

**Related [Object Classes:](#)** [model](#)

If this parameter value is set to `true`, a diagnostics file containing all the network line outage distributions factors in CSV format will be written to the current directory.

## write\_mps\_file

A selector for writing an .mps file of the model.

**Uses [Parameter Value Lists:](#)** [write\\_mps\\_file\\_list](#)

**Related [Object Classes:](#)** [model](#)

This parameter is deprecated and will be removed in a future version.

This parameter controls when to write a diagnostic model file in MPS format. If set to `write_mps_always`, the model will always be written in MPS format to the current directory. If set to `write\_mps\_on\_no\_solve`, the MPS file will be written when the model solve terminates with a status of false. If set to `write\_mps\_never`, no file will be written

## write\_ptdf\_file

A boolean flag for whether the LODF values should be written to a results file.

**Default value:** false

**Uses [Parameter Value Lists](#):** [boolean\\_value\\_list](#)

**Related [Object Classes](#):** [model](#)

If this parameter value is set to `true`, a diagnostics file containing all the network power transfer distributions factors in CSV format will be written to the current directory.

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[« Relationship Classes](#)

[Parameter Value Lists »](#)

# Parameter Value Lists

## balance\_type\_list

**Default value:** balance\_type\_none

The [balance\\_type\\_list](#) parameter value list contains the possible values for the [balance\\_type](#) parameter.

## boolean\_value\_list

**Default value:** true

A list of boolean values (True or False).

## commodity\_physics\_list

**Default value:** commodity\_physics\_ptdf

[commodity\\_physics\\_list](#) holds the possible values for the commodity parameter [commodity\\_physics](#) parameter. See [commodity\\_physics](#) for more details

## connection\_investment\_variable\_type\_list

**Default value:** connection\_investment\_variable\_type\_integer

The [connection\\_investment\\_variable\\_type\\_list](#) holds the possible values for the type of a [connection](#)'s investment variable which may be chosen between integer or continuous.

## connection\_type\_list

**Default value:** connection\_type\_normal

`connection_type_list` holds the possible values for the [connection\\_type](#) parameter. See [connection\\_type](#) for more details

## constraint\_sense\_list

**Default value:** >=

The [constraint\\_sense\\_list](#) parameter value list contains the possible values for the [constraint\\_sense](#) parameter.

## duration\_unit\_list

**Default value:** minute

The [duration\\_unit\\_list](#) parameter value list contains the possible values for the [duration\\_unit](#) parameter.

## model\_type\_list

**Default value:** spineopt\_other

`model_type_list` holds the possible values for the model parameter [model\\_type](#) parameter. See [model\\_type\\_list](#) for more details

## node\_opf\_type\_list

**Default value:** node\_opf\_type\_reference

Houses the different possible values for the [node\\_opf\\_type](#) parameter. To identify the reference node, set `node_opf_type = :node_opf_type_reference`, while `node_opf_type = node_opf_type_normal` is the default value for non-reference nodes.

See also [powerflow](#).

## unit\_investment\_variable\_type\_list

**Default value:** unit\_investment\_variable\_type\_integer

`unit_investment_variable_type_list` holds the possible values for the type of a `unit`'s investment variable which may be chosen from integer, binary or continuous.

## `unit_online_variable_type_list`

**Default value:** `unit_online_variable_type_linear`

`unit_online_variable_type_list` holds the possible values for the type of a `unit`'s commitment status variable which may be chosen from binary, integer, or linear.

## `variable_type_list`

**Default value:** `variable_type_integer`

The `variable_type_list` parameter value list contains the possible values for the `connection_investment_variable_type` and `storage_investment_variable_type` parameters.

## `write_mps_file_list`

**Default value:** `write_mps_on_no_solve`

`_list`

This parameter value list is deprecated and will be removed in a future version.

Houses the different values for the `write_mps_file` parameter. Possible values include `write_mps_always`, `write\_mps\_on\_no\_solve`, and `write\_mps\_never`.

# Variables

## binary\_gas\_connection\_flow

**Math symbol:**  $v_{binary\_gas\_connection\_flow}$

**Indices:** (connection=conn, node=n, direction=d, stochastic\_scenario=s, t=t)

**Indices function:** binary\_gas\_connection\_flow\_indices

Binary variable with the indices node  $n$  over the connection  $conn$  in the direction  $to\_node$  for the stochastic scenario  $s$  at timestep  $t$  describing if the direction of gas flow for a pressure drive gas transfer is in the indicated direction.

## connection\_flow

**Math symbol:**  $v_{connection\_flow}$

**Indices:** (connection=conn, node=n, direction=d, stochastic\_scenario=s, t=t)

**Indices function:** connection\_flow\_indices

Commodity flow associated with node  $n$  over the connection  $conn$  in the direction  $d$  for the stochastic scenario  $s$  at timestep  $t$

## connection\_intact\_flow

**Math symbol:**  $v_{connection\_intact\_flow}$



**Indices:** (connection=conn, node=n, direction=d, stochastic\_scenario=s, t=t)

**Indices function:** connection\_intact\_flow\_indices

Represents the ptdf-based flow on connections where all investment candidate connections are present in the network.

## connections\_decommissioned

**Math symbol:**  $v_{connections\_decommissioned}$

**Indices:** (connection=conn, stochastic\_scenario=s, t=t)

**Indices function:** connections\_invested\_available\_indices

Number of decomissioned connections  $conn$  for the stochastic scenario  $s$  at timestep  $t$

## connections\_invested

**Math symbol:**  $v_{connections\_invested}$

**Indices:** (connection=conn, stochastic\_scenario=s, t=t)

**Indices function:** connections\_invested\_available\_indices

Number of connections  $conn$  invested at timestep  $t$  in for the stochastic scenario  $s$

## connections\_invested\_available

**Math symbol:**  $v_{connections\_invested\_available}$

**Indices:** (connection=conn, stochastic\_scenario=s, t=t)

**Indices function:** connections\_invested\_available\_indices

Number of invested connections  $conn$  that are available still the stochastic scenario  $s$  at timestep  $t$

## mp\_objective\_lowerbound\_indices

**Math symbol:**  $v_{mp\_objective\_lowerbound\_indices}$

**Indices:** (t=t)

**Indices function:** mp\_objective\_lowerbound\_indices

Updating lowerbound for master problem of Benders decomposition

## node\_injection

**Math symbol:**  $v_{node\_injection}$

**Indices:** (node=n, stochastic\_scenario=s, t=t)

**Indices function:** node\_injection\_indices

Commodity injections at node  $n$  for the stochastic scenario  $s$  at timestep  $t$

## node\_pressure

**Math symbol:**  $v_{node\_pressure}$

**Indices:** (node= $n$ , stochastic\_scenario= $s$ , t= $t$ )

**Indices function:** node\_pressure\_indices

Pressure at a node  $n$  for a specific stochastic scenario  $s$  and timestep  $t$ . See also: [has\\_pressure](#)

## node\_slack\_neg

**Math symbol:**  $v_{node\_slack\_neg}$

**Indices:** (node= $n$ , stochastic\_scenario= $s$ , t= $t$ )

**Indices function:** node\_slack\_indices

Negative slack variable at node  $n$  for the stochastic scenario  $s$  at timestep  $t$

## node\_slack\_pos

**Math symbol:**  $v_{node\_slack\_pos}$

**Indices:** (node= $n$ , stochastic\_scenario= $s$ , t= $t$ )

**Indices function:** node\_slack\_indices

Positive slack variable at node  $n$  for the stochastic scenario  $s$  at timestep  $t$

## node\_state

**Math symbol:**  $v_{node\_state}$

**Indices:** (node= $n$ , stochastic\_scenario= $s$ , t= $t$ )

**Indices function:** node\_state\_indices

Storage state at node  $n$  for the stochastic scenario  $s$  at timestep  $t$

## node\_voltage\_angle

**Math symbol:**  $v_{node\_voltage\_angle}$

**Indices:** (node= $n$ , stochastic\_scenario= $s$ , t= $t$ )

**Indices function:** node\_voltage\_angle\_indices

Voltage angle at a node  $n$  for a specific stochastic scenario  $s$  and timestep  $t$ . See also: [has\\_voltage\\_angle](#)

## nonspin\_ramp\_down\_unit\_flow

**Math symbol:**  $v_{nonspin\_ramp\_down\_unit\_flow}$

**Indices:** (unit= $u$ , node= $n$ , direction= $d$ , stochastic\_scenario= $s$ , t= $t$ )

**Indices function:** nonspin\_ramp\_down\_unit\_flow\_indices

Non-spinning downward reserve commodity flows of unit  $u$  at node  $n$  in the direction  $d$  for the stochastic scenario  $s$  at timestep  $t$

## nonspin\_ramp\_up\_unit\_flow

**Math symbol:**  $v_{nonspin\_ramp\_up\_unit\_flow}$

**Indices:** (unit= $u$ , node= $n$ , direction= $d$ , stochastic\_scenario= $s$ , t= $t$ )

**Indices function:** nonspin\_ramp\_up\_unit\_flow\_indices

Non-spinning upward reserve commodity flows of unit  $u$  at node  $n$  in the direction  $d$  for the stochastic scenario  $s$  at timestep  $t$

## nonspin\_units\_shut\_down

**Math symbol:**  $v_{nonspin\_units\_shut\_down}$

**Indices:** (unit= $u$ , node= $n$ , stochastic\_scenario= $s$ , t= $t$ )

**Indices function:** nonspin\_units\_shut\_down\_indices

Number of units  $u$  held available for non-spinning downward reserve provision via shutdown to node  $n$  for the stochastic scenario  $s$  at timestep  $t$

## nonspin\_units\_started\_up

**Math symbol:**  $v_{nonspin\_units\_started\_up}$

**Indices:** (unit= $u$ , node= $n$ , stochastic\_scenario= $s$ , t= $t$ )

**Indices function:** nonspin\_units\_started\_up\_indices

Number of units  $u$  held available for non-spinning upward reserve provision via startup to node  $n$  for the stochastic scenario  $s$  at timestep  $t$

## ramp\_down\_unit\_flow

**Math symbol:**  $v_{ramp\_down\_unit\_flow}$

**Indices:** (unit= $u$ , node= $n$ , direction= $d$ , stochastic\_scenario= $s$ , t= $t$ )

**Indices function:** ramp\_down\_unit\_flow\_indices

Spinning downward ramp commodity flow associated with node  $n$  of unit  $u$  with node  $n$  over the connection  $conn$  in the direction  $d$  for the stochastic scenario  $s$  at timestep  $t$

## ramp\_up\_unit\_flow

**Math symbol:**  $v_{ramp\_up\_unit\_flow}$

**Indices:** (unit= $u$ , node= $n$ , direction= $d$ , stochastic\_scenario= $s$ , t= $t$ )

**Indices function:** ramp\_up\_unit\_flow\_indices

Spinning upward ramp commodity flow associated with node  $n$  of unit  $u$  with node  $n$  over the connection  $conn$  in the direction  $d$  for the stochastic scenario  $s$  at timestep  $t$

## shut\_down\_unit\_flow

**Math symbol:**  $v_{shut\_down\_unit\_flow}$

**Indices:** (unit= $u$ , node= $n$ , direction= $d$ , stochastic\_scenario= $s$ , t= $t$ )

**Indices function:** shut\_down\_unit\_flow\_indices

Downward ramp commodity flow during shutdown associated with node  $n$  of unit  $u$  with node  $n$  over the connection  $conn$  in the direction  $d$  for the stochastic scenario  $s$  at timestep  $t$

---

## start\_up\_unit\_flow

**Math symbol:**  $v_{start\_up\_unit\_flow}$

**Indices:** (unit= $u$ , node= $n$ , direction= $d$ , stochastic\_scenario= $s$ , t= $t$ )

**Indices function:** start\_up\_unit\_flow\_indices

Upward ramp commodity flow during start-up associated with node  $n$  of unit  $u$  with node  $n$  over the connection  $conn$  in the direction  $d$  for the stochastic scenario  $s$  at timestep  $t$

## storages\_decommissioned

**Math symbol:**  $v_{storages\_decommissioned}$

**Indices:** (node= $n$ , stochastic\_scenario= $s$ , t= $t$ )

**Indices function:** storages\_invested\_available\_indices

Number of decomissioned storage nodes  $n$  for the stochastic scenario  $s$  at timestep  $t$

## storages\_invested

**Math symbol:**  $v_{storages\_invested}$

**Indices:** (node= $n$ , stochastic\_scenario= $s$ , t= $t$ )

**Indices function:** storages\_invested\_available\_indices

Number of storage nodes  $n$  invested in at timestep  $t$  for the stochastic scenario  $s$

---

## storages\_invested\_available

**Math symbol:**  $v_{storages\_invested\_available}$

**Indices:** (node= $n$ , stochastic\_scenario= $s$ , t= $t$ )

**Indices function:** storages\_invested\_available\_indices

Number of invested storage nodes  $n$  that are available still the stochastic scenario  $s$  at timestep  $t$

## unit\_flow

**Math symbol:**  $v_{unit\_flow}$

**Indices:** (unit= $u$ , node= $n$ , direction= $d$ , stochastic\_scenario= $s$ , t= $t$ )

**Indices function:** unit\_flow\_indices

Commodity flow associated with node  $n$  over the unit  $u$  in the direction  $d$  for the stochastic scenario  $s$  at timestep  $t$

## unit\_flow\_op

**Math symbol:**  $v_{unit\_flow\_op}$

**Indices:** (unit= $u$ , node= $n$ , direction= $d$ , i= $i$ , stochastic\_scenario= $s$ , t= $t$ )

**Indices function:** unit\_flow\_op\_indices

Contribution of the unit flow associated with operating point  $i$

---



## units\_available

**Math symbol:**  $v_{units\_available}$

**Indices:** (unit= $u$ , stochastic\_scenario= $s$ , t= $t$ )

**Indices function:** units\_on\_indices

Number of available units  $u$  for the stochastic scenario  $s$  at timestep  $t$

## units\_invested

**Math symbol:**  $v_{units\_invested}$

**Indices:** (unit= $u$ , stochastic\_scenario= $s$ , t= $t$ )

**Indices function:** units\_invested\_available\_indices

Number of units  $u$  for the stochastic scenario  $s$  invested in at timestep  $t$

## units\_invested\_available

**Math symbol:**  $v_{units\_invested\_available}$

**Indices:** (unit= $u$ , stochastic\_scenario= $s$ , t= $t$ )

**Indices function:** units\_invested\_available\_indices

Number of invested units  $u$  that are available still the stochastic scenario  $s$  at timestep  $t$

---

## units\_mothballed

**Math symbol:**  $v_{units\_mothballed}$

**Indices:** (unit= $u$ , stochastic\_scenario= $s$ , t= $t$ )

**Indices function:** units\_invested\_available\_indices

Number of units  $u$  for the stochastic scenario  $s$  mothballed at timestep  $t$

## units\_on

**Math symbol:**  $v_{units\_on}$

**Indices:** (unit= $u$ , stochastic\_scenario= $s$ , t= $t$ )

**Indices function:** units\_on\_indices

Number of online units  $u$  for the stochastic scenario  $s$  at timestep  $t$

## units\_shut\_down

**Math symbol:**  $v_{units\_shut\_down}$

**Indices:** (unit= $u$ , stochastic\_scenario= $s$ , t= $t$ )

**Indices function:** units\_on\_indices

Number of units  $u$  for the stochastic scenario  $s$  that switched to offline status at timestep  $t$

---

# units\_started\_up

**Math symbol:**  $v_{units\_started\_up}$

**Indices:** (unit= $u$ , stochastic\_scenario= $s$ , t= $t$ )

**Indices function:** units\_on\_indices

Number of units  $u$  for the stochastic scenario  $s$  that switched to online status at timestep  $t$

---

[« Parameter Value Lists](#)

[Constraints »](#)

# Constraints

## Balance constraint

### Nodal balance

In **SpineOpt**, **node** is the place where an energy balance is enforced. As universal aggregators, they are the glue that brings all components of the energy system together. An energy balance is created for each **node** for all **node\_stochastic\_time\_indices**, unless the **balance\_type** parameter of the node takes the value **balance\_type\_none** or if the node in question is a member of a node group, for which the **balance\_type** is **balance\_type\_group**. The parameter **nodal\_balance\_sense** defaults to equality, but can be changed to allow overproduction (**nodal\_balance\_sense** **>=**) or underproduction (**nodal\_balance\_sense** **<=**). The energy balance is enforced by the following constraint:

$$\begin{aligned}
 & v_{node\_injection}(n, s, t) \\
 & + \sum_{\substack{(conn, n', d_{in}, s, t) \in connection\_flow\_indices: \\ d_{out} == to\_node}} v_{connection\_flow}(conn, n', d_{in}, s, t) \\
 & - \sum_{\substack{(conn, n', d_{out}, s, t) \in connection\_flow\_indices: \\ d_{out} == from\_node}} v_{connection\_flow}(conn, n', d_{out}, s, t) \\
 & + v_{node\_slack\_pos}(n, s, t) \\
 & - v_{node\_slack\_neg}(n, s, t) \\
 & \{>=, ==, <=\} \\
 & 0 \\
 & \forall (n, s, t) \in node\_stochastic\_time\_indices : \\
 & p_{balance\_type}(n) \neq balance\_type\_none \\
 & \nexists ng \in groups(n) : balance\_type\_group
 \end{aligned}$$

The constraint consists of the **node injections**, the net **connection flows** and **node slack variables**.

### Node injection

The node injection itself represents all local production and consumption, represented by the sum of all connected unit flows and the nodal demand. The node injection is created for each node in the network (unless the node is only used for parameter aggregation purposes, see [Introduction to groups of objects](#)).

$$\begin{aligned}
 & v_{node\_injection}(n, s, t) \\
 & == \\
 & + \sum_{\substack{(u, n', d_{in}, s, t) \in unit\_flow\_indices: \\ d_{out} == to\_node}} v_{unit\_flow}(u, n', d_{in}, s, t) \\
 & - \sum_{\substack{(u, n', d_{out}, s, t) \in unit\_flow\_indices: \\ d_{out} == from\_node}} v_{unit\_flow}(u, n', d_{out}, s, t) \\
 & - p_{demand}(n, s, t) \\
 & \forall (n, s, t) \in node\_stochastic\_time\_indices
 \end{aligned}$$

### Node injection with storage capability

If a node corresponds to a storage node, the parameter **has\_state** should be set to **true** for this node. In this case the nodal injection will translate to the following constraint:

$$\begin{aligned}
& v_{node\_injection}(n, s, t) \\
& == \\
& (v_{node\_state}(n, s, t\_before) \\
& - v_{node\_state}(n, s, t) \cdot p_{state\_coeff}(n, s, t)) \\
& / \Delta t_{after} \\
& - v_{node\_state}(n, s, t) \cdot p_{frac\_state\_loss}(n, s, t) \\
& + \sum_{\substack{(n2, s, t) \in node\_state\_indices: \\ \exists diff\_coeff(n2, n)}} v_{node\_state}(n2, s, t) \\
& - \sum_{\substack{(n2, s, t) \in node\_state\_indices: \\ \exists diff\_coeff(n, n2)}} v_{node\_state}(n2, s, t) \\
& + \sum_{\substack{(u, n', d_{in}, s, t) \in unit\_flow\_indices: \\ d_{out} == to\_node}} v_{unit\_flow}(u, n', d_{in}, s, t) \\
& - \sum_{\substack{(u, n', d_{out}, s, t) \in unit\_flow\_indices: \\ d_{out} == from\_node}} v_{unit\_flow}(u, n', d_{out}, s, t) \\
& - demand(n, s, t) \\
& \forall (n, t) \in node\_time\_indices : p_{has\_state}(n) \\
& \forall s \in stochastic\_scenario\_path \\
& t_{before} \in t\_before\_t(t\_after = t)
\end{aligned}$$

Note that for simplicity, the stochastic path is assumed to be known. In the constraint `constraint_node_injection.jl` the active stochastic paths of all involved variables is retrieved beforehand.

## Node state capacity

To limit the storage content, the  $v_{node\_state}$  variable needs be constrained by the following equation:

$$\begin{aligned}
& v_{node\_state}(n, s, t) \\
& \leq p_{node\_state\_cap}(n, s, t) \\
& \forall (n, s, t) \in node\_stochastic\_time\_indices : \\
& p_{has\_state}(n)
\end{aligned}$$

The discharging and charging behavior of storage nodes can be described through unit(s), representing the link between the storage node and the supply node. Note that the dis-/charging efficiencies and capacities are properties of these units. See [the capacity constraint](#) and [the unit flow ratio constraints](#)

## Cyclic condition on node state variable

To ensure that the node state at the end of the optimization is at least the same value as the initial value at the beginning of the optimization (or higher), the cyclic node state constraint can be used by setting the `cyclic_condition` of a `node_temporal_block` to `true`. This triggers the following cyclic constraint:

$$\begin{aligned}
&v_{node\_state}(n, s, t) \\
&\geq v_{node\_state}(n, s, t) \\
&\forall (n, tb) \in p_{cyclic\_condition}(n, tb) : \\
&\{p_{cyclic\_condition}(n, tb) == true, \\
&phas\_state(n)\} \\
&\forall (n', t_{initial}) \in node\_time\_indices : \\
&\{n' == n, \\
&t_{initial} == t\_before\_t(t\_after = first(t \in tb)), \\
&\forall (n', t_{last}) \in node\_time\_indices : \\
&n' == n, \\
&t_{last} == last(t \in tb)) \\
&\forall s \in stochastic\_path
\end{aligned}$$

## Unit operation

In the following, the operational constraints on the variables associated with units will be elaborated on. The static constraints, in contrast to the dynamic constraints, are addressing constraints without sequential time-coupling. It should however be noted that static constraints can still perform temporal aggregation.

## Static constraints

The fundamental static constraints for units within SpineOpt relate to the relationships between commodity flows from and to units and to limits on the unit flow capacity.

### Conversion constraint / limiting flow shares in process / relationship in process

A [unit](#) can have different commodity flows associated with it. The most simple relationship between these flows is a linear relationship between input and/or output nodes/node groups. SpineOpt holds constraints for each combination of flows and also for the type of relationship, i.e. whether it is a maximum, minimum or fixed ratio between commodity flows. Note that node groups can be used in order to aggregate flows, i.e. to give a ratio between a combination of units flows.

### Ratios between output and input flows of a unit

By defining the parameters [fix\\_ratio\\_out\\_in\\_unit\\_flow](#), [max\\_ratio\\_out\\_in\\_unit\\_flow](#) or [min\\_ratio\\_out\\_in\\_unit\\_flow](#), a ratio can be set between **outgoing** and **incoming** flows from and to a unit. Whenever there is only a single input node and a single output node, this relationship relates to the notion of an efficiency. Also, the ratio equation can for instance be used to relate emissions to input primary fuel flows. In the most general form of the equation, two node groups are defined (an input node group  $ng_{in}$  and an output node group  $ng_{out}$ ), and a linear relationship is expressed between both node groups. Note that whenever the relationship is specified between groups of multiple nodes, there remains a degree of freedom regarding the composition of the input node flows within group  $ng_{in}$  and the output node flows within group  $ng_{out}$ .

The constraint given below enforces a fixed, maximum or minimum ratio between outgoing and incoming [unit\\_flow](#). Note that the potential node groups, that the parameters [fix\\_ratio\\_out\\_in\\_unit\\_flow](#), [max\\_ratio\\_out\\_in\\_unit\\_flow](#) and [min\\_ratio\\_out\\_in\\_unit\\_flow](#) defined on, are getting internally expanded to the members of the node group within the `unit_flow_indices`.

$$\begin{aligned}
& \sum_{\substack{(u,n,d,s,t_{out}) \in \text{unit\_flow\_indices:} \\ (u,n,d,s,t_{out}) \in (u,ng_{out},:to\_node,s,t)}} v_{\text{unit\_flow}}(u,n,d,s,t_{out}) \cdot \Delta t_{out} \\
& \{ \\
& \quad == p_{\text{fix\_ratio\_out\_in\_unit\_flow}}(u,ng_{out},ng_{in},s,t), \\
& \quad <= p_{\text{max\_ratio\_out\_in\_unit\_flow}}(u,ng_{out},ng_{in},s,t), \\
& \quad >= p_{\text{min\_ratio\_out\_in\_unit\_flow}}(u,ng_{out},ng_{in},s,t) \\
& \} \\
& \cdot \sum_{\substack{(u,n,d,s,t_{in}) \in \text{unit\_flow\_indices:} \\ (u,n,d,s,t_{in}) \in (u,ng_{in},:from\_node,s,t)}} v_{\text{unit\_flow}}(u,n,d,s,t_{in}) \cdot \Delta t_{in} \\
& + P_{\{fix,max,min\}\text{-units\_on\_coefficient\_out\_in}}(u,ng_{out},ng_{in},s,t) \\
& \sum_{\substack{(u,s,t_{\text{units\_on}}) \in \text{units\_on\_indices:} \\ (u,s,t_{\text{units\_on}}) \in (u,s,t)}} v_{\text{units\_on}}(u,s,t_{\text{units\_on}}) \\
& \cdot \min(\Delta t_{\text{units\_on}}, \Delta t) \\
& \forall (u,ng_{out},ng_{in}) \in \text{ind}(p_{\{fix,max,min\}\text{-ratio\_out\_in\_unit\_flow}}), \\
& \forall t \in \text{time\_slices}, \forall s \in \text{stochastic\_path}
\end{aligned}$$

Note that a right-hand side constant coefficient associated with the variable `units_on` can optionally be included, triggered by the existence of the `fix_units_on_coefficient_out_in`, `max_units_on_coefficient_out_in`, `min_units_on_coefficient_out_in`, respectively.

### Ratios between input and output flows of a unit

Similarly to the ratio between outgoing and incoming unit flows, a ratio can also be defined in reverse between **incoming** and **outgoing** flows.

$$\begin{aligned}
& \sum_{\substack{(u,n,d,s,t_{in}) \in \text{unit\_flow\_indices:} \\ (u,n,d,s,t_{in}) \in (u,ng_{in},:from\_node,s,t)}} v_{\text{unit\_flow}}(u,n,d,s,t_{in}) \cdot \Delta t_{in} \\
& \{ \\
& \quad == p_{\text{fix\_ratio\_in\_out\_unit\_flow}}(u,ng_{in},ng_{out},s,t), \\
& \quad <= p_{\text{max\_ratio\_in\_out\_unit\_flow}}(u,ng_{in},ng_{out},s,t), \\
& \quad >= p_{\text{min\_ratio\_in\_out\_unit\_flow}}(u,ng_{in},ng_{out},s,t) \\
& \} \\
& \cdot \sum_{\substack{(u,n,d,s,t_{out}) \in \text{unit\_flow\_indices:} \\ (u,n,d,s,t_{out}) \in (u,ng_{out},:to\_node,s,t)}} v_{\text{unit\_flow}}(u,n,d,s,t_{out}) \cdot \Delta t_{out} \\
& + P_{\{fix,max,min\}\text{-units\_on\_coefficient\_in\_out}}(u,ng_{in},ng_{out},s,t) \\
& \sum_{\substack{(u,s,t_{\text{units\_on}}) \in \text{units\_on\_indices:} \\ (u,s,t_{\text{units\_on}}) \in (u,s,t)}} v_{\text{units\_on}}(u,s,t_{\text{units\_on}}) \\
& \cdot \min(\Delta t_{\text{units\_on}}, \Delta t) \\
& \forall (u,ng_{in},ng_{out}) \in \text{ind}(p_{\{fix,max,min\}\text{-ratio\_in\_out\_unit\_flow}}), \\
& \forall t \in \text{time\_slices}, \forall s \in \text{stochastic\_path}
\end{aligned}$$

Note that a right-hand side constant coefficient associated with the variable `units_on` can optionally be included, triggered by the existence of the `fix_units_on_coefficient_in_out`, `max_units_on_coefficient_in_out`, `min_units_on_coefficient_in_out`, respectively.

### Ratios between input and input flows of a unit

Similarly to the [ratio between outgoing and incoming units flows](#), one can also define a fixed, maximum or minimum ratio between **incoming** flows of a units.

$$\begin{aligned}
& \sum_{\substack{(u,n,d,s,t_{in1}) \in \text{unit\_flow\_indices:} \\ (u,n,d,s,t_{in1}) \in (u,ng_{in1},:from\_node,s,t)}} v_{unit\_flow}(u,n,d,s,t_{in1}) \cdot \Delta t_{in1} \\
& \{ \\
& \quad == p_{fix\_ratio\_in\_in\_unit\_flow}(u,ng_{in1},ng_{in2},s,t), \\
& \quad <= p_{max\_ratio\_in\_in\_unit\_flow}(u,ng_{in1},ng_{in2},s,t), \\
& \quad >= p_{min\_ratio\_in\_in\_unit\_flow}(u,ng_{in1},ng_{in2},s,t) \\
& \} \\
& \cdot \sum_{\substack{(u,n,d,s,t_{in2}) \in \text{unit\_flow\_indices:} \\ (u,n,d,s,t_{in2}) \in (u,ng_{in2},:from\_node,s,t)}} v_{unit\_flow}(u,n,d,s,t_{in2}) \cdot \Delta t_{in2} \\
& + p_{\{fix,max,min\}\_units\_on\_coefficient\_in\_in}(u,ng_{in1},ng_{in2},s,t) \\
& \quad \sum_{\substack{(u,s,t_{units\_on}) \in \text{units\_on\_indices:} \\ (u,s,t_{units\_on}) \in (u,s,t)}} v_{units\_on}(u,s,t_{units\_on}) \\
& \quad \cdot \min(\Delta t_{units\_on}, \Delta t) \\
& \forall (u,ng_{in1},ng_{in2}) \in ind(p_{\{fix,max,min\}\_ratio\_in\_in\_unit\_flow}), \\
& \forall t \in \text{time\_slices}, \forall s \in \text{stochastic\_path}
\end{aligned}$$

Note that a right-hand side constant coefficient associated with the variable `units_on` can optionally be included, triggered by the existence of the `fix_units_on_coefficient_in_in`, `max_units_on_coefficient_in_in`, `min_units_on_coefficient_in_in`, respectively.

### Ratios between output and output flows of a unit

Similarly to the [ratio between outgoing and incoming units flows](#), one can also define a fixed, maximum or minimum ratio between **out**going flows of a units.

$$\begin{aligned}
& \sum_{\substack{(u,n,d,s,t_{out1}) \in \text{unit\_flow\_indices:} \\ (u,n,d,s,t_{out1}) \in (u,ng_{out1},:to\_node,s,t)}} v_{unit\_flow}(u,n,d,s,t_{out1}) \cdot \Delta t_{out1} \\
& \{ \\
& \quad == p_{fix\_ratio\_out\_out\_unit\_flow}(u,ng_{out1},ng_{out2},s,t), \\
& \quad <= p_{max\_ratio\_out\_out\_unit\_flow}(u,ng_{out1},ng_{out2},s,t), \\
& \quad >= p_{min\_ratio\_out\_out\_unit\_flow}(u,ng_{out1},ng_{out2},s,t) \\
& \} \\
& \cdot \sum_{\substack{(u,n,d,s,t_{out2}) \in \text{unit\_flow\_indices:} \\ (u,n,d,s,t_{out2}) \in (u,ng_{out2},:to\_node,s,t)}} v_{unit\_flow}(u,n,d,s,t_{out2}) \cdot \Delta t_{out2} \\
& + p_{\{fix,max,min\}\_units\_on\_coefficient\_out\_out}(u,ng_{out1},ng_{out2},s,t) \\
& \quad \sum_{\substack{(u,s,t_{units\_on}) \in \text{units\_on\_indices:} \\ (u,s,t_{units\_on}) \in (u,s,t)}} v_{units\_on}(u,s,t_{units\_on}) \\
& \quad \cdot \min(\Delta t_{units\_on}, \Delta t) \\
& \forall (u,ng_{out1},ng_{out2}) \in ind(p_{\{fix,max,min\}\_ratio\_out\_out\_unit\_flow}), \\
& \forall t \in \text{time\_slices}, \forall s \in \text{stochastic\_path}
\end{aligned}$$

Note that a right-hand side constant coefficient associated with the variable `units_on` can optionally be included, triggered by the existence of the `fix_units_on_coefficient_out_out`, `max_units_on_coefficient_out_out`, `min_units_on_coefficient_out_out`, respectively.

### Bounds on the unit capacity

In a multi-commodity setting, there can be different commodities entering/leaving a certain technology/unit. These can be energy-related commodities (e.g., electricity, natural gas, etc.), emissions, or other commodities (e.g., water, steel). The `unit_capacity` be specified for at least one `unit_to_node` or `unit_from_node` relationship, in order to trigger a constraint on the maximum commodity flows to this location in each time step. When desirable, the capacity can be specified for a group of nodes (e.g. combined capacity for multiple products).



$$\begin{aligned}
& \sum_{\substack{(u,n,d,s,t') \in \text{unit\_flow\_indices:} \\ (u,n,d,s,t') \in (u,ng,d,s,t)}} v_{\text{unit\_flow}}(u, n, d, s, t') \cdot \Delta t' \\
& \leq p_{\text{unit\_capacity}}(u, ng, d, s, t) \\
& \cdot p_{\text{unit\_conv\_cap\_to\_flow}}(u, ng, d, s, t) \\
& \cdot \sum_{\substack{(u,s,t_{\text{units\_on}}) \in \text{units\_on\_indices:} \\ (u,\Delta t_{\text{units\_on}}) \in (u,t)}} v_{\text{units\_on}}(u, s, t_{\text{units\_on}}) \\
& \cdot \min(t_{\text{units\_on}}, \Delta t) \\
& \forall (u, ng, d) \in \text{ind}(p_{\text{unit\_capacity}}), \\
& \forall t \in \text{time\_slices}, \\
& \forall s \in \text{stochastic\_path}
\end{aligned}$$

Note that the conversion factor [unit\\_conv\\_cap\\_to\\_flow](#) has a default value of 1, but can be adjusted in case the unit of measurement for the capacity is different to the unit flows unit of measurement.

## Dynamic constraints

### Commitment constraints

For modeling certain technologies/units, it is important to not only have [unit\\_flow](#) variables of different commodities, but also model the online ("commitment") status of the unit/technology at every time step. Therefore, an additional variable [units\\_on](#) is introduced. This variable represents the number of online units of that technology (for a normal unit commitment model, this variable might be a binary, for investment planning purposes, this might also be an integer or even a continuous variable). To define the type of a commitment variable, see [online\\_variable\\_type](#). Commitment variables will be introduced by the following constraints (with corresponding parameters):

- constraint on [units\\_on](#)
- constraint on [units\\_available](#)
- constraint on the unit state transition
- constraint on the minimum operating point
- constraint on minimum down time
- constraint on minimum up time
- constraint on ramp rates

### Bound on online units

The number of online units need to be restricted to the number of available units:

$$\begin{aligned}
& v_{\text{units\_on}}(u, s, t) \\
& \leq v_{\text{units\_available}}(u, s, t) \\
& \forall (u, s, t) \in \text{units\_on\_indices}
\end{aligned}$$

### Bound on available units

The number of available units itself is constrained by the parameters [unit\\_availability\\_factor](#) and [number\\_of\\_units](#), and the variable number of invested units [units\\_invested\\_available](#):

$$\begin{aligned}
& v_{\text{units\_available}}(u, s, t) \\
& = p_{\text{unit\_availability\_factor}}(u, s, t) \\
& \cdot (p_{\text{number\_of\_units}}(u, s, t) \\
& + \sum_{(u,s,t) \in \text{units\_invested\_available\_indices}} v_{\text{units\_invested\_available}}(u, s, t)) \\
& \forall (u, s, t) \in \text{units\_on\_indices}
\end{aligned}$$

The investment formulation is described in chapter [Investments](#).

### Unit state transition

The units on status is constrained by shutting down and starting up actions. This transition is defined as follows:

$$\begin{aligned}
& v_{units\_on}(u, s, t_{after}) \\
& - v_{units\_started\_up}(u, s, t_{after}) \\
& + v_{units\_shut\_down}(u, s, t_{after}) \\
& == v_{units\_on}(u, s, t_{before}) \\
& \forall (u, s, t_{after}) \in units\_on\_indices, \\
& \forall t_{before} \in t\_before\_t(t\_after = t_{after}) : t_{before} \in units\_on\_indices
\end{aligned}$$

### Constraint on minimum operating point

The minimum operating point of a unit can be based on the [unit\\_flows](#) of input or output nodes/node groups ng:

$$\begin{aligned}
& \sum_{\substack{(u,n,d,s,t') \in unit\_flow\_indices: \\ (u,n,d,t') \in (u,ng,d,t)}} v_{unit\_flow}(u, n, d, s, t') \cdot \Delta t' \\
& >= p_{minimum\_operating\_point}(u, ng, d, s, t) \\
& \cdot p_{unit\_capacity}(u, ng, d, s, t) \\
& \cdot p_{conv\_cap\_to\_flow}(u, ng, d, s, t) \\
& \cdot \sum_{\substack{(u,s,t_{units\_on}) \in units\_on\_indices: \\ (u,\Delta t_{units\_on}) \in (u,t)}} v_{units\_on}(u, s, t_{units\_on}) \\
& \cdot \min(\Delta t_{units\_on}, \Delta t) \\
& \forall (u, ng, d) \in ind(p_{minimum\_operating\_point}), \\
& \forall t \in t\_lowest\_resolution(node\_temporal\_block(node = members(ng))), \\
& \forall s \in stochastic\_path
\end{aligned}$$

Note that this constraint is always generated for the lowest resolution of all involved members of the node group ng, i.e. the lowest resolution of the involved units flows. This is also why the term  $\min(\Delta t_{units\_on}, \Delta t)$  is added for the units on variable, in order to dis-/aggregate the units on resolution to the resolution of the unit flows.

### Minimum down time (basic version)

In order to impose a minimum offline time of a unit, before it can be started up again, the [min\\_down\\_time](#) parameter needs to be defined, which triggers the generation of the following constraint:

$$\begin{aligned}
& v_{units\_available}(u, s, t) \\
& - v_{units\_on}(u, s, t) \\
& >= \sum_{\substack{(u,s,t') \in units\_on\_indices: \\ t' >= t - p_{min\_down\_time}(u,s,t) \quad t' <= t}} v_{units\_shut\_down}(u, s, t') \\
& \forall (u, s, t) \in units\_on\_indices
\end{aligned}$$

Note that for the use reserves the generated minimum down time constraint will include [startups for non-spinning reserves](#).

### Minimum up time (basic version)

Similarly to the [minimum down time constraint](#), a minimum time that a unit needs to remain online after a startup can be imposed by defining the [min\\_up\\_time](#) parameter. This will trigger the generation of the following constraint:

$$\begin{aligned}
& v_{units\_on}(u, s, t) \\
& >= \sum_{\substack{(u,s,t') \in units\_on\_indices: \\ t' >= t - p_{min\_up\_time}(u,s,t), \\ t' <= t}} v_{units\_started\_up}(u, s, t') \\
& \forall (u, s, t) \in units\_on\_indices
\end{aligned}$$

This constraint can be extended to the use of nonspinning reserves. See [also](#).

### Ramping and reserve constraints

To include ramping and reserve constraints, it is a pre requisite that [minimum operating points](#) and [maximum capacity constraints](#) are enforced as described.

For dispatchable units, additional ramping constraints can be introduced. For setting up ramping characteristics of units see [Ramping and Reserves](#). First, the unit flows are split into their online, start-up, shut-down and non-spinning ramping contributions.

### Splitting unit flows into ramps

$$\begin{aligned}
& + \sum_{\substack{(u,n,d,s,t_{after}) \in \text{unit\_flow\_indices:} \\ (u,n,d,t_{after}) \in (u,n,d,t_{after}) \\ p_{is\_reserve}(n)}} v_{\text{unit\_flow}}(u, n, d, s, t_{after}) \\
& + \sum_{\substack{(u,n,d,s,t_{after}) \in \text{unit\_flow\_indices:} \\ (u,n,d,t_{after}) \in (u,n,d,t_{after}) \\ p_{is\_reserve}(n) \\ p_{upward\_reserve}(n)}} v_{\text{unit\_flow}}(u, n, d, s, t_{after}) \\
& - \sum_{\substack{(u,n,d,s,t_{after}) \in \text{unit\_flow\_indices:} \\ (u,n,d,t_{after}) \in (u,n,d,t_{after}) \\ p_{is\_reserve}(n) \\ p_{downward\_reserve}(n)}} v_{\text{unit\_flow}}(u, n, d, s, t_{after}) \\
& - \sum_{\substack{(u,n,d,s,t_{before}) \in \text{unit\_flow\_indices:} \\ (u,n,d,t_{before}) \in (u,n,d,t_{before}) \\ p_{is\_reserve}(n)}} v_{\text{unit\_flow}}(u, n, d, s, t_{before}) \\
& == \\
& + \sum_{\substack{(u,n,d,s,t_{after}) \in \text{ramp\_up\_unit\_flow\_indices:} \\ (u,n,d,t_{after}) \in (u,n,d,t_{after})}} v_{\text{ramp\_up\_unit\_flow}}(u, n, d, s, t_{after}) \\
& + \sum_{\substack{(u,n,d,s,t_{after}) \in \text{start\_up\_unit\_flow\_indices:} \\ (u,n,d,t_{after}) \in (u,n,d,t_{after})}} v_{\text{start\_up\_unit\_flow}}(u, n, d, s, t_{after}) \\
& + \sum_{\substack{(u,n,d,s,t_{after}) \in \text{nonspin\_ramp\_up\_unit\_flow\_indices:} \\ (u,n,d,t_{after}) \in (u,n,d,t_{after})}} v_{\text{nonspin\_ramp\_up\_unit\_flow}}(u, n, d, s, t_{after}) \\
& - \sum_{\substack{(u,n,d,s,t_{after}) \in \text{ramp\_down\_unit\_flow\_indices:} \\ (u,n,d,t_{after}) \in (u,n,d,t_{after})}} v_{\text{ramp\_down\_unit\_flow}}(u, n, d, s, t_{after}) \\
& - \sum_{\substack{(u,n,d,s,t_{after}) \in \text{shut\_down\_unit\_flow\_indices:} \\ (u,n,d,t_{after}) \in (u,n,d,t_{after})}} v_{\text{shut\_down\_unit\_flow}}(u, n, d, s, t_{after}) \\
& - \sum_{\substack{(u,n,d,s,t_{after}) \in \text{nonspin\_ramp\_down\_unit\_flow\_indices:} \\ (u,n,d,t_{after}) \in (u,n,d,t_{after})}} v_{\text{nonspin\_ramp\_down\_unit\_flow}}(u, n, d, s, t_{after}) \\
& \forall (u, n, d, s, t_{after}) \in ( \\
& \text{ramp\_up\_unit\_flow\_indices,} \\
& \text{start\_up\_unit\_flow\_indices,} \\
& \text{nonspin\_ramp\_up\_unit\_flow\_indices,} \\
& \text{ramp\_down\_unit\_flow\_indices,} \\
& \text{shut\_down\_unit\_flow\_indices,} \\
& \text{nonspin\_ramp\_down\_unit\_flow\_indices}) \\
& \forall t_{before} \in t_{before\_t}(t_{after} = t_{after}) : t_{before} \in \text{unit\_flow\_indices}
\end{aligned}$$

Note that each *individual* tuple of the `unit_flow_indices` is split into its ramping contributions, if any of the ramping variables exist for this tuple. How to set-up ramps for units is described in [Ramping and Reserves](#).

### Constraint on spinning upwards ramp\_up

The maximum online ramp up ability of a unit can be constraint by the [ramp\\_up\\_limit](#), expressed as a share of the [unit\\_capacity](#). With this constraint, online (i.e. spinning) ramps can be applied to groups of commodities (e.g. electricity + balancing capacity). Moreover, balancing product might have specific ramping requirements, which can herewith also be enforced.

$$\begin{aligned}
& + \sum_{\substack{(u,n,d,s,t) \in \text{ramp\_up\_unit\_flow\_indices:} \\ (u,n,d) \in (u,ng,d)}} v_{\text{ramp\_up\_unit\_flow}}(u, n, d, s, t) \\
& \leq \\
& + \sum_{\substack{(u,s,t') \in \text{units\_on\_indices:} \\ (u,s) \in (u,s) \\ t' \in t\_overlap\_t(t)}} (v_{\text{units\_on}}(u, s, t') - v_{\text{units\_started\_up}}(u, s, t')) \\
& \min(\Delta t', \Delta t) \\
& \cdot p_{\text{ramp\_up\_limit}}(u, ng, d, s, t) \\
& \cdot p_{\text{unit\_capacity}}(u, ng, d, s, t) \\
& \cdot p_{\text{conv\_cap\_to\_flow}}(u, ng, d, s, t) \\
& \forall (u, ng, d) \in \text{ind}(p_{\text{ramp\_up\_limit}}) \\
& \forall s \in \text{stochastic\_path}, \forall t \in \text{time\_slice}
\end{aligned}$$

Note that only online units that are not started up during this timestep are considered.

### Constraint on minimum upward start up ramp\_up

To enforce a lower bound on the ramp of a unit during start-up, the [min\\_startup\\_ramp](#) given as a share of the [unit\\_capacity](#) needs to be defined, which triggers the constraint below. Usually, only non-reserve commodities can have a start-up ramp. However, it is possible to include them, by adding them to the ramp defining node `ng`.

$$\begin{aligned}
& + \sum_{\substack{(u,n,d,s,t) \in \text{start\_up\_unit\_flow\_indices:} \\ (u,n,d) \in (u,ng,d)}} v_{\text{start\_up\_unit\_flow}}(u, n, d, s, t) \\
& \geq \\
& + \sum_{\substack{(u,s,t') \in \text{units\_on\_indices:} \\ (u,s) \in (u,s) \\ t' \in t\_overlap\_t(t)}} v_{\text{units\_started\_up}}(u, s, t') \\
& \cdot p_{\text{min\_startup\_ramp}}(u, ng, d, s, t) \\
& \cdot p_{\text{unit\_capacity}}(u, ng, d, s, t) \\
& \cdot p_{\text{conv\_cap\_to\_flow}}(u, ng, d, s, t) \\
& \forall (u, ng, d) \in \text{ind}(p_{\text{min\_startup\_ramp}}) \\
& \forall s \in \text{stochastic\_path}, \forall t \in \text{time\_slice}
\end{aligned}$$

### Constraint on maximum upward start up ramp\_up

This constraint enforces an upper limit on the unit ramp during startup process, triggered by the existence of the [max\\_startup\\_ramp](#), which should be given as a share of the [unit\\_capacity](#). Typically, only ramp flows to non-reserve nodes are considered during the start-up process. However, it is possible to include them, by adding them to the ramp defining node `ng`.

$$\begin{aligned}
& + \sum_{\substack{(u,n,d,s,t) \in \text{start\_up\_unit\_flow\_indices:} \\ (u,n,d) \in (u,ng,d)}} v_{\text{start\_up\_unit\_flow}}(u, n, d, s, t) \\
& \leq \\
& + \sum_{\substack{(u,s,t') \in \text{units\_on\_indices:} \\ (u,s) \in (u,s) \\ t' \in t\_overlap\_t(t)}} v_{\text{units\_started\_up}}(u, s, t') \\
& \cdot p_{\text{max\_startup\_ramp}}(u, ng, d, s, t) \\
& \cdot p_{\text{unit\_capacity}}(u, ng, d, s, t) \\
& \cdot p_{\text{conv\_cap\_to\_flow}}(u, ng, d, s, t) \\
& \forall (u, ng, d) \in \text{ind}(p_{\text{max\_startup\_ramp}}) \\
& \forall s \in \text{stochastic\_path}, \forall t \in \text{time\_slice}
\end{aligned}$$

### Constraint on upward non-spinning start ups

For non-spinning reserve provision, offline units can be scheduled to provide nonspinning reserves, if they have recovered their minimum down time. If nonspinning reserves are used for a unit, the minimum down-time constraint takes the following form:

$$\begin{aligned}
& v_{units\_available}(u, s, t) \\
& - v_{units\_on}(u, s, t) \\
& \geq \sum_{\substack{(u,s,t') \in units\_on\_indices: \\ t' > t - p_{min\_down\_time}(u,s,t) \\ t' \leq t}} v_{units\_shut\_down}(u, s, t') \\
& + \sum_{\substack{(u',n',s',t') \in nonspin\_units\_started\_up\_indices: \\ (u',s',t') \in (u,s,t)}} v_{nonspin\_units\_started\_up}(u', n', s', t') \\
& \forall (u, s, t) \in units\_on\_indices : \\
& (u, n, s, t) \in nonspin\_units\_started\_up\_indices
\end{aligned}$$

### Minimum nonspinning ramp up

The nonspinning ramp flows of a units [nonspin\\_ramp\\_up\\_unit\\_flow](#) are dependent on the units holding available for nonspinning reserve provision, i.e. [nonspin\\_units\\_started\\_up](#). A lower bound on these nonspinning reserves can be enforced by defining the [min\\_res\\_startup\\_ramp](#) parameter (given as a fraction of the [unit\\_capacity](#)).

$$\begin{aligned}
& + \sum_{\substack{(u,n,d,s,t) \in nonspin\_ramp\_up\_unit\_flow\_indices: \\ (u,n,d) \in (u,ng,d)}} v_{nonspin\_ramp\_up\_unit\_flow}(u, n, d, s, t) \\
& \geq \\
& + \sum_{\substack{(u,n,s,t) \in nonspin\_units\_started\_up\_indices: \\ (u,n) \in (u,ng)}} v_{nonspin\_units\_started\_up}(u, n, s, t) \\
& \cdot p_{min\_res\_startup\_ramp}(u, ng, d, s, t) \\
& \cdot p_{unit\_capacity}(u, ng, d, s, t) \\
& \cdot p_{conv\_cap\_to\_flow}(u, ng, d, s, t) \\
& \forall (u, ng, d) \in ind(p_{min\_res\_startup\_ramp}) \\
& \forall s \in stochastic\_path, \forall t \in time\_slice
\end{aligned}$$

### Maximum nonspinning ramp up

The nonspinning ramp flows of a units [nonspin\\_ramp\\_up\\_unit\\_flow](#) are dependent on the units holding available for nonspinning reserve provision, i.e. [nonspin\\_units\\_started\\_up](#). An upper bound on these nonspinning reserves can be enforced by defining the [max\\_res\\_startup\\_ramp](#) parameter (given as a fraction of the [unit\\_capacity](#)).

$$\begin{aligned}
& + \sum_{\substack{(u,n,d,s,t) \in nonspin\_ramp\_up\_unit\_flow\_indices: \\ (u,n,d) \in (u,ng,d)}} v_{nonspin\_ramp\_up\_unit\_flow}(u, n, d, s, t) \\
& \leq \\
& + \sum_{\substack{(u,n,s,t) \in nonspin\_units\_started\_up\_indices: \\ (u,n) \in (u,ng)}} v_{nonspin\_units\_started\_up}(u, n, s, t) \\
& \cdot p_{max\_res\_startup\_ramp}(u, ng, d, s, t) \\
& \cdot p_{unit\_capacity}(u, ng, d, s, t) \\
& \cdot p_{conv\_cap\_to\_flow}(u, ng, d, s, t) \\
& \forall (u, ng, d) \in ind(p_{max\_res\_startup\_ramp}) \\
& \forall s \in stochastic\_path, \forall t \in time\_slice
\end{aligned}$$

### Constraint on spinning downward ramps

Similarly to the online [ramp up capability](#) of a unit, it is also possible to impose an upper bound on the online ramp down ability of unit by defining a [ramp\\_down\\_limit](#), expressed as a share of the [unit\\_capacity](#).

$$\begin{aligned}
& + \sum_{\substack{(u,n,d,s,t) \in \text{ramp\_down\_unit\_flow\_indices:} \\ (u,n,d) \in (u,ng,d)}} v_{\text{ramp\_down\_unit\_flow}}(u, n, d, s, t) \\
& \leq \\
& + \sum_{\substack{(u,s,t') \in \text{units\_on\_indices:} \\ (u,s) \in (u,s) \\ t' \in t\_overlap\_t(t)}} (v_{\text{units\_on}}(u, s, t') - v_{\text{units\_started\_up}}(u, s, t')) \\
& \cdot p_{\text{ramp\_down\_limit}}(u, ng, d, s, t) \\
& \cdot p_{\text{unit\_capacity}}(u, ng, d, s, t) \\
& \cdot p_{\text{conv\_cap\_to\_flow}}(u, ng, d, s, t) \\
& \forall (u, ng, d) \in \text{ind}(p_{\text{ramp\_down\_limit}}) \\
& \forall s \in \text{stochastic\_path}, \forall t \in \text{time\_slice}
\end{aligned}$$

### Lower bound on downward shut-down ramps

This constraint enforces a lower bound on the unit ramp during shutdown process. Usually, units will only provide shutdown ramps to non-reserve nodes. However, it is possible to include them, by adding them to the ramp defining node `ng`. The constraint is triggered by the existence of the `min_shutdown_ramp` parameter.

$$\begin{aligned}
& + \sum_{\substack{(u,n,d,s,t) \in \text{shut\_down\_unit\_flow\_indices:} \\ (u,n,d) \in (u,ng,d)}} v_{\text{shut\_down\_unit\_flow}}(u, n, d, s, t) \\
& \leq \\
& + \sum_{\substack{(u,s,t') \in \text{units\_on\_indices:} \\ (u,s) \in (u,s) \\ t' \in t\_overlap\_t(t)}} v_{\text{units\_shut\_down}}(u, s, t') \\
& \cdot p_{\text{min\_shutdown\_ramp}}(u, ng, d, s, t) \\
& \cdot p_{\text{unit\_capacity}}(u, ng, d, s, t) \\
& \cdot p_{\text{conv\_cap\_to\_flow}}(u, ng, d, s, t) \\
& \forall (u, ng, d) \in \text{ind}(p_{\text{min\_shutdown\_ramp}}) \\
& \forall s \in \text{stochastic\_path}, \forall t \in \text{time\_slice}
\end{aligned}$$

### Upper bound on downward shut-down ramps

This constraint enforces an upper bound on the unit ramp during shutdown process. Usually, units will only provide shutdown ramps to non-reserve nodes. However, it is possible to include them, by adding them to the ramp defining node `ng`. The constraint is triggered by the existence of the `max_shutdown_ramp` parameter.

$$\begin{aligned}
& + \sum_{\substack{(u,n,d,s,t) \in \text{shut\_down\_unit\_flow\_indices:} \\ (u,n,d) \in (u,ng,d)}} v_{\text{shut\_down\_unit\_flow}}(u, n, d, s, t) \\
& \leq \\
& + \sum_{\substack{(u,s,t') \in \text{units\_on\_indices:} \\ (u,s) \in (u,s) \\ t' \in t\_overlap\_t(t)}} v_{\text{units\_shut\_down}}(u, s, t') \\
& \cdot p_{\text{max\_shutdown\_ramp}}(u, ng, d, s, t) \\
& \cdot p_{\text{unit\_capacity}}(u, ng, d, s, t) \\
& \cdot p_{\text{conv\_cap\_to\_flow}}(u, ng, d, s, t) \\
& \forall (u, ng, d) \in \text{ind}(p_{\text{max\_shutdown\_ramp}}) \\
& \forall s \in \text{stochastic\_path}, \forall t \in \text{time\_slice}
\end{aligned}$$

### Constraint on upward non-spinning shut-downs

For non-spinning downward reserves, online units can be scheduled for reserve provision through shut down if they have recovered their minimum up time. If nonspinning reserves are used the minimum up-time constraint becomes:

$$\begin{aligned}
& v_{units\_on}(u, s, t) \\
& \geq \sum_{\substack{(u,s,t') \in units\_on\_indices: \\ t' > t - p_{min\_up\_time}(u,s,t) \quad t' \leq t}} v_{units\_started\_up}(u, s, t') \\
& + \sum_{\substack{(u',n',s',t') \in nonspin\_units\_shut\_down\_indices: \\ (u',s',t') \in (u,s,t)}} v_{nonspin\_units\_shut\_down}(u', n', s', t') \\
& \forall (u, s, t) \in units\_on\_indices : \\
& u \in nonspin\_units\_started\_up\_indices
\end{aligned}$$

## Lower bound on the nonspinning downward reserve provision

A lower bound on the nonspinning reserve provision of a unit can be imposed by defining the [min\\_res\\_shutdown\\_ramp](#) parameter, which leads to the creation of the following constraint in the model:

$$\begin{aligned}
& + \sum_{\substack{(u,n,d,s,t) \in nonspin\_ramp\_down\_unit\_flow\_indices: \\ (u,n,d,s,t) \in (u,n,d,s,t)}} v_{nonspin\_ramp\_down\_unit\_flow}(u, n, d, s, t) \\
& \leq \\
& + \sum_{\substack{(u,n,s,t) \in nonspin\_units\_shut\_down\_indices: \\ (u,n,s,t) \in (u,n,s,t)}} v_{nonspin\_units\_shut\_down}(u, n, s, t) \\
& \cdot p_{min\_res\_shutdown\_ramp}(u, ng, d, s, t) \\
& \cdot p_{unit\_capacity}(u, ng, d, s, t) \\
& \cdot p_{conv\_cap\_to\_flow}(u, ng, d, s, t) \\
& \forall (u, ng, d) \in ind(p_{min\_res\_shutdown\_ramp}) \\
& \forall s \in stochastic\_path, \forall t \in time\_slice
\end{aligned}$$

## Upper bound on the nonspinning downward reserve provision

An upper limit on the nonspinning reserve provision of a unit can be imposed by defining the [max\\_res\\_shutdown\\_ramp](#) parameter, which leads to the creation of the following constraint in the model:

$$\begin{aligned}
& + \sum_{\substack{(u,n,d,s,t) \in nonspin\_ramp\_down\_unit\_flow\_indices: \\ (u,n,d,s,t) \in (u,n,d,s,t)}} v_{nonspin\_ramp\_down\_unit\_flow}(u, n, d, s, t) \\
& \leq \\
& + \sum_{\substack{(u,n,s,t) \in nonspin\_units\_shut\_down\_indices: \\ (u,n,s,t) \in (u,n,s,t)}} v_{nonspin\_units\_shut\_down}(u, n, s, t) \\
& \cdot p_{max\_res\_shutdown\_ramp}(u, ng, d, s, t) \\
& \cdot p_{unit\_capacity}(u, ng, d, s, t) \\
& \cdot p_{conv\_cap\_to\_flow}(u, ng, d, s, t) \\
& \forall (u, ng, d) \in ind(p_{max\_res\_shutdown\_ramp}) \\
& \forall s \in stochastic\_path, \forall t \in time\_slice
\end{aligned}$$

## Constraint on minimum node state for reserve provision

Storage nodes can also contribute to the provision of reserves. The amount of balancing contributions is limited by the ramps of the storage unit (see above) and by the node state:

$$\begin{aligned}
& v_{node\_state}(n_{stor}, s, t) \\
& \geq p_{node\_state\_min}(n_{stor}, s, t) \\
& + \sum_{\substack{(u,n_{res},d,s,t) \in unit\_flow\_indices: \\ u \in unit\_flow\_indices; n=n_{stor} \\ p_{is\_reserve\_node}(n_{res})}} v_{unit\_flow}(u, n_{res}, d, s, t) \\
& \cdot p_{minimum\_reserve\_activation\_time}(n_{res}) \\
& \forall (n_{stor}, s, t) \in node\_stochastic\_time\_indices : p_{has\_state}(n)
\end{aligned}$$

## Bounds on the unit capacity including ramping constraints

(Comment 2021-04-29: Currently under development)

## Operating segments

### Operating segments of units

The `unit_flow_op` operating segment variable is bounded by the difference between successive `operating_points` adjusted for `unit_capacity`

$$\begin{aligned} & v_{unit\_flow\_op}(u, n, op, s, t) \\ & \leq p_{unit\_capacity}(u, n, d, s, t) \\ & \quad \cdot v_{units\_available}(u, s, t) \\ & \quad \cdot p_{unit\_conv\_cap\_to\_flow}(u, n, d, s, t) \\ & \quad \cdot \left( p_{operating\_points}(u, n, op, s, t) \right. \\ & \quad \left. - \begin{cases} 0 & \text{if } op = 1 \\ p_{operating\_points}(u, n, op - 1, s, t) & \text{otherwise} \end{cases} \right) \\ & \quad \forall (u, n, d, s, t) \in unit\_flow\_op\_indices \end{aligned}$$

## Bounding unit flows by summing over operating segments

`unit_flow` is constrained to be the sum of all operating segment variables, `unit_flow_op`

$$\begin{aligned} & v_{unit\_flow}(u, n, s, t) \\ & = \sum_{op} v_{unit\_flow\_op}(u, n, op, s, t) \\ & \quad \forall (u, n, d) \in operating\_point\_indices \\ & \quad \forall (u, n, d, s, t) \in unit\_flow\_op\_indices \end{aligned}$$

## Unit piecewise incremental heat rate

$$\begin{aligned} & v_{unit\_flow}(u, n_{in}, d, s, t) \\ & = \sum_{op} \left( v_{unit\_flow\_op}(u, n_{out}, d, op, s, t) \right. \\ & \quad \cdot p_{unit\_incremental\_heat\_rate}(u, n_{in}, n_{out}, op, s, t) \\ & \quad + v_{units\_on}(u, s, t) \cdot p_{unit\_idle\_heat\_rate}(u, n_{in}, n_{out}, s, t) \\ & \quad \left. + v_{units\_started\_up}(u, s, t) \cdot p_{unit\_start\_flow}(u, n_{in}, n_{out}, s, t) \right) \\ & \quad \forall (u, n_{in}, n_{out}, s, t) \in unit\_pw\_heat\_rate\_indices \end{aligned}$$

## Bounds on commodity flows

### Upper bound on cumulated unit flows

To impose a limit on the cumulative amount of certain commodity flows, a cumulative bound can be set by defining the parameter `max_cum_in_unit_flow_bound` for entire optimization window:

$$\begin{aligned} & \sum_{\substack{(u, n, d, s, t') \in unit\_flow\_indices: \\ (u, n, d, t') \in (ug, ng, d)}} v_{unit\_flow}(u, n, d, s, t') \cdot \Delta t' \\ & \leq p_{max\_cum\_unit\_flow\_bound}(ug, ng, d, s, t) \\ & \quad \forall (ug, ng, d) \in ind(p_{max\_cum\_unit\_flow\_bound}) \end{aligned}$$

(Comment 2021-04-29: Currently under development)

## Network constraints



## Static constraints

### Capacity constraint on connections

In a multi-commodity setting, there can be different commodities entering/leaving a certain connection. These can be energy-related commodities (e.g., electricity, natural gas, etc.), emissions, or other commodities (e.g., water, steel). The [connection\\_capacity](#) should be specified for at least one [connection\\_to\\_node](#) or [connection\\_from\\_node](#) relationship, in order to trigger a constraint on the maximum commodity flows to this location in each time step. When desirable, the capacity can be specified for a group of nodes (e.g. combined capacity for multiple products). Note that the conversion factor [connection\\_conv\\_cap\\_to\\_flow](#) has a default value of 1, but can be adjusted in case the unit of measurement for the capacity is different to the connection flows unit of measurement.

$$\begin{aligned} & \sum_{\substack{(conn,n,d,s,t') \in \text{connection\_flow\_indices:} \\ (conn,n,d,s,t') \in (conn,ng,d,s,t)}} v_{\text{connection\_flow}}(conn, n, d, s, t') \cdot \Delta t' \\ & - \sum_{\substack{(conn,n,d_{\text{reverse}},s,t') \in \text{connection\_flow\_indices:} \\ (conn,n,s,t') \in (conn,ng,s,t) \\ d_{\text{reverse}} \neq d}} v_{\text{connection\_flow}}(conn, n, d_{\text{reverse}}, s, t') \cdot \Delta t' \\ & \leq p_{\text{connection\_capacity}}(conn, ng, d, s, t) \\ & \cdot p_{\text{connection\_availability\_factor}}(conn, s, t) \\ & \cdot p_{\text{connection\_conv\_cap\_to\_flow}}(conn, ng, d, s, t) \Delta t \\ & \forall (conn, ng, d) \in \text{ind}(p_{\text{connection\_capacity}}) : \\ & \nexists p_{\text{candidate\_connections}}(conn) \\ & \forall t \in \text{time\_slices}, \\ & \forall s \in \text{stochastic\_path} \end{aligned}$$

If the connection is a [candidate\\_connections](#), i.e. can be invested in, the connection capacity constraint translates to:

$$\begin{aligned} & \sum_{\substack{(conn,n,d,s,t') \in \text{connection\_flow\_indices:} \\ (conn,n,d,s,t') \in (conn,ng,d,s,t)}} v_{\text{connection\_flow}}(conn, n, d, s, t') \cdot \Delta t' \\ & - \sum_{\substack{(conn,n,d_{\text{reverse}},s,t') \in \text{connection\_flow\_indices:} \\ (conn,n,s,t') \in (conn,ng,s,t) \\ d_{\text{reverse}} \neq d}} v_{\text{connection\_flow}}(conn, n, d_{\text{reverse}}, s, t') \cdot \Delta t' \\ & \leq p_{\text{connection\_capacity}}(conn, ng, d, s, t) \\ & \cdot p_{\text{connection\_availability\_factor}}(conn, s, t) \\ & \cdot p_{\text{connection\_conv\_cap\_to\_flow}}(conn, ng, d, s, t) \Delta t \\ & \cdot \sum_{\substack{(conn,s,t') \in \text{connections\_invested\_available\_indices:} \\ (conn,s,t') \in (conn,s,t\_in\_t(t_{\text{short}})}} v_{\text{connections\_invest\_available}}(conn, s, t) \quad \forall (conn, ng, d) \in \text{ind}(p_{\text{connection\_capacity}}) : \\ & \exists p_{\text{candidate\_connections}}(conn) \\ & \forall t \in \text{time\_slices}, \\ & \forall s \in \text{stochastic\_path} \end{aligned}$$

### Fixed ratio between outgoing and incoming flows of a connection

By defining the parameters [fix\\_ratio\\_out\\_in\\_connection\\_flow](#), [max\\_ratio\\_out\\_in\\_connection\\_flow](#) or [min\\_ratio\\_out\\_in\\_connection\\_flow](#), a ratio can be set between **outgoing** and **incoming** flows from and to a connection.

In the most general form of the equation, two node groups are defined (an input node group  $ng_{in}$  and an output node group  $ng_{out}$ ), and a linear relationship is expressed between both node groups. Note that whenever the relationship is specified between groups of multiple nodes, there remains a degree of freedom regarding the composition of the input node flows within group  $ng_{in}$  and the output node flows within group  $ng_{out}$ .

The constraint given below enforces a fixed, maximum or minimum ratio between outgoing and incoming [connection\\_flow](#). Note that the potential node groups, that the parameters [fix\\_ratio\\_out\\_in\\_connection\\_flow](#), [max\\_ratio\\_out\\_in\\_connection\\_flow](#) and [min\\_ratio\\_out\\_in\\_connection\\_flow](#) are defined on, are getting internally expanded to the members of the node group within the [connection\\_flow\\_indices](#).

$$\begin{aligned}
& \sum_{\substack{(conn,n,d,s,t_{out}) \in connection\_flow\_indices: \\ (conn,n,d,s,t_{out}) \in (conn,ng_{out},:to\_node,s,t)}} v_{connection\_flow}(conn,n,d,s,t_{out}) \cdot \Delta t_{out} \\
& \{ \\
& \quad == p_{fix\_ratio\_out\_in\_connection\_flow}(conn,ng_{out},ng_{in},s,t), \\
& \quad <= p_{max\_ratio\_out\_in\_connection\_flow}(conn,ng_{out},ng_{in},s,t), \\
& \quad >= p_{min\_ratio\_out\_in\_connection\_flow}(conn,ng_{out},ng_{in},s,t) \\
& \} \\
& \cdot \sum_{\substack{(conn,n,d,s,t_{in}) \in connection\_flow\_indices: \\ (conn,n,d,s,t_{in}) \in (conn,ng_{in},:from\_node,s,t)}} v_{connection\_flow}(conn,n,d,s,t_{in}) \cdot \Delta t_{in} \\
& \forall (conn,ng_{out},ng_{in}) \in ind(p_{\{fix,max,min\}\_ratio\_out\_in\_connection\_flow}), \\
& \forall t \in time\_slices, \forall s \in stochastic\_path
\end{aligned}$$

## Specific network representation

In the following, the different specific network representations are introduced. While the [Static constraints](#) find application in any of the different networks, the following equations are specific to the discussed use cases. Currently, SpineOpt incorporated equations for pressure driven gas networks, nodal lossless DC power flows and PTDF based lossless DC power flow.

### Pressure driven gas transfer

For gas pipelines it can be relevant a pressure driven gas transfer can be modelled, i.e. to account for linepack flexibility. Generally speaking, the main challenges related to pressure driven gas transfers are the non-convexities associated with the Weymouth equation. In SpineOpt, a convexified MILP representation has been implemented, which as been presented in [Schwele - Coordination of Power and Natural Gas Systems: Convexification Approaches for Linepack Modeling](#). The approximation approach is based on the Taylor series expansion around fixed pressure points.

In addition to the already known variables, such as [connection\\_flow](#) and [node\\_state](#), the start and end points of a gas pipeline connection are associated with the variable [node\\_pressure](#). The variable is triggered by the [has\\_pressure](#) parameter. For more details on how to set up a gas pipeline, see also the advanced concept section [on pressure driven gas transfer](#).

### Maximum node pressure

In order to impose an upper limit on the maximum pressure at a node the [maximum node pressure constraint](#) can be included, by defining the parameter [max\\_node\\_pressure](#) which triggers the following constraint:

$$\begin{aligned}
& \sum_{\substack{(n,s,t') \in node\_pressure\_indices: \\ (n,s,t') \in (n,s,t)}} v_{node\_pressure}(n,s,t') \cdot \Delta t' \\
& <= p_{max\_node\_pressure}(ng,s,t) \cdot \Delta t \\
& \forall (ng) \in ind(p_{max\_node\_pressure}), \\
& \forall t \in time\_slices, \\
& \forall s \in stochastic\_path
\end{aligned}$$

As indicated in the equation, the parameter [max\\_node\\_pressure](#) can also be defined on a node group, in order to impose an upper limit on the aggregated [node\\_pressure](#) within one node group.

### Minimum node pressure

In order to impose a lower limit on the pressure at a node the [maximum node pressure constraint](#) can be included, by defining the parameter [min\\_node\\_pressure](#) which triggers the following constraint:

$$\begin{aligned}
& \sum_{\substack{(n,s,t') \in node\_pressure\_indices: \\ (n,s,t') \in (n,s,t)}} v_{node\_pressure}(n,s,t') \cdot \Delta t' \\
& >= p_{min\_node\_pressure}(ng,s,t) \cdot \Delta t \\
& \forall (ng) \in ind(p_{min\_node\_pressure}), \\
& \forall t \in time\_slices, \\
& \forall s \in stochastic\_path
\end{aligned}$$

As indicated in the equation, the parameter `min_node_pressure` can also be defined on a node group, in order to impose a lower limit on the aggregated `node_pressure` within one node group.

### Constraint on the pressure ratio between two nodes

If a compression station is located in between two nodes, the connection is considered to be active and a compression ratio between the two nodes can be imposed. The parameter `compression_factor` needs to be defined on a `connection_node_node` relationship, where the first node corresponds the origin node, before the compression, while the second node corresponds to the destination node, after compression. The existence of this parameter will trigger the following constraint:

$$\begin{aligned}
& \sum_{\substack{(n,s,t') \in \text{node\_pressure\_indices:} \\ (n,s,t') \in (ng2,s,t)}} v_{\text{node\_pressure}}(n,s,t') \cdot \Delta t' \\
& \leq p_{\text{compression\_factor}}(\text{conn}, ng1, ng2, s, t) \\
& \sum_{\substack{(n,s,t') \in \text{node\_pressure\_indices:} \\ (n,s,t') \in (ng1,s,t)}} v_{\text{node\_pressure}}(n,s,t') \cdot \Delta t' \\
& \forall (\text{conn}, ng1, ng2) \in \text{ind}(p_{\text{compression\_factor}}), \\
& \forall t \in \text{time\_slices}, \\
& \forall s \in \text{stochastic\_path}
\end{aligned}$$

### Outer approximation through fixed pressure points

The Weymouth relates the average flows through a connection to the difference between the adjacent squared node pressures.

$$\begin{aligned}
& ((v_{\text{connection\_flow}}(\text{conn}, n_{\text{orig}}, : \text{from\_node}, s, t) + v_{\text{connection\_flow}}(\text{conn}, n_{\text{dest}}, : \text{to\_node}, s, t))/2 \\
& - (v_{\text{connection\_flow}}(\text{conn}, n_{\text{dest}}, : \text{from\_node}, s, t) + v_{\text{connection\_flow}}(\text{conn}, n_{\text{orig}}, : \text{to\_node}, s, t))/2) \\
& \cdot \\
& |((v_{\text{connection\_flow}}(\text{conn}, n_{\text{orig}}, : \text{from\_node}, s, t) + v_{\text{connection\_flow}}(\text{conn}, n_{\text{dest}}, : \text{to\_node}, s, t))/2 \\
& - (v_{\text{connection\_flow}}(\text{conn}, n_{\text{dest}}, : \text{from\_node}, s, t) + v_{\text{connection\_flow}}(\text{conn}, n_{\text{orig}}, : \text{to\_node}, s, t))/2)| \\
& = K(\text{conn}) \cdot (v_{\text{node\_pressure}}(n_{\text{orig}}, s, t)^2 - v_{\text{node\_pressure}}(n_{\text{dest}}, s, t)^2)
\end{aligned}$$

Which can be rewritten as

$$\begin{aligned}
& ((v_{\text{connection\_flow}}(\text{conn}, n_{\text{orig}}, : \text{from\_node}, s, t) + v_{\text{connection\_flow}}(\text{conn}, n_{\text{dest}}, : \text{to\_node}, s, t))/2 \\
& - (v_{\text{connection\_flow}}(\text{conn}, n_{\text{dest}}, : \text{from\_node}, s, t) + v_{\text{connection\_flow}}(\text{conn}, n_{\text{orig}}, : \text{to\_node}, s, t))/2) \\
& = \sqrt{K(\text{conn}) \cdot (v_{\text{node\_pressure}}(n_{\text{orig}}, s, t)^2 - v_{\text{node\_pressure}}(n_{\text{dest}}, s, t)^2)} \\
& \forall (v_{\text{connection\_flow}}(\text{conn}, n_{\text{orig}}, : \text{from\_node}, s, t) + v_{\text{connection\_flow}}(\text{conn}, n_{\text{dest}}, : \text{to\_node}, s, t))/2 > 0
\end{aligned}$$

and

$$\begin{aligned}
& ((v_{\text{connection\_flow}}(\text{conn}, n_{\text{dest}}, : \text{from\_node}, s, t) + v_{\text{connection\_flow}}(\text{conn}, n_{\text{orig}}, : \text{to\_node}, s, t))/2 \\
& - (v_{\text{connection\_flow}}(\text{conn}, n_{\text{orig}}, : \text{from\_node}, s, t) + v_{\text{connection\_flow}}(\text{conn}, n_{\text{dest}}, : \text{to\_node}, s, t))/2) \\
& = \sqrt{K(\text{conn}) \cdot (v_{\text{node\_pressure}}(n_{\text{dest}}, s, t)^2 - v_{\text{node\_pressure}}(n_{\text{orig}}, s, t)^2)} \\
& \forall (v_{\text{connection\_flow}}(\text{conn}, n_{\text{orig}}, : \text{from\_node}, s, t) + v_{\text{connection\_flow}}(\text{conn}, n_{\text{dest}}, : \text{to\_node}, s, t))/2 < 0
\end{aligned}$$

where `K` corresponds to the natural gas flow constant.

The cone described by the Weymouth equation can be outer approximated by a number of tangent planes, using a set of fixed pressure points, as illustrated in [Schwele - Integration of Electricity, Natural Gas and Heat Systems With Market-based Coordination](#). The bigM method is used to replace the sign function.

The linearized version of the Weymouth equation implemented in SpineOpt is given as follows:

$$\begin{aligned}
& ((v_{\text{connection\_flow}}(\text{conn}, n_{\text{orig}}, : \text{from\_node}, s, t) + v_{\text{connection\_flow}}(\text{conn}, n_{\text{dest}}, : \text{to\_node}, s, t)) / 2 \\
& \leq p_{\text{fixed\_pressure\_constant\_1}}(\text{conn}, n_{\text{orig}}, n_{\text{dest}}, j, s, t) \cdot v_{\text{node\_pressure}}(n_{\text{orig}}, s, t) \\
& - p_{\text{fixed\_pressure\_constant\_0}}(\text{conn}, n_{\text{orig}}, n_{\text{dest}}, j, s, t) \cdot v_{\text{node\_pressure}}(n_{\text{dest}}, s, t) \\
& + p_{\text{big\_m}} \cdot (1 - v_{\text{binary\_gas\_connection\_flow}}(\text{conn}, n_{\text{dest}}, : \text{to\_node}, s, t)) \\
& \forall (\text{conn}, n_{\text{orig}}, n_{\text{dest}}) \in \text{ind}(p_{\text{fixed\_pressure\_constant\_1}}) \\
& \forall j \in 1 : n(p_{\text{fixed\_pressure\_constant\_1}}(\text{connection}=\text{conn}, \text{node1}=n_{\text{orig}}, \text{node2}=n_{\text{dest}})) : \\
& p_{\text{fixed\_pressure\_constant\_1}}(\text{conn}, n_{\text{orig}}, n_{\text{dest}}, i = j) = 0 \\
& \forall t \in \text{time\_slices}, \\
& \forall s \in \text{stochastic\_path}
\end{aligned}$$

The parameters [fixed\\_pressure\\_constant\\_1](#) and [fixed\\_pressure\\_constant\\_0](#) should be defined in the database. For each considered fixed pressure point, they can be calculated as follows:

$$\begin{aligned}
& \begin{aligned} & p_{\text{fixed\_pressure\_constant\_1}}(\text{conn}, n_{\text{orig}}, n_{\text{dest}}, j) \cdot K(\text{conn}) \cdot p_{\text{fixed\_pressure}}(n_{\text{orig}}, j) / \\ & \sqrt{p_{\text{fixed\_pressure}}(n_{\text{orig}}, j)^2 - p_{\text{fixed\_pressure}}(n_{\text{dest}}, j)^2} \end{aligned} \\
& \begin{aligned} & p_{\text{fixed\_pressure\_constant\_0}}(\text{conn}, n_{\text{orig}}, n_{\text{dest}}, j) \cdot K(\text{conn}) \cdot p_{\text{fixed\_pressure}}(n_{\text{dest}}, j) / \\ & \sqrt{p_{\text{fixed\_pressure}}(n_{\text{orig}}, j)^2 - p_{\text{fixed\_pressure}}(n_{\text{dest}}, j)^2} \end{aligned} \\
& \text{where } K \text{ corresponds to the natural gas flow constant.}
\end{aligned}$$

The [big\\_m](#) parameter combined with the variable [binary\\_gas\\_connection\\_flow](#) together with the equations [on unitary gas flow](#) and on the [maximum gas flow](#) ensure that the bound on the average flow through the fixed pressure points becomes active, if the flow is in a positive direction for the observed set of connection, node1 and node2.

### Enforcing unidirectional flow

As stated above, the flow through a connection can only be in one direction at a time. Whether a flow is active in a certain direction is indicated by the [binary\\_gas\\_connection\\_flow](#) variable, which takes a value of 1 if the direction of flow is positive. To ensure that the [binary\\_gas\\_connection\\_flow](#) in the opposite direction then takes the value 0, the following constraint is enforced:

$$\begin{aligned}
& v_{\text{binary\_gas\_connection\_flow}}(\text{conn}, n_{\text{orig}}, : \text{to\_node}, s, t) \\
& = (1 - v_{\text{binary\_gas\_connection\_flow}}(\text{conn}, n_{\text{dest}}, : \text{to\_node}, s, t)) \\
& \forall (n, d, s, t) \in \text{binary\_gas\_connection\_flow\_indices}
\end{aligned}$$

### Gas connection flow capacity

To enforce that the average flow of a connection is only in one direction, the flow in the opposite direction is forced to be 0 by the following equation. For the connection flow in the direction of flow the parameter [big\\_m](#) should be chosen large enough to not become binding.

$$\begin{aligned}
& ((v_{\text{connection\_flow}}(\text{conn}, n_{\text{orig}}, : \text{from\_node}, s, t) + v_{\text{connection\_flow}}(\text{conn}, n_{\text{dest}}, : \text{to\_node}, s, t)) / 2 \\
& \leq p_{\text{big\_m}} \cdot v_{\text{binary\_gas\_connection\_flow}}(\text{conn}, n_{\text{dest}}, : \text{to\_node}, s, t) \\
& \forall (\text{conn}, n_{\text{orig}}, n_{\text{dest}}) \in \text{ind}(p_{\text{fixed\_pressure\_constant\_1}}) \\
& \forall t \in \text{time\_slices}, \\
& \forall s \in \text{stochastic\_path}
\end{aligned}$$

### Linepack storage flexibility

In order to account for linepack flexibility, i.e. storage capability of a connection, the linepack storage is linked to the average pressure of the adjacent nodes by the following equation, triggered by the parameter [connection\\_linepack\\_constant](#):

$$\begin{aligned}
& v_{\text{node\_state}}(n_{\text{stor}}, s, t) \Delta t \\
& = p_{\text{connection\_linepack\_constant}}(\text{conn}, n_{\text{stor}}, n_{\text{ngroup}}) / 2 \sum_{\substack{(n, s, t') \in \text{node\_pressure\_indices:} \\ (n, s, t') \in (n_{\text{g}}, s, t)}} v_{\text{node\_pressure}}(n, s, t') \cdot \Delta t' \\
& \forall (\text{conn}, n_{\text{stor}}, n_{\text{ngroup}}) \in \text{ind}(p_{\text{connection\_linepack\_constant}}) \\
& \forall t \in \text{time\_slices}, \\
& \forall s \in \text{stochastic\_path}
\end{aligned}$$

Note that the parameter [connection\\_linepack\\_constant](#) should be defined on a [connection\\_\\_node\\_\\_node](#) relationship, where the first node corresponds to the linepack storage node, whereas the second node corresponds to the node group of both start and end nodes of the pipeline.

### Nodebased lossless DC power flow

For the implementation of the nodebased loss DC powerflow model, a new variable [node\\_voltage\\_angle](#) is introduced. See also [has\\_voltage\\_angle](#). For further explanation on setting up a database for nodal lossless DC power flow, see the advanced concept chapter on [Lossless nodal DC power flows](#).

## Maximum node voltage angle

In order to impose an upper limit on the maximum voltage angle at a node the [maximum node voltage angle constraint](#) can be included, by defining the parameter [max\\_voltage\\_angle](#) which triggers the following constraint:

$$\sum_{\substack{(n,s,t') \in \text{node\_voltage\_angle\_indices:} \\ (n,s,t') \in (n,s,t)}} v_{\text{node\_voltage\_angle}}(n, s, t') \cdot \Delta t' \\ \leq p_{\text{max\_voltage\_angle}}(ng, s, t) \cdot \Delta t \\ \forall (ng) \in \text{ind}(p_{\text{max\_voltage\_angle}}), \\ \forall t \in \text{time\_slices}, \\ \forall s \in \text{stochastic\_path}$$

As indicated in the equation, the parameter [max\\_voltage\\_angle](#) can also be defined on a node group, in order to impose an upper limit on the aggregated [node\\_voltage\\_angle](#) within one node group.

## Minimum node voltage angle

In order to impose a lower limit on the voltage angle at a node the [maximum node voltage angle constraint](#) can be included, by defining the parameter [min\\_voltage\\_angle](#) which triggers the following constraint:

$$\sum_{\substack{(n,s,t') \in \text{node\_voltage\_angle\_indices:} \\ (n,s,t') \in (n,s,t)}} v_{\text{node\_voltage\_angle}}(n, s, t') \cdot \Delta t' \\ \geq p_{\text{min\_voltage\_angle}}(ng, s, t) \cdot \Delta t \\ \forall (ng) \in \text{ind}(p_{\text{min\_voltage\_angle}}), \\ \forall t \in \text{time\_slices}, \\ \forall s \in \text{stochastic\_path}$$

As indicated in the equation, the parameter [min\\_voltage\\_angle](#) can also be defined on a node group, in order to impose a lower limit on the aggregated [node\\_voltage\\_angle](#) within one node group.

## Voltage angle to connection flows

To link the flow over a connection to the voltage angles of the adjacent nodes, the following constraint is imposed. Note that this constraint is only generated if the parameter [connection\\_reactance](#) is defined for a [connection\\_\\_node\\_\\_node](#) relationship and if a [fix\\_ratio\\_out\\_in\\_connection\\_flow](#) is defined for the corresponding connection, node, node tuples.

$$\begin{aligned}
& + \sum_{\substack{(conn, n', d, s, t) \in \text{connection\_flow\_indices:} \\ d_{from} == \text{from\_node} \\ n' \in n_{from}}} v_{\text{connection\_flow}}(conn, n', d, s, t) \\
& - \sum_{\substack{(conn, n', d, s, t) \in \text{connection\_flow\_indices:} \\ d_{from} == \text{from\_node} \\ n' \in n_{to}}} v_{\text{connection\_flow}}(conn, n', s, t) \\
& = \\
& 1/p_{\text{connection\_reactance}}(conn) \cdot p_{\text{connection\_reactance\_base}}(conn) \\
& \cdot \left( \sum_{\substack{(n, s, t') \in \text{node\_voltage\_angle\_indices:} \\ (n, s, t') \in (n_{from}, s, t)}} v_{\text{node\_voltage\_angle}}(n, s, t') \cdot \Delta t' \right. \\
& \left. - \sum_{\substack{(n, s, t') \in \text{node\_voltage\_angle\_indices:} \\ (n, s, t') \in (n_{to}, s, t)}} v_{\text{node\_voltage\_angle}}(n, s, t') \cdot \Delta t' \right) \\
& (conn, n_{to}, n_{from}) \in \text{indices}(p_{\text{fix\_ratio}}, ut_i, n_{\text{connection\_flow}}) \\
& \forall t \in \text{time\_slices}, \\
& \forall s \in \text{stochastic\_path}
\end{aligned}$$

## PTDF based DC lossless powerflow

### Connection intact flow PTDF

The power transfer distribution factors are a property of the network reactances. `ptdf(n, c)` represents the fraction of an injection at `node n` that will flow on `connection c`. The flow on `connection c` is then the sum over all nodes of `ptdf(n, c) * net_injection(c)`. `connection_intact_flow` represents the flow on each line of the network will all candidate connections with PTDF-based flow present in the network.

$$\begin{aligned}
& + v_{\text{connection\_intact\_flow}}(c, n_{to}, d_{to}, s, t) \\
& - v_{\text{connection\_intact\_flow}}(c, n_{to}, d_{from}, s, t) \\
& == \sum_{n_{inj}} \left( v_{\text{node\_injection}}(n_{inj}, s, t) \cdot p_{\text{ptdf}}(c, n_{inj}) \right) \\
& \forall (c, n_{to}, s, t) \in \text{connection\_ptdf\_flow\_indices}
\end{aligned}$$

### N-1 post contingency connection flow limits

The N-1 security constraint for the post-contingency flow on monitored connection, `c_mon`, upon the outage of contingency connection, `c_conn`, is formed using line outage distribution factors (`lodf`). `lodf(c_conn, c_mon)` represents the fraction of the pre-contingency flow on connection `c_conn` that will flow on `c_mon` if `c_conn` is disconnected. If `connection c_conn` is disconnected, the post-contingency flow on monitored connection `connection c_mon` is the pre-contingency `connection_flow` on `c_mon` plus the line outage distribution factor (`lodf`) times the pre-contingency `connection_flow` on `c_conn`. This post-contingency flow should be less than the `connection_emergency_capacity` of `c_mon`.

$$\begin{aligned}
& + v_{\text{connection\_flow}}(c_{mon}, n_{mon\_to}, d_{to}, s, t) \\
& - v_{\text{connection\_flow}}(c_{mon}, n_{mon\_to}, d_{from}, s, t) \\
& + p_{\text{lodf}}(c_{conn}, c_{mon}) \cdot \left( \begin{aligned} & + v_{\text{connection\_flow}}(c_{conn}, n_{conn\_to}, d_{to}, s, t) \\ & - v_{\text{connection\_flow}}(c_{conn}, n_{conn\_to}, d_{from}, s, t) \end{aligned} \right) \\
& < \min(p_{\text{connection\_emergency\_capacity}}(c_{mon}, n_{conn\_to}, d_{to}, s, t), p_{\text{connection\_emergency\_capacity}}(c_{mon}, n_{conn\_to}, d_{from}, s, t)) \\
& \forall (c_{mon}, c_{conn}, s, t) \in \text{constraint\_connection\_flow\_lodf\_indices}
\end{aligned}$$

## Investments

### Investments in units

#### Economic lifetime of a unit

Enforces the minimum duration of a unit's investment decision. Once a unit has been invested-in, it must remain invested-in for `unit_investment_lifetime`.

$$\begin{aligned}
 &v_{units\_invested\_available}(u, s, t) \\
 &>= \sum_{\substack{(u,s,t') \in units\_invested\_available\_indices: \\ t' >= t - p_{unit\_investment\_lifetime}(u,s,t), \\ t' <= t}} v_{units\_invested}(u, s, t') \\
 &\forall (u, s, t) \in unit\_investment\_lifetime\_indices
 \end{aligned}$$

## Technical lifetime of a unit

(Comment 2021-04-29: Currently under development)

## Available Investment Units

The number of available invested-in units at any point in time is less than the number of investment candidate units.

$$\begin{aligned}
 &v_{units\_invested\_available}(u, s, t) \\
 &< p_{candidate\_units}(u, s, t) \\
 &\forall u \in candidate\_units\_indices, \\
 &\forall (u, s, t) \in units\_invested\_available\_indices
 \end{aligned}$$

## Investment transfer

`units_invested` represents the point-in-time decision to invest in a unit or not while `units_invested_available` represents the invested-in units that are available in a specific timeslice. This constraint enforces the relationship between `units_invested`, `units_invested_available` and `units_mothballed` in adjacent timeslices.

$$\begin{aligned}
 &v_{units\_invested\_available}(u, s, t_{after}) \\
 &- v_{units\_invested}(u, s, t_{after}) \\
 &+ v_{units\_mothballed}(u, s, t_{after}) \\
 &== v_{units\_invested\_available}(u, s, t_{before}) \\
 &\forall (u, s, t_{after}) \in units\_invested\_available\_indices, \\
 &\forall t_{before} \in t\_before\_t(t\_after = t_{after}) : t_{before} \in units\_invested\_available\_indices
 \end{aligned}$$

## Investments in connections

### Available invested-in connections

The number of available invested-in connections at any point in time is less than the number of investment candidate connections.

$$\begin{aligned}
 &v_{connections\_invested\_available}(c, s, t) \\
 &< p_{candidate\_connections}(c, s, t) \\
 &\forall c \in candidate\_connections\_indices, \\
 &\forall (c, s, t) \in connections\_invested\_available\_indices
 \end{aligned}$$

## Transfer of previous investments

`connections_invested` represents the point-in-time decision to invest in a connection or not while `connections_invested_available` represents the invested-in connections that are available in a specific timeslice. This constraint enforces the relationship between `connections_invested`, `connections_invested_available` and `connections_decommissioned` in adjacent timeslices.

$$\begin{aligned}
& v_{connections\_invested\_available}(c, s, t_{after}) \\
& - v_{connections\_invested}(c, s, t_{after}) \\
& + v_{connections\_decommissioned}(c, s, t_{after}) \\
& == v_{connections\_invested\_available}(c, s, t_{before}) \\
& \forall (c, s, t_{after}) \in connections\_invested\_available\_indices, \\
& \forall t_{before} \in t\_before\_t(t_{after} = t_{after}) : t_{before} \in connections\_invested\_available\_indices
\end{aligned}$$

## Intact network ptdf-based flows on connections

Enforces the relationship between [connection\\_intact\\_flow](#) (flow with all investments assumed in force) and [connection\\_flow](#). [connection\\_intact\\_flow](#) is the flow on all lines with all investments assumed in place. This constraint ensures that the [connection\\_flow](#) is [connection\\_intact\\_flow](#) plus additional flow contributions from investment connections that are not invested in.

$$\begin{aligned}
& + v_{connection\_flow}(c, n_{to}, d_{from}, s, t) \\
& - v_{connection\_flow}(c, n_{to}, d_{to}, s, t) \\
& - v_{connection\_intact\_flow}(c, n_{to}, d_{from}, s, t) \\
& + v_{connection\_intact\_flow}(c, n_{to}, d_{to}, s, t) \\
& == \\
& \sum_{c_{candidate}, n_{to\_candidate}} p_{lopf}(c_{candidate}, c) \cdot \left( \begin{aligned} & + v_{connection\_flow}(c_{candidate}, n_{to\_candidate}, d_{from}, s, t) \\ & - v_{connection\_flow}(c_{candidate}, n_{to\_candidate}, d_{to}, s, t) \\ & - v_{connection\_intact\_flow}(c_{candidate}, n_{to\_candidate}, d_{from}, s, t) \\ & + v_{connection\_intact\_flow}(c_{candidate}, n_{to\_candidate}, d_{to}, s, t) \end{aligned} \right) \\
& \forall (c, n_{to}, s, t) \in connection\_flow\_intact\_flow\_indices
\end{aligned}$$

## Intact connection flows capacity

Similarly to [constraint\\_connection\\_flow\\_capacity](#), limits [connection\\_intact\\_flow](#) according to [connection\\_capacity](#)

$$\begin{aligned}
& \sum_{\substack{(conn, n, d, s, t') \in connection\_intact\_flow\_indices: \\ (conn, n, d, s, t') \in (conn, ng, d, s, t)}} v_{connection\_intact\_flow}(conn, n, d, s, t') \cdot \Delta t' \\
& - \sum_{\substack{(conn, n, d_{reverse}, s, t') \in connection\_intact\_flow\_indices: \\ (conn, n, s, t') \in (conn, ng, s, t) \\ d_{reverse} \neq d}} v_{connection\_intact\_flow}(conn, n, d_{reverse}, s, t') \cdot \Delta t' \\
& \leq p_{connection\_capacity}(conn, ng, d, s, t) \\
& \cdot p_{connection\_availability\_factor}(conn, s, t) \\
& \cdot p_{connection\_conv\_cap\_to\_flow}(conn, ng, d, s, t) \Delta t \\
& \forall (conn, ng, d) \in ind(p_{connection\_capacity}) : \\
& \forall t \in time\_slices, \\
& \forall s \in stochastic\_path
\end{aligned}$$

## Fixed ratio between outgoing and incoming intact flows of a connection

For ptdf-based lossless DC power flow, ensures that the output flow to the `to_node` equals the input flow from the `from_node`.

$$\begin{aligned}
& + v_{connection\_intact\_flow}(c, n_{out}, d_{to}, s, t) \\
& == \\
& + v_{connection\_intact\_flow}(c, n_{in}, d_{from}, s, t) \\
& \forall (c, n_{in}, n_{out}, s, t) \in connection\_intact\_flow\_indices
\end{aligned}$$

## Lower bound on candidate connection flow

For candidate connections with PTDf-based powerflow, together with [constraint\\_candidate\\_connection\\_flow\\_ub](#), this constraint ensures that [connection\\_flow](#) is zero if the candidate connection is not invested-in and equals [connection\\_intact\\_flow](#) otherwise.



$$\begin{aligned}
& + v_{\text{connection\_flow}}(c, n, d, s, t) \\
& \geq \\
& + v_{\text{connection\_intact\_flow}}(c, n, d, s, t) \\
& - p_{\text{connection\_capacity}}(c, n, d, s, t) \cdot (p_{\text{candidate\_connections}}(c, s, t) - v_{\text{connections\_invested\_available}}(c, s, t)) \\
& \forall (c, n, d, s, t) \in \text{constraint\_candidate\_connection\_flow\_lb\_indices}
\end{aligned}$$

## Upper bound on candidate connection flow

For candidate connections with PTDF-based poweflow, together with [constraint\\_candidate\\_connection\\_flow\\_lb](#), this constraint ensures that [connection\\_flow](#) is zero if the candidate connection is not invested-in and equals [connection\\_intact\\_flow](#) otherwise.

$$\begin{aligned}
& + v_{\text{connection\_flow}}(c, n, d, s, t) \\
& \leq \\
& + v_{\text{connection\_intact\_flow}}(c, n, d, s, t) \\
& \forall (c, n, d, s, t) \in \text{constraint\_candidate\_connection\_flow\_ub\_indices}
\end{aligned}$$

## Economic lifetime of a connection

Enforces the minimum duration of a [connection](#)'s investment decision. Once a [connection](#) has been invested-in, it must remain invested-in for [connection\\_investment\\_lifetime](#).

$$\begin{aligned}
& v_{\text{connections\_invested\_available}}(c, s, t) \\
& \geq \sum_{\substack{(c, s, t') \in \text{connections\_invested\_available\_indices:} \\ t' \geq t - p_{\text{connection\_investment\_lifetime}}(c, s, t), \\ t' \leq t}} v_{\text{connections\_invested}}(c, s, t') \\
& \forall (c, s, t) \in \text{connection\_investment\_lifetime\_indices}
\end{aligned}$$

## Technical lifetime of a connection

(Comment 2021-04-29: Currently under development)

## Investments in storages

Note: can we actually invest in nodes that are not storages? (e.g. new location)

## Available invested storages

The number of available invested-in storages at node n at any point in time is less than the number of investment candidate storages at that node.

$$\begin{aligned}
& v_{\text{storages\_invested\_available}}(n, s, t) \\
& < p_{\text{candidate\_storages}}(n, s, t) \\
& \forall (n) \in \text{candidate\_storages\_indices}, \\
& \forall (n, s, t) \in \text{storages\_invested\_available\_indices}
\end{aligned}$$

## Storage capacity transfer?

[storages\\_invested](#) represents the point-in-time decision to invest in storage at a node, n or not while [storages\\_invested\\_available](#) represents the invested-in storages that are available at a node in a specific timeslice. This constraint enforces the relationship between [storages\\_invested](#), [storages\\_invested\\_available](#) and [storages\\_decommissioned](#) in adjacent timeslices.

$$\begin{aligned}
& v_{\text{storages\_invested\_available}}(n, s, t_{\text{after}}) \\
& - v_{\text{storages\_invested}}(n, s, t_{\text{after}}) \\
& + v_{\text{storages\_decommissioned}}(n, s, t_{\text{after}}) \\
& == v_{\text{storages\_invested\_available}}(n, s, t_{\text{before}}) \\
& \forall (n, s, t_{\text{after}}) \in \text{storages\_invested\_available\_indices}, \\
& \forall t_{\text{before}} \in t_{\text{before\_t}}(t_{\text{after}} = t_{\text{after}}) : t_{\text{before}} \in \text{storages\_invested\_available\_indices}
\end{aligned}$$

## Economic lifetime of a storage

Enforces the minimum duration of a storage investment decision at node  $n$ . Once a storage has been invested-in, it must remain invested-in for `storage_investment_lifetime`.

$$\begin{aligned}
 &v_{storages\_invested\_available}(n, s, t) \\
 &\geq \sum_{\substack{(n, s, t') \in storages\_invested\_available\_indices: \\ t' \geq t - p_{storage\_investment\_lifetime}(n, s, t), \\ t' \leq t}} v_{storages\_invested}(n, s, t') \\
 &\forall (n, s, t) \in storage\_investment\_lifetime\_indices
 \end{aligned}$$

## Technical lifetime of a storage

(Comment 2021-04-29: Currently under development)

## Capacity transfer

(Comment 2021-04-29: Currently under development)

## Early retirement of capacity

(Comment 2021-04-29: Currently under development)

## Benders decomposition

This section describes the high-level formulation of the benders-decomposed problem.

Taking the simple example of minimising capacity and operating cost for a fleet of units with a linear cost coefficient ' $operational\_cost_u$ ' :

$$\begin{aligned}
 &minimise : \\
 &+ \sum_u p_{unit\_investment\_cost}(u) \cdot v_{units\_invested}(u) \\
 &+ \sum_{u, n, t} p_{operational\_cost} \cdot v_{unit\_flow}(u, n, t) \\
 &subject\ to : \\
 &flow_{u, n, t} \leq p_{unit\_capacity}(u, n, t) \cdot (v_{units\_available} + v_{units\_invested\_available}(u, n, t)) \\
 &\sum_{u, n, t} v_{unit\_flow}(u, t) = p_{demand}(n, t)
 \end{aligned}$$

So this is a single problem that can't be decoupled over  $t$  because the investment variables `units_invested_available` couple the timesteps together. If `units_invested_available` were a constant in the problem, then all  $t$ 's could be solved individually. This is the basic idea in Benders decomposition. We decompose the problem into a master problem and sub problems with the master problem optimising the coupling investment variables which are treated as constants in the sub problems.

The master problem in the initial benders iteration is simply to minimise total investment costs:

$$\begin{aligned}
 &minimise Z : \\
 &Z \geq \sum_u p_{unit\_investment\_cost}(u) \cdot v_{units\_invested}(u)
 \end{aligned}$$

The solution to this problem yields values for the investment variables which are fixed as ' $\overline{units\_invested_u}$ ' in the sub problem and will be zero in the first iteration.

The sub problem for benders iteration  $b$  then becomes :

*minimise :*

$$obj_b = + \sum_{u,n,t} p_{operational\_cost}(u) \cdot v_{unit\_flow}(u, n, t)$$

*subject to :*

$$v_{unit\_flow}(u, n, t) \leq p_{unit\_capacity}(u) \cdot (v_{units\_available}(u, t) + p_{units\_invested\_available}(u, t)) \quad \mu_{b,u,t}$$

$$\sum_{u,n,t} v_{unit\_flow}(u, n, t) = p_{demand}(n, t)$$

This sub problem can be solved individually for each  $t$ . This is pretty trivial in this small example but if we consider a single  $t$  to be a single rolling horizon instead, decoupling the investment variables means that each rolling horizon can be solved individually rather than having to solve the entire model horizon as a single problem.

$$\mu_{u,t}$$

is the marginal value of the capacity constraint and can be interpreted as the decrease in the objective function at time  $t$  for an additional MW of flow from unit  $u$ . This information is used to construct a benders cut which represents the reduction in the sub problem objective function which is possible in this benders iteration by adjusting the variable  $units\_investment$ . This is effectively the decrease in operating costs possible by adding another unit of type  $u$  and is expressed as :

$$obj_b + \sum_{u,t} p_{unit\_capacity}(u, n, t) \cdot \mu_{b,u,t} \cdot (v_{units\_invested}(u, t) - p_{units\_invested}(u, b, t))$$

In the first benders iteration, the value of the investment variables will have been zero so  $p_{units\_invested}(u, b, t)$  will have the value of zero and thus the expression represents the total reduction in cost from an addition of a new unit of type  $u$ . This Benders cut is added to the master problem which then becomes, for each subsequent benders iteration,  $b$ :

*minimise Z :*

$$Z \geq \sum_{u,t} p_{unit\_investment\_cost}(u) \cdot v_{units\_invested}(u, t)$$

*subject to :*

$$Z \geq + \sum_u p_{unit\_investment\_cost}(u) \cdot v_{units\_invested}(u, t)$$

$$+ \sum_{u,t} p_{unit\_capacity}(u, t) \cdot \mu_{b,u,t} \cdot (v_{units\_invested}(u, t) - p_{units\_invested}(u, b, t)) \quad \forall b$$

Note the benders cuts are added as inequalities because they represent an upper bound on the value we are going to get from adjusting the master problem variables in that benders iteration. If we consider the example of renewable generation - because it's marginal cost is zero, on the first benders iteration, it could look like there would be a lot of value in increasing the capacity because of the marginal values from the sub problems. However, when the capacity variables are increased accordingly and curtailment occurs in the sub-problems, the marginal values will be zero when curtailment occurs and so, other resources may become optimal in subsequent iterations.

This is a simple example but it illustrates the general strategy. The algorithm pseudo code looks something like this:

```

initialise master problem
initialise sub problem
solve first master problem
create master problem variable time series
solve rolling spine opt model
save zipped marginal values
while master problem not converged
    update master problem
    solve master problem
    update master problem variable timeseries for benders iteration b
    rewind sub problem
    update sub problem
    solve rolling spine opt model
    save zipped marginal values
    test for convergence
end

```

## Benders cuts

The benders cuts for the problem including all investments in candidate connections, storages and units is given below.

$$\begin{aligned} &v_{objective\_lower\_bound}(b) \\ &\geq \\ &+ \sum_{u,s,t} p_{units\_invested\_available\_mv}(u,t,b) \cdot [v_{units\_invested\_available}(u,s,t) - p_{units\_invested\_available\_bi}(u,t,b)] \\ &+ \sum_{c,s,t} p_{connections\_invested\_available\_mv}(c,t,b) \cdot [v_{connections\_invested\_available}(c,s,t) - p_{connections\_invested\_available\_bi}(c,t,b)] \\ &+ \sum_{n,s,t} p_{storages\_invested\_available\_mv}(n,t,b) \cdot [v_{storages\_invested\_available}(n,s,t) - p_{storages\_invested\_available\_bi}(n,t,b)] \end{aligned}$$

where

$$p_{units\_invested\_available\_mv}$$

is the reduced cost of the [units\\_invested\\_available](#) fixed sub-problem variable, representing the reduction in operating costs possible from an investment in a [unit](#) of this type,  $p_{connections\_invested\_available\_mv}$  is the reduced cost of the [connections\\_invested\\_available](#) fixed sub-problem variable, representing the reduction in operating costs possible from an investment in a [connection](#) of this type,  $p_{storages\_invested\_available\_mv}$  is the reduced cost of the [storages\\_invested\\_available](#) fixed sub-problem variable, representing the reduction in operating costs possible from an investment in a [storage](#) of this type,  $p_{units\_invested\_available\_bi}(u,t,b)$  is the value of the fixed sub problem variable [units\\_invested\\_available](#)(u,t) in benders iteration b,  $p_{connections\_invested\_available\_bi}(c,t,b)$  is the value of the fixed sub problem variable [connections\\_invested\\_available](#)(c,t) in benders iteration b and  $p_{storages\_invested\_available\_bi}(n,t,b)$  is the value of the fixed sub problem variable [storages\\_invested\\_available](#)(n,t) in benders iteration b

## User constraints

### Unit constraint

The [unit\\_constraint](#) is a generic data-driven [custom constraint](#), which allows for defining constraints involving multiple [units](#), [nodes](#), or [connections](#). The [constraint\\_sense](#) parameter changes the sense of the [unit\\_constraint](#), while the [right\\_hand\\_side](#) parameter allows for defining the constant terms of the constraint.

Coefficients for the different [variables](#) appearing in the [unit\\_constraint](#) are defined using relationships, like e.g. [unit\\_from\\_node\\_unit\\_constraint](#) and [connection\\_to\\_node\\_unit\\_constraint](#) for [unit\\_flow](#) and [connection\\_flow](#) variables, or [unit\\_unit\\_constraint](#) and [node\\_unit\\_constraint](#) for [units\\_on](#), [units\\_started\\_up](#), and [node\\_state](#) variables.

For more information, see the dedicated article on [Unit Constraints](#)

$$\begin{aligned}
& + \sum_{u,n \in \text{unit\_node\_unit\_constraint}(uc),t,s} \\
& \left\{ \begin{array}{ll} \sum_{op} v_{\text{unit\_flow\_op}}(u,n,d,op,s,t) \cdot p_{\text{unit\_flow\_coefficient}}(u,n,op,uc,s,t) & \text{if } |\text{operating\_points}(u)| > 1 \\ v_{\text{unit\_flow}}(u,n,d,s,t) \cdot p_{\text{unit\_flow\_coefficient}}(u,n,uc,s,t) & \text{otherwise} \end{array} \right. \\
& + \sum_{u \in \text{unit\_unit\_constraint}(uc),t,s} v_{\text{units\_started\_up}}(u,s,t) \cdot p_{\text{units\_started\_up\_coefficient}}(u,uc,s,t) \\
& + \sum_{u \in \text{unit\_unit\_constraint}(uc),t,s} v_{\text{units\_on}}(u,s,t) \cdot p_{\text{units\_on\_coefficient}}(u,uc,s,t) \\
& + \sum_{c,n \in \text{connection\_node\_unit\_constraint}(uc),t,s} v_{\text{connection\_flow}}(c,n,d,s,t) \cdot p_{\text{connection\_flow\_coefficient}}(c,n,uc,s,t) \\
& + \sum_{n \in \text{node\_unit\_constraint}(uc),t,s} v_{\text{node\_state}}(n,s,t) \cdot p_{\text{node\_state\_coefficient}}(n,uc,s,t) \\
& + \sum_{n \in \text{node\_unit\_constraint}(uc),t,s} p_{\text{demand}}(n,s,t) \cdot p_{\text{demand\_coefficient}}(n,uc,s,t) \\
& \left\{ \begin{array}{ll} == & \text{if } p_{\text{constraint\_sense}}(uc) = "==" \\ >= & \text{if } p_{\text{constraint\_sense}}(uc) = ">=" \\ <= & \text{otherwise} \end{array} \right. \\
& + p_{\text{right\_hand\_side}}(uc,t,s) \\
& \forall uc,t,s \in \text{constraint\_unit\_constraint\_indices}
\end{aligned}$$

## Objective function

The objective function of SpineOpt expresses the minimization of the total system costs associated with maintaining and operating the considered energy system.

$$\begin{aligned} \min obj = & v_{unit\_investment\_costs} + v_{connection\_investment\_costs} + v_{storage\_investment\_costs} \\ & + v_{fixed\_om\_costs} + v_{variable\_om\_costs} + v_{fuel\_costs} + v_{start\_up\_costs} \\ & + v_{shut\_down\_costs} + v_{ramp\_costs} + v_{res\_proc\_costs} \\ & + v_{renewable\_curtailment\_costs} + v_{connection\_flow\_costs} + v_{taxes} + v_{objective\_penalties} \end{aligned}$$

Note that each cost term is reflected here as a separate variable that can be expressed mathematically by the equations below. All cost terms are weighted by the associated scenario and temporal block weights. To enhance readability and avoid writing a product of weights in every cost term, all weights are combined in a single weight parameter  $p_{weight}(\dots)$ . As such, the indices associated with each weight parameter indicate which weights are included.

## Unit investment costs

To take into account unit investments in the objective function, the parameter [unit\\_investment\\_cost](#) can be defined. For all tuples of (unit, scenario, timestep) in the set `units_invested_available_indices` for which this parameter is defined, an investment cost term is added to the objective function if a unit is invested in during the current optimization window. The total unit investment costs can be expressed as:

$$\begin{aligned} v_{unit\_investment\_costs} \\ = & \sum_{\substack{(u,s,t) \in \text{units\_invested\_available\_indices:} \\ u \in \text{ind}(p_{unit\_investment\_cost})}} v_{units\_invested}(u, s, t) \cdot p_{unit\_investment\_cost}(u, s, t) \cdot p_{weight}(u, s, t) \end{aligned}$$

## Connection investment costs

To take into account connection investments in the objective function, the parameter [connection\\_investment\\_cost](#) can be defined. For all tuples of (connection, scenario, timestep) in the set `connections_invested_available_indices` for which this parameter is defined, an investment cost term is added to the objective function if a connection is invested in during the current optimization window. The total connection investment costs can be expressed as:

$$\begin{aligned} v_{connection\_investment\_costs} \\ = & \sum_{\substack{(conn,s,t) \in \text{connections\_invested\_available\_indices:} \\ conn \in \text{ind}(p_{connection\_investment\_cost})}} v_{connections\_invested}(conn, s, t) \cdot p_{connection\_investment\_cost}(conn, s, t) \cdot p_{weight}(conn, s, t) \end{aligned}$$

## Storage investment costs

To take into account storage investments in the objective function, the parameter [storage\\_investment\\_cost](#) can be defined. For all tuples of (node, scenario, timestep) in the set `storages_invested_available_indices` for which this parameter is defined, an investment cost term is added to the objective function if a node storage is invested in during the current optimization window. The total storage investment costs can be expressed as:

$$\begin{aligned} v_{storage\_investment\_costs} \\ = & \sum_{\substack{(n,s,t) \in \text{storages\_invested\_available\_indices:} \\ n \in \text{ind}(p_{storage\_investment\_cost})}} v_{storages\_invested}(n, s, t) \cdot p_{storage\_investment\_cost}(n, s, t) \cdot p_{weight}(n, s, t) \end{aligned}$$

## Fixed O&M costs

Fixed operation and maintenance costs associated with a specific unit can be accounted for by defining the parameters [fom\\_cost](#) and [unit\\_capacity](#). For all tuples of (unit, {node,node\_group}, direction) for which these parameters are defined, and for which tuples (unit, scenario, timestep) exist in the set `units_on_indices`, a fixed O&M cost term is added to the objective function. Note that, as the `units_on_indices` are used to retrieve the relevant time slices, the unit of the [fom\\_cost](#) parameter should be given per resolution of the [units\\_on](#). The total fixed O&M costs can be expressed as:

$$\begin{aligned}
v_{fixed\_om\_costs} &= \sum_{\substack{(u,n,d) \in ind(p_{unit\_capacity}): \\ u \in ind(p_{fom\_cost})}} \sum_{(u,s,t) \in units\_on\_indices} p_{unit\_capacity}(u, n, d, s, t) \cdot p_{number\_of\_units}(u, s, t) \cdot p_{fom\_cost}(u, s, t) \cdot p_{weight}(t) \cdot p_{duration}(t)
\end{aligned}$$

## Variable O&M costs

Variable operation and maintenance costs associated with a specific unit can be accounted for by defining the parameter ([vom\\_cost](#)). For all tuples of (unit, {node,node\_group}, direction, scenario, timestep) in the set `unit_flow_indices` for which this parameter is defined, a variable O&M cost term is added to the objective function. As the parameter [vom\\_cost](#) is a dynamic parameter, the cost term is multiplied with the duration of each timestep. The total variable O&M costs can be expressed as:

$$\begin{aligned}
v_{variable\_om\_costs} &= \sum_{\substack{(u,n,d,s,t) \in unit\_flow\_indices: \\ (u,n,d) \in ind(p_{vom\_cost})}} v_{unit\_flow}(u, n, d, s, t) \cdot p_{vom\_cost}(u, n, d, s, t) \cdot p_{weight}(n, s, t) \cdot p_{duration}(t)
\end{aligned}$$

## Fuel costs

Fuel costs associated with a specific unit can be accounted for by defining the parameter [fuel\\_cost](#). For all tuples of (unit, {node,node\_group}, direction, scenario, timestep) in the set `unit_flow_indices` for which this parameter is defined, a fuel cost term is added to the objective function. As the parameter [fuel\\_cost](#) is a dynamic parameter, the cost term is multiplied with the duration of each timestep. The total fuel costs can be expressed as:

$$\begin{aligned}
v_{fuel\_costs} &= \sum_{\substack{(u,n,d,s,t) \in unit\_flow\_indices: \\ (u,n,d) \in ind(p_{fuel\_cost})}} v_{unit\_flow}(u, n, d, s, t) \cdot p_{fuel\_cost}(u, n, d, s, t) \cdot p_{weight}(n, s, t) \cdot p_{duration}(t)
\end{aligned}$$

## Connection flow costs

To account for operational costs associated with flows over a specific connection, the [connection\\_flow\\_cost](#) parameter can be defined. For all tuples of (conn, {node,node\_group}, direction, scenario, timestep) in the set `connection_flow_indices` for which this parameter is defined, a connection flow cost term is added to the objective function. The total connection flow costs can be expressed as:

$$\begin{aligned}
v_{connection\_flow\_costs} &= \sum_{\substack{(conn,n,d,s,t) \in connection\_flow\_indices: \\ conn \in ind(p_{connection\_flow\_cost})}} v_{connection\_flow}(conn, n, d, s, t) \cdot p_{connection\_flow\_cost}(conn, s, t) \cdot p_{weight}(n, s, t) \cdot p_{duration}(t)
\end{aligned}$$

## Start up costs

Start up costs associated with a specific unit can be included by defining the [start\\_up\\_cost](#) parameter. For all tuples of (unit, scenario, timestep) in the set `units_on_indices` for which this parameter is defined, a start up cost term is added to the objective function. The total start up costs can be expressed as:

$$\begin{aligned}
v_{start\_up\_costs} &= \sum_{\substack{(u,s,t) \in units\_on\_indices: \\ u \in ind(p_{start\_up\_cost})}} v_{units\_started\_up}(u, s, t) \cdot p_{start\_up\_cost}(u, s, t) \cdot p_{weight}(u, s, t)
\end{aligned}$$

## Shut down costs

Shut down costs associated with a specific unit can be included by defining the [shut\\_down\\_cost](#) parameter. For all tuples of (unit, scenario, timestep) in the set `units_on_indices` for which this parameter is defined, a shut down cost term is added to the objective function. The total shut down costs can be expressed as:

$$\begin{aligned}
v_{shut\_down\_costs} &= \sum_{\substack{(u,s,t) \in units\_on\_indices: \\ u \in ind(p_{shut\_down\_cost})}} v_{units\_shut\_down}(u, s, t) \cdot p_{shut\_down\_cost}(u, s, t) \cdot p_{weight}(u, s, t)
\end{aligned}$$

# Ramping costs

To account for the ramping costs (up and down) associated with a specific unit, the parameters `ramp_up_cost` and `ramp_down_cost` can be defined. For all tuples of (unit, {node,node\_group}, direction, scenario, timestep) in the sets `ramp_up_unit_flow_indices` and `ramp_down_unit_flow_indices` for which `ramp_up_cost` and `ramp_down_cost` are defined, respectively, a ramping cost term is added to the objective function. The total ramping costs can be expressed as:

$$\begin{aligned} v_{\text{ramp\_costs}} &= \sum_{\substack{(u,n,d,s,t) \in \text{ramp\_up\_unit\_flow\_indices:} \\ (u,n,d) \in \text{ind}(p_{\text{ramp\_up\_cost}})}} v_{\text{ramp\_up\_unit\_flow}}(u,n,d,s,t) \cdot p_{\text{ramp\_up\_cost}}(u,n,d,s,t) \cdot p_{\text{weight}}(n,s,t) \cdot p_{\text{duration}}(t) \\ &+ \sum_{\substack{(u,n,d,s,t) \in \text{ramp\_down\_unit\_flow\_indices:} \\ (u,n,d) \in \text{ind}(p_{\text{ramp\_down\_cost}})}} v_{\text{ramp\_down\_unit\_flow}}(u,n,d,s,t) \cdot p_{\text{ramp\_down\_cost}}(u,n,d,s,t) \cdot p_{\text{weight}}(n,s,t) \cdot p_{\text{duration}}(t) \end{aligned}$$

# Reserve procurement costs

The procurement costs for reserves provided by a specific unit can be accounted for by defining the `reserve_procurement_cost` parameter. For all tuples of (unit, {node,node\_group}, direction, scenario, timestep) in the set `unit_flow_indices` for which this parameter is defined, a reserve procurement cost term is added to the objective function. The total reserve procurement costs can be expressed as:

$$\begin{aligned} v_{\text{res\_proc\_costs}} &= \sum_{\substack{(u,n,d,s,t) \in \text{unit\_flow\_indices:} \\ (u,n,d) \in \text{ind}(p_{\text{reserve\_procurement\_cost}})}} v_{\text{unit\_flow}}(u,n,d,s,t) \cdot p_{\text{reserve\_procurement\_cost}}(u,n,d,s,t) \cdot p_{\text{weight}}(n,s,t) \cdot p_{\text{duration}}(t) \end{aligned}$$

# Renewable curtailment costs

The curtailment costs of renewable units can be accounted for by defining the parameters `curtailment_cost` and `unit_capacity`. For all tuples of (unit, {node,node\_group}, direction) for which these parameters are defined, and for which tuples (unit, scenario, timestep\_long) exist in the set `units_on_indices`, and for which tuples (unit, {node,node\_group}, direction, scenario, timestep\_short) exist in the set `unit_flow_indices`, a renewable curtailment cost term is added to the objective function. The total renewable curtailment costs can be expressed as:

$$\begin{aligned} v_{\text{renewable\_curtailment\_costs}} &= \sum_{\substack{(u,n,d) \in \text{ind}(p_{\text{unit\_capacity}}): \\ u \in \text{ind}(p_{\text{curtailment\_cost}})}} \sum_{(u,s,t_{\text{long}}) \in \text{units\_on\_indices}} \sum_{(u,n,s,t_{\text{short}}) \in \text{unit\_flow\_indices}} \\ & (v_{\text{units\_available}}(u,s,t_{\text{long}}) \cdot p_{\text{unit\_capacity}}(u,n,d,s,t_{\text{short}}) \cdot p_{\text{unit\_conv\_cap\_to\_flow}}(u,n,d,s,t_{\text{short}}) \\ & - v_{\text{unit\_flow}}(u,n,d,s,t_{\text{short}})) \\ & \cdot p_{\text{curtailment\_cost}}(u,s,t_{\text{short}}) \cdot p_{\text{weight}}(n,s,t_{\text{short}}) \cdot p_{\text{duration}}(t_{\text{short}}) \end{aligned}$$

# Taxes

To account for taxes on certain commodity flows, the tax unit flow parameters (i.e., `tax_net_unit_flow`, `tax_out_unit_flow` and `tax_in_unit_flow`) can be defined. For all tuples of (unit, {node,node\_group}, direction, scenario, timestep) in the set `unit_flow_indices` for which these parameters are defined, a tax term is added to the objective function. The total taxes can be expressed as:

$$\begin{aligned} v_{\text{taxes}} &= \sum_{\substack{(u,n,d,s,t) \in \text{unit\_flow\_indices:} \\ n \in \text{ind}(p_{\text{tax\_net\_unit\_flow}}) \& d = \text{to\_node}}} v_{\text{unit\_flow}}(u,n,d,s,t) \cdot p_{\text{tax\_net\_unit\_flow}}(n,s,t) \cdot p_{\text{weight}}(n,s,t) \cdot p_{\text{duration}}(t) \\ &- \sum_{\substack{(u,n,d,s,t) \in \text{unit\_flow\_indices:} \\ n \in \text{ind}(p_{\text{tax\_net\_unit\_flow}}) \& d = \text{from\_node}}} v_{\text{unit\_flow}}(u,n,d,s,t) \cdot p_{\text{tax\_net\_unit\_flow}}(n,s,t) \cdot p_{\text{weight}}(n,s,t) \cdot p_{\text{duration}}(t) \\ &+ \sum_{\substack{(u,n,d,s,t) \in \text{unit\_flow\_indices:} \\ n \in \text{ind}(p_{\text{tax\_out\_unit\_flow}}) \& d = \text{from\_node}}} v_{\text{unit\_flow}}(u,n,d,s,t) \cdot p_{\text{tax\_out\_unit\_flow}}(n,s,t) \cdot p_{\text{weight}}(n,s,t) \cdot p_{\text{duration}}(t) \\ &+ \sum_{\substack{(u,n,d,s,t) \in \text{unit\_flow\_indices:} \\ n \in \text{ind}(p_{\text{tax\_in\_unit\_flow}}) \& d = \text{to\_node}}} v_{\text{unit\_flow}}(u,n,d,s,t) \cdot p_{\text{tax\_in\_unit\_flow}}(n,s,t) \cdot p_{\text{weight}}(n,s,t) \cdot p_{\text{duration}}(t) \end{aligned}$$



# Objective penalties

Penalty cost terms associated with the slack variables of a specific constraint can be accounted for by defining a [node\\_slack\\_penalty](#) parameter. For all tuples of  $(\{node, node\_group\}, scenario, timestep)$  in the set `node_slack_indices` for which this parameter is defined, a penalty term is added to the objective function. The total objective penalties can be expressed as:

$$\begin{aligned} &v_{objective\_penalties} \\ &= \sum_{(u,s,t) \in node\_slack\_indices} \left[ v_{node\_slack\_neg}(n, s, t) - v_{node\_slack\_pos}(n, s, t) \right] \cdot p_{node\_slack\_penalty}(n, s, t) \cdot p_{weight}(n, s, t) \cdot p_{duration}(t) \end{aligned}$$

# Temporal Framework

Spine Model aims to provide a high degree of flexibility in the temporal dimension across different components of the created model. This means that the user has some freedom to choose how the temporal aspects of different components of the model are defined. This freedom increases the variety of problems that can be tackled in Spine: from very coarse, long term models, to very detailed models with a more limited horizon, or a mix of both. The choice of the user on how this flexibility is used will lead to the temporal structure of the model.

The main components of flexibility consist of the following parts:

- The horizon that is modeled: end and start time
- Temporal resolution
- Possibility of a rolling optimization window
- Support for commonly used methods such as representative days

Part of the temporal flexibility in Spine is due to the fact that these options mentioned above can be implemented differently across different components of the model, which can be very useful when different markets are coupled in a single model. The resolution and horizon of the gas market can for example be taken differently than that of the electricity market. This documentation aims to give the reader insight in how these aspects are defined, and which objects are used for this.

We start by introducing the relevant objects with their parameters, and the relevant relationship classes for the temporal structure. Afterwards, we will discuss how this setting creates flexibility and will present some of the practical approaches to create a variety of temporal structures.

## Objects, relationships, and their parameters

In this section, the objects and relationships will be discussed that form the temporal structure together.

### Objects relevant for the temporal framework

For the objects, the relevant parameters will also be introduced, along with the type of values that are allowed, following the format below:

- 'parameter\_name' : "Allowed value type"

**model object**

Each `model` object holds general information about the model at hand. Here we only discuss the time related parameters:

- `model_start` and `model_end` : "Date time value"

These two parameters define the model horizon. A Datetime value is to be taken for both parameters, in which case they directly mark respectively the beginning and end of the modeled time horizon.

- `duration_unit` (optional): "minute or hour"

This parameters gives the unit of duration that is used in the model calculations. The default value for this parameter is 'minute'. E.g. if the `duration_unit` is set to `hour`, a `Duration` of one minute gets converted into  $1/60$  hours for the calculations.

- `roll_forward` (optional): "duration value"

This parameter defines how much the optimization window rolls forward in a rolling horizon optimization and should be expressed as a duration. In the practical approaches presented below, the rolling window optimization will be explained in more detail.

## `temporal_block` object

A temporal block defines the properties of the optimization that is to be solved in the current window. Most importantly, it holds the necessary information about the resolution and horizon of the optimization.

- `resolution` (optional): "duration value" or "array of duration values"

This parameter specifies the resolution of the temporal block, or in other words: the length of the timesteps used in the optimization run.

- `block_start` (optional): "duration value" or "Date time value"

Indicates the start of this temporal block.

- `block_end`(optional): "duration value" or "Date time value"

Indicates the end of this temporal block.

## Relationships relevant for the temporal framework

### `model_temporal_block` relationship

In this relationship, a model instance is linked to a temporal block. If this relationship doesn't exist - the temporal block is disregarded from this optimization model.

### `model_default_temporal_block` relationship

Defines the default temporal block used for model objects, which will be replaced when a specific relationship is defined for a model in `model__temporal_block`.

### `node__temporal_block` relationship

This relationship will link a node to a temporal block.

### `units_on__temporal_block` relationship

This relationship links the `units_on` variable of a unit to a temporal block and will therefore govern the time-resolution of the unit's online/offline status.

### `unit__investment_temporal_block` relationship

This relationship sets the temporal dimensions for investment decisions of a certain unit. The separation between this relationship and the `units_on__temporal_block`, allows the user for example to give a much finer resolution to a unit's on- or offline status than to its investment decisions.

### `model__default_investment_temporal_block` relationship

Defines the default temporal block used for investment decisions, which will be replaced when a specific relationship is defined for a unit in `unit__investment_temporal_block`.

## General principle of the temporal framework

The general principle of the Spine modeling temporal structure is that different temporal blocks can be defined and linked to different objects in a model. This leads to great flexibility in the temporal structure of the model as a whole. To illustrate this, we will discuss some of the possibilities that arise in this framework.

### One single `temporal_block`

#### Single solve with single block

The simplest case is a single solve of the entire time horizon (so `roll_forward` not defined) with a fixed resolution. In this case, only one temporal block has to be defined with a fixed resolution. Each node has to be linked to this `temporal_block`.

Alternatively, a variable resolution can be defined by choosing an array of durations for the `resolution` parameter. The sum of the durations in the array then have to match the length of the temporal block. The example below illustrates an optimization that spans one day for which the resolution is hourly in the beginning and then gradually decreases to a 6h resolution at the end.

- `temporal_block_1`
  - `block_start`: 0h (Alternative `DateTime`: e.g. 2030-01-01T00:00:00)
  - `block_end`: 1D (Alternative `DateTime`: e.g. 2030-01-02T00:00:00)

- `resolution: [1h 1h 1h 1h 2h 2h 2h 4h 4h 6h]`

Note that, as mentioned above, the `block_start` and `block_end` parameters can also be entered as absolute values, i.e. `DateTime` values.

## Rolling window optimization with single block

A model with a single `temporal_block` can also be optimized in a rolling horizon framework. In this case, the `roll_forward` parameter has to be defined in the `model` object. The `roll_forward` parameter will then determine how much the optimization moves forward with every step, while the size of the temporal block will determine how large a time frame is optimized in each step. To see this more clearly, let's take a look at an example.

Suppose we want to model a horizon of one week, with a rolling window size of one day. The `roll_forward` parameter will then be a duration value of `1d`. If we take the `temporal_block` parameters `block_start` and `block_end` to be the duration values `0h` and `1d` respectively, the model will optimize each day of the week separately. However, we could also take the `block_end` parameter to be `2d`. Now the model will start by optimizing day 1 and day 2 together, after which it keeps only the values obtained for the first day, and moves forward to optimize the second and third day together.

Again, a variable resolution can be implemented for the rolling window optimization. The sum of the durations must in this case match the size of the optimized window.

## Advanced usage: multiple `temporal_block` objects

### Single solve with multiple blocks

#### Disconnected time periods

Multiple temporal blocks can be used to optimize disconnected periods. Let's take a look at an example in which two temporal blocks are defined.

- `temporal_block_1`
  - `block_start: 0h`
  - `block_end: 4h`
- `temporal_block_2`
  - `block_start: 12h`
  - `block_end: 16h`

This example will lead to an optimization of the first four hours of the model horizon, and also of hour 12 to 16. By defining exactly the same relationships for the two temporal blocks, an optimization of disconnected periods is achieved for exactly the same model components. This leads to the possibility of implementing the widely used representative days method. If desired, it is possible to choose a different temporal resolution for the different `temporal_blocks`.

It is worth noting that dynamic [variables](#) like [node\\_state](#) and [units\\_on](#) merit special attention when using disconnected time periods. By default, when trying to access [variables](#) Variables outside the defined [temporal\\_blocks](#), *SpineOpt.jl* assumes such variables exist but allows them to take any values within specified bounds. If fixed initial conditions for the disconnected periods are desired, one needs to use parameters such as [fix\\_node\\_state](#) or [fix\\_units\\_on](#).

## Different regions/commodities in different resolutions

Multiple temporal blocks can also be used to model different regions or different commodities with a different resolution. This is especially useful when there is a certain region or commodity of interest, while other elements are connected to this but require less detail. For this kind of usage, the relationships that are defined for the temporal blocks will be different, as shown in the example below.

- `temporal_blocks`
  - `temporal_block_1`
    - `resolution: 1h`
  - `temporal_block_2`
    - `resolution: 2h`
- `nodes`
  - `node_1`
  - `node_2`
- `node_temporal_block` relationships
  - `node_1_temporal_block_1`
  - `node_2_temporal_block_2`

Similarly, the on- and offline status of a unit can be modeled with a lower resolution than the actual output of that unit, by defining the `units_on_temporal_block` relationship with a different temporal block than the one used for the `node_temporal_block` relationship (of the node to which the unit is connected).

## Rolling horizon with multiple blocks

### Rolling horizon with different window sizes

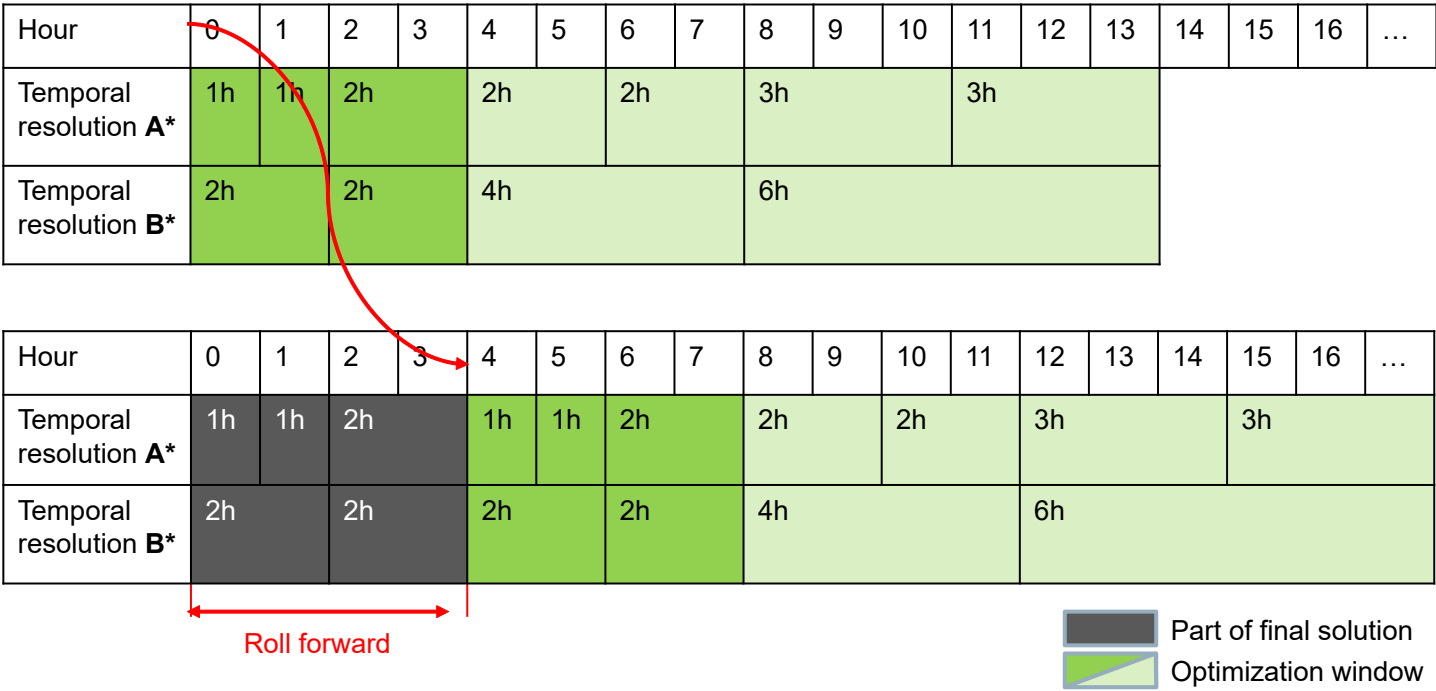
Similar to what has been discussed above in [Different regions/commodities in different resolutions](#), different commodities or regions can be modeled with a different resolution in the rolling horizon setting. The way to do it is completely analogous. Furthermore, when using the rolling horizon framework, a different window size can be chosen for the different modeled components, by simply using a different `block_end` parameter. However, using different [block\\_ends](#) e.g. for interconnected regions should be treated with care, as the variables for each region will only be generated for their respective [temporal\\_block](#), which in most cases will lead to inconsistent linking constraints.

## Putting it all together: rolling horizon with variable resolution that differs for different model components

Below is an example of an advanced use case in which a rolling horizon optimization is used, and different model components are optimized with a different resolution. By choosing the relevant parameters in the following way:

- model
  - roll\_forward: 4h
- temporal\_blocks
  - temporal\_block\_A
    - resolution: [1h 1h 2h 2h 2h 3h 3h]
    - block\_end: 14h
  - temporal\_block\_B
    - resolution: [2h 2h 4h 6h]
    - block\_end: 14h
- nodes
  - node\_1
  - node\_2
- node\_temporal\_block relationships
  - node\_1\_temporal\_block\_A
  - node\_2\_temporal\_block\_B

The two model components that are considered have a different resolution, and their own resolution is also varying within the optimization window. Note that in this case the two optimization windows have the same size, but this is not strictly necessary. The image below visualizes the first two window optimizations of this model.



\* A/B can correspond to different zones/regions (e.g., BE and FR) and/or a different market sector (e.g., Gas & Electricity)





# Stochastic Framework

Scenario-based stochastics in unit commitment and economic dispatch models typically only consider branching scenario trees. However, sometimes the available stochastic data doesn't span over the entire desired modelling horizon, or all the modelled phenomena. Especially with increasing interest in energy system integration and sector coupling, stochastic data of consistent quality and/or length might be hard to come by.

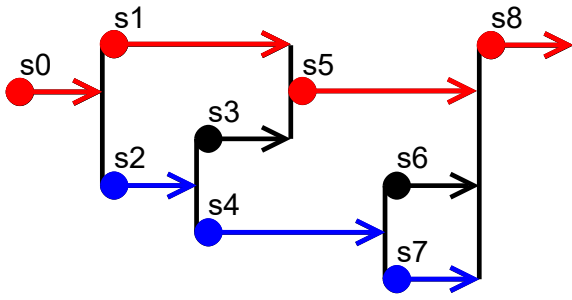
While these data issues can be circumvented by either cloning stochastic data across multiple scenario branches or generating dummy forecasts, they can result in inflated problem sizes. Furthermore, Ensuring realistic correlations between generated forecasts is extremely difficult, especially across multiple energy sectors.

The stochastic framework in *SpineOpt.jl* aims to support stochastic directed acyclic graphs (DAGs) instead of only branching trees, allowing for scenarios to converge later on in the modelled horizon. In addition, the framework allows for slightly different stochastic scenario graphs for different variables, making it easier to define e.g. variables common between all stochastic scenarios.

## Key concepts

Here, we briefly describe the key concepts required to understand the stochastic framework:

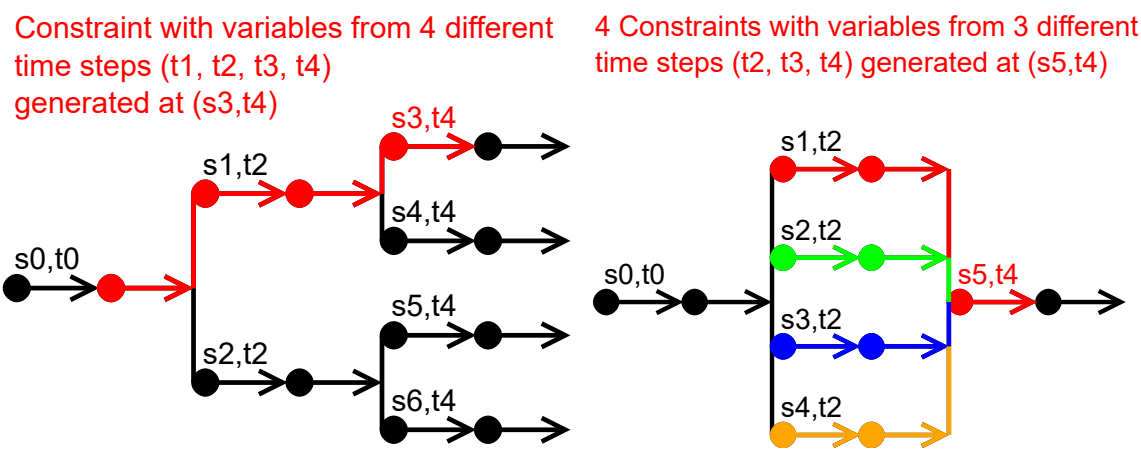
1. **stochastic\_scenario** is essentially just a label for an alternative period of time, describing one possibility of what may come to pass. Even in deterministic modelling with *SpineOpt.jl*, a single **stochastic\_scenario** is required for labelling the deterministic timeline.
2. **Stochastic DAG** is the directed acyclic graph describing the **parent\_stochastic\_scenario\_child\_stochastic\_scenario** relationships between the **stochastic scenarios**. The key difference between a *stochastic DAG* and a traditional *stochastic tree* is that the scenarios are allowed to have multiple parents, making it possible to converge scenarios into each other in addition to branching.
3. **Stochastic path** is a unique sequence of **stochastic scenarios** for traversing the *stochastic DAG*. Every (*finite*) *stochastic DAG* has a limited number of *full stochastic paths* that traverse it from roots (*scenarios without parents*) to leaves (*scenarios without children*). Here, we use the term *stochastic path* to refer to any subset of scenarios within a *full stochastic path*.
4. **stochastic\_structure** is essentially a "realization" of the *stochastic DAG*, with additional information like the **stochastic\_scenario\_end** and **weight\_relative\_to\_parents Parameters**. These become relevant when we start discussing interactions between different **stochastic structures**.



The above figure presents an example *stochastic DAG* with the individual stochastic scenarios labelled from `s0`–`s8`. An example *full stochastic path* [`s0`, `s1`, `s5`, `s8`] is highlighted in red, while an example *stochastic path* [`s2`, `s4`, `s7`] is highlighted in blue.

# General idea in brief

The major issue with *stochastic DAGs* compared to *stochastic trees*, is that indexing constraints that include variables from multiple time steps (*henceforth referred to as "dynamic constraints"*) needs rethinking. With *stochastic trees*, constraints can always be unambiguously indexed using `(stochastic_scenario, last_time_step)`, since all stochastic scenarios only have a single parent. However, this is no longer the case for *stochastic DAGs*, as illustrated in the figures below:



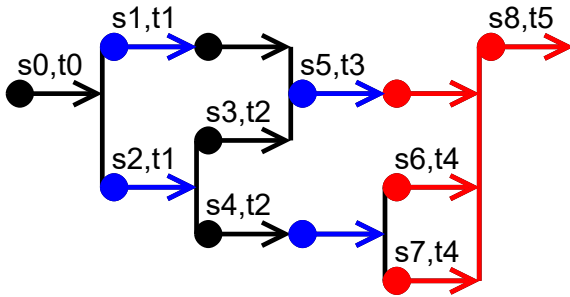
The example on the left illustrates the "traditional" indexing in branching *stochastic trees*, where backtracking through the tree always leads to unambiguous `(stochastic_scenario, time_step)` indices. The example on the right shows a similar situation in a *stochastic DAG*, where backtracking through the DAG leads to four different `(stochastic_scenario, time_step)` indices, and thus requires four constraints to be generated and indexed.

# Stochastic path indexing

As discussed in the previous section, dynamic constraints in *stochastic DAGs* cannot be unambiguously indexed using a single `(stochastic_scenario, time_step)`. However, they *can* be unambiguously indexed using `(stochastic_path, time_step)`, where the *stochastic path* is the unique sequence of stochastic scenarios traversing the DAG. Since there are only a limited number of ways to traverse the DAG, represented by the *full stochastic paths*, we can identify the number of unique paths necessary for constraint generation as follows:

1. Identify all unique *full stochastic paths*, meaning all the possible ways of traversing the DAG from roots to leaves.
2. Find all the *stochastic scenarios* that are active on all the *time steps* included in the constraint.
3. Find all the unique *stochastic paths* by intersecting the set of active *stochastic scenarios* with the *full stochastic paths*.
4. Generate constraints over each unique *stochastic path* found in step 3.

## Example dynamic constraint generation



The above figure shows examples of two different dynamic constraints generated in a *stochastic DAG*: the red constraint including variables from timesteps  $t_4$ – $t_5$  and the blue constraint including variables from timesteps  $t_1$ ,  $t_3$ . The *full stochastic paths* for traversing the above DAG are as follows:

1.  $[s_0, s_1, s_5, s_8]$
2.  $[s_0, s_2, s_3, s_5, s_8]$
3.  $[s_0, s_2, s_4, s_6, s_8]$
4.  $[s_0, s_2, s_4, s_7, s_8]$

For the red constraint, the *stochastic scenarios*  $s_5$ – $s_8$  are active on the *time steps*  $t_4$ – $t_5$ . All the above *full stochastic paths* include at least two of the active *stochastic scenarios*, but full paths 1 and 2 both produce an identical path  $[s_5, s_8]$ , so the set of unique *stochastic paths* for the red constraint becomes:

1.  $[s_5, s_8]$
2.  $[s_6, s_8]$
3.  $[s_7, s_8]$

There are no paths  $[s_5, s_6]$ ,  $[s_5, s_7]$ ,  $[s_6, s_7]$  since following the DAG one cannot start from  $s_5$  and end up in  $s_6$ , even though these *stochastic scenarios* are active.

The blue constraint illustrates a case where the time step range is non-continuous. The active *stochastic scenarios* on  $t_1$ ,  $t_3$  are  $s_1$ ,  $s_2$ ,  $s_4$ ,  $s_5$ , so again by comparing these to the *full stochastic paths* we get:

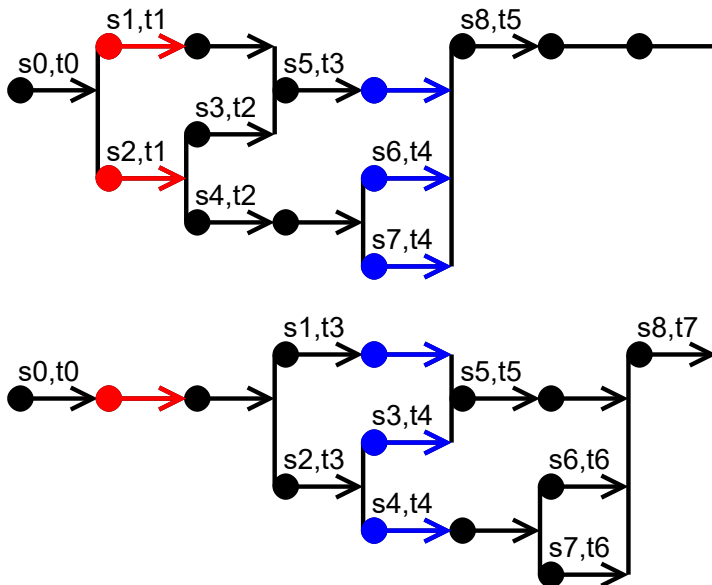
1.  $[s_1, s_5]$
2.  $[s_2, s_5]$

### 3. [ s2, s4 ]

In this case, the *full stochastic paths* 3 and 4 both produce the path [ s2, s4 ], so only three unique constraints need to be generated. Again, the path [ s1, s4 ] is invalid, since the DAG cannot be traversed from s1 to s4.

## Interaction between different stochastic structures

*Stochastic path* indexing in constraints also allows for "distorting" the *stochastic DAG* in different parts of the model. As long as the *stochastic DAG* itself isn't changed, meaning the [parent\\_stochastic\\_scenario\\_child\\_stochastic\\_scenario](#) relationships and the resulting *full stochastic paths*, we can actually define different *stochastic structures* and still be able to handle constraint generation between them. This is due to the fact that when determining the *stochastic paths*, it makes no difference whether we're looking at the same *stochastic structure* at different *time steps*, or at two *stochastic structures*, one of which has been delayed, on the same *time step*. This is illustrated by the figure below:



The above represents constraint generation over two *stochastic structures*, where the lower *stochastic structure* has been delayed in respect to the one above. Nevertheless, the procedure for finding the *stochastic paths* for the constraints remains identical to the previous example:

1. Identify all unique *full stochastic paths*, meaning all the possible ways of traversing the DAG. As long as the DAG remains the same between all the involved *stochastic structures*, the pathing remains the same.
2. Find all the *stochastic scenarios* that are active on all the *stochastic structures* and *time steps* included in the constraint.
3. Find all the unique stochastic paths by intersecting the set of active scenarios with the *full stochastic paths*.
4. Generate constraints over each unique stochastic path found in step 3.

# Stochastics in the model data structure

While the [Key concepts](#) and [General idea in brief](#) sections go over the stochastic framework in *SpineOpt.jl* in a more general sense, here we'll go over how to set up stochastics using *SpineOpt.jl* data structure. Simple step-by-step examples are also provided in the [Example of deterministic stochastics](#), [Example of branching stochastics](#), and [Example of converging stochastics](#) sections further below. We won't go into too much detail about the related [Object Classes](#), [Relationship Classes](#), or [Parameters](#), since those can be found in their respective sections. Introductions to these concepts can also be found in the [Structural object classes](#) and [Structural relationship classes](#) sections, if necessary.

## Setting up the stochastic framework

As with all things in *SpineOpt.jl*, you'll want to start with adding the desired number of objects to the relevant [Object Classes](#), as one cannot define relationships over objects that don't exist. For the stochastic framework, this means creating at least one [stochastic\\_scenario](#) and [stochastic\\_structure](#) object each. This needs to be done even if your model is fully deterministic, as even the deterministic structure needs to be labelled for *SpineOpt.jl* to recognize that it exists.

Next, if your model has multiple [stochastic\\_scenario](#) objects, you'll want to define how they are related using the [parent\\_stochastic\\_scenario\\_\\_child\\_stochastic\\_scenario](#) relationship. This relationship essentially defines the *stochastic DAG*, as well as all the possible *stochastic paths*, explained in the [Key concepts](#) section. Unless the [parent\\_stochastic\\_scenario\\_\\_child\\_stochastic\\_scenario](#) relationship is defined, there won't be a *stochastic DAG*, and all [stochastic\\_scenario](#) objects will be assumed to be completely independent of each other.

Now that you've set up the desired [stochastic\\_scenario](#) and [stochastic\\_structure](#) objects, as well as defined the *stochastic DAG* using the [parent\\_stochastic\\_scenario\\_\\_child\\_stochastic\\_scenario](#) relationship, it's time to define the properties of the [stochastic\\_structure](#) objects using the [stochastic\\_structure\\_\\_stochastic\\_scenario](#) relationship, and the [stochastic\\_scenario\\_end](#) and [weight\\_relative\\_to\\_parents](#) [Parameters](#) therein. You'll always have to define at least one [stochastic\\_structure\\_\\_stochastic\\_scenario](#) relationship, as the [stochastic\\_structure](#) object is what connects the [Systemic object classes](#) to the stochastic framework. [stochastic\\_structure\\_\\_stochastic\\_scenario](#) relationship holds two key [Parameters](#):

- [weight\\_relative\\_to\\_parents](#) defines the coefficient the corresponding [stochastic\\_scenario](#) has in the [Objective function](#), and needs to be defined for each [stochastic\\_scenario](#) included in the [stochastic\\_structure](#). The weight is relative to the parents of the [stochastic\_scenario], and is calculated as presented below.

```
# For root `stochastic_scenarios` (meaning no parents)

weight(scenario) = weight_relative_to_parents(scenario)

# If not a root `stochastic_scenario`
```

```
weight(scenario) = sum([weight(parent) * weight_relative_to_parents(scenario)] for p
```

- **stochastic\_scenario\_end** is a `Duration` type parameter that tells when the **stochastic\_scenario** ends in relation to the start of the current optimization. When defined, the **stochastic\_scenario** ends at the defined point in time, and spawns its children according to **parent\_stochastic\_scenario\_child\_stochastic\_scenario**, if any. **Note that this means the children are included in the **stochastic\_structure**, even without an explicit relationship!** If **stochastic\_scenario\_end** isn't defined, the **stochastic\_scenario** is assumed to go on indefinitely.

Finally, with all the pieces in place, we'll need to connect the defined **stochastic\_structure** objects to the desired objects in the **Systemic object classes** using the **Structural relationship classes** like **node\_stochastic\_structure** etc. Here, we essentially tell which parts of the modelled system use which **stochastic\_structure**. Since creating each of these relationships individually can be a bit of a pain, there are a few **Meta relationship classes** like the **model\_default\_stochastic\_structure**, that can be used to set **model**-wide defaults that are used if specific relationships are missing.

## Example of deterministic stochastics

Here, we'll demonstrate step-by-step how to create the simplest possible stochastic frame: the fully deterministic one.

1. Create a **stochastic\_scenario** called e.g. `realization` and a **stochastic\_structure** called e.g. `deterministic`.
2. We can skip the **parent\_stochastic\_scenario\_child\_stochastic\_scenario** relationship, since there isn't a *stochastic DAG* in this example, and the default behaviour of each **stochastic\_scenario** being independent works for our purposes (*only one **stochastic\_scenario** anyhow*).
3. Create the **stochastic\_structure\_stochastic\_scenario** relationship for `(deterministic, realization)`, and set its **weight\_relative\_to\_parents** parameter to 1. We don't need to define the **stochastic\_scenario\_end** parameter, as we want the `realization` to go on indefinitely.
4. Relate the `deterministic` **stochastic\_structure** to all the desired system objects using the appropriate **Structural relationship classes**, or use the **model**-level default **Meta relationship classes**.

## Example of branching stochastics

Here, we'll demonstrate step-by-step how to create a simple branching stochastic tree, where one scenario branches into three at a specific point in time.

1. Create four **stochastic\_scenario** objects called e.g. `realization`, `forecast1`, `forecast2`, and `forecast3`, and a **stochastic\_structure** called e.g. `branching`.
2. Define the *stochastic DAG* by creating the **parent\_stochastic\_scenario\_child\_stochastic\_scenario** relationships for `(realization, forecast1)`, `(realization, forecast2)`, and `(realization, forecast3)`.

forecast3).

3. Create the `stochastic_structure__stochastic_scenario` relationship for (branching, realization), (branching, forecast1), (branching, forecast2), and (branching, forecast3).
4. Set the `weight_relative_to_parents` parameter to 1 and the `stochastic_scenario_end` parameter e.g. to 6h for the `stochastic_structure__stochastic_scenario` relationship (branching, realization). Now, the realization `stochastic_scenario` will end after 6 hours of time steps, and its children (forecast1, forecast2, and forecast3) will become active.
5. Set the `weight_relative_to_parents Parameters` for the (branching, forecast1), (branching, forecast2), and (branching, forecast3) `stochastic_structure__stochastic_scenario` relationships to whatever you desire, e.g. 0.33 for equal probabilities across all forecasts.
6. Relate the branching `stochastic_structure` to all the desired system objects using the appropriate [Structural relationship classes](#), or use the `model`-level default [Meta relationship classes](#).

## Example of converging stochastics

Here, we'll demonstrate step-by-step how to create a simple *stochastic DAG*, where both branching and converging occurs. This example relies on the previous [Example of branching stochastics](#), but adds another `stochastic_scenario` at the end, which is a child of the forecast1, forecast2, and forecast3 scenarios.

1. Follow the steps 1-5 in the previous [Example of branching stochastics](#), except call the `stochastic_structure` something different, e.g. converging.
2. Create a new `stochastic_scenario` called e.g. converged\_forecast.
3. Alter the *stochastic DAG* by creating the `parent_stochastic_scenario__child_stochastic_scenario` relationships for (forecast1, converged\_forecast), (forecast2, converged\_forecast), and (forecast3, converged\_forecast). Now all three forecasts will converge into a single converged\_forecast.
4. Add the `stochastic_structure__stochastic_scenario` relationship for (converging, converged\_forecast), and set its `weight_relative_to_parents` parameter to 1. Now, all the probability mass in forecast1, forecast2, and forecast3 will be summed up back to the converged\_forecast.
5. Set the `stochastic_scenario_end Parameters` of the `stochastic_structure__stochastic_scenario` relationships (converging, forecast1), (converging, forecast2), and (converging, forecast3) to e.g. 12h, so that all three scenarios end at the same time and the converged\_forecast becomes active.
6. Relate the converging `stochastic_structure` to all the desired system objects using the appropriate [Structural relationship classes](#), or use the `model`-level default [Meta relationship classes](#).

## Working with stochastic updating data



Now that we've discussed how to set up stochastics for SpineOpt, let's focus on stochastic data. The most complex form of input data SpineOpt can currently handle is both *stochastic* and *updating*, meaning that the values the parameter takes can depend on both the `stochastic_scenario`, and the *analysis time (first time step)* of each solve. However, just *stochastic* or just *updating* cases are supported as well, using the same input data format.

In SpineOpt, stochastic data uses the `Map` data type from [SpineInterface.jl](#). Essentially, `Maps` are general indexed data containers, which SpineOpt tries to interpret as stochastic data. Every time SpineOpt calls a parameter, it passes the `stochastic_scenario` and *analysis time* as keyword arguments to the parameter, but depending on the parameter type, it doesn't necessarily do anything with that information. For `Map` type parameters, those keyword arguments are used for navigating the indices of the `Map` to try and find the corresponding value. If the `Map` doesn't include the `stochastic_scenario` index it's looking for, it assumes there's no *stochastic* information in the `Map` and carries on to search for *analysis time* indices. This logic is useful for defining both *stochastic* and *updating* data, as well as either case by itself, as shown in the following examples.

## Example of stochastic data

By *stochastic data*, we mean parameter values that depend only on the `stochastic_scenario`. In such a case, the input data must be formatted as a `Map` with the following structure

<code>stochastic_scenario</code>	value
scenario1	value1
scenario2	value2

where `stochastic_scenario` indices are simply `Strings` corresponding to the names of the `stochastic_scenario` objects. The *values* can be whatever data types [SpineInterface.jl](#) supports, like `Constants`, `DateTimes`, `Durations`, or `TimeSeries`. In the above example, the parameter will take `value1` in `scenario1`, and `value2` in `scenario2`. Note that since there's no *analysis time* index in this example, the *values* are used regardless of the *analysis time*.

## Example of updating data

By *updating data*, we mean parameter values that depend only on the *analysis time*. In such a case, the input data must be formatted as a `Map` with the following structure

analysis time	value
2000-01-01T00:00:00	value1
2000-01-01T12:00:00	value2



where the *analysis time* indices are `DateTime` values. The *values* can be whatever data types [SpineInterface.jl](#) supports, like `Constants`, `Datetimes`, `Durations`, or `TimeSeries`. In the above example, the parameter will take `value1` if the first time step of the current simulation is between `2000-01-01T00:00:00` and `2000-01-01T12:00:00`, and `value2` if the first time step of the simulation is after `2000-01-01T12:00:00`. Note that since there's no `stochastic_scenario` index in this example, the *values* are used regardless of the `stochastic_scenario`.

## Example of stochastic updating data

By *stochastic updating data*, we mean parameter values that depend on both the `stochastic_scenario` and the *analysis time*. In such a case, the input data must be formatted as a `Map` with the following structure

<code>stochastic_scenario</code>	<code>analysis time</code>	<code>value</code>
<code>scenario1</code>	<code>2000-01-01T00:00:00</code>	<code>value1</code>
<code>scenario1</code>	<code>2000-01-01T12:00:00</code>	<code>value2</code>
<code>scenario2</code>	<code>2000-01-01T00:00:00</code>	<code>value3</code>
<code>scenario2</code>	<code>2000-01-01T12:00:00</code>	<code>value4</code>

where the `stochastic_scenario` indices are simply `Strings` corresponding to the names of the `stochastic_scenario` objects, and the *analysis time* indices are `DateTime` values. The *values* can be whatever data types [SpineInterface.jl](#) supports, like `Constants`, `Datetimes`, `Durations`, or `TimeSeries`. In the above example, the parameter will take `value1` if the first time step of the current simulation is between `2000-01-01T00:00:00` and `2000-01-01T12:00:00` and the parameter is called in `scenario1`, and `value3` in `scenario2`. If the first time step of the current simulation is after `2000-01-01T12:00:00`, the parameter will take `value2` in `scenario1`, and `value4` in `scenario2`.

## Constraint generation with stochastic path indexing

Every time a constraint might refer to variables either on different time steps or on different `stochastic scenarios` (meaning different `nodes` or `units`), the constraint needs to use stochastic path indexing in order to be correctly generated for arbitrary stochastic DAGs. In practise, this means following the procedure outlined below:

1. Identify all unique *full stochastic paths*, meaning all the possible ways of traversing the DAG. This is done along with generating the stochastic structure, so no real impact on constraint generation.
2. Find all the `stochastic scenarios` that are active on all the `stochastic structures` and `time slices` included in the constraint.
3. Find all the unique stochastic paths by intersecting the set of active scenarios with the *full stochastic paths*.

4. Generate constraints over each unique stochastic path found in step 3.

Steps 2 and 3 are the crucial ones, and are currently handled by separate `constraint_<constraint_name>_indices` functions. Essentially, these functions go through all the variables on all the time steps included in the constraint, collect the set of active `stochastic_scenarios` on each time step, and then determine the unique active stochastic paths on each time step. The functions pre-form the index set over which the constraint is then generated in the `add_constraint_<constraint_name>` functions.

# Unit commitment

To incorporate technical detail about (clustered) unit-commitment statuses of units, the online, started and shutdown status of units can be tracked and constrained in SpineOpt. In the following, relevant relationships and parameters are introduced and the general working principle is described.

## Key concepts for unit commitment

Here, we briefly describe the key concepts involved in the representation of (clustered) unit commitment models:

- [units\\_on](#) is an optimization variable that holds information about the on- or offline status of a unit. Unit commitment restrictions will govern how this variable can change through time.
- [units\\_on\\_\\_temporal\\_block](#) is a relationship linking the `units_on` variable of this unit to a specific [temporal\\_block](#) object. The temporal block holds information on the temporal scope and resolution for which the variable should be optimized.
- [online\\_variable\\_type](#) is a method parameter and can take the values `unit_online_variable_type_binary`, `unit_online_variable_type_integer`, `unit_online_variable_type_linear`. If the binary value is chosen, the units status is modelled as a binary (classic UC). For clustered unit commitment units, the integer type is applicable. Note that if the parameter is not defined, the default will be linear. If the units status is not crucial, this can reduce the computational burden.
- [number\\_of\\_units](#) defines how many units of a certain unit type are available. Typically this parameter takes a binary (UC) or integer (clustered UC) value. To avoid confusion the following distinction will be made in this document: `unit` will be used to identify a Spine unit object, which can have multiple `members`. Together with the `unit_availability_factor`, this will determine the maximum number of members that can be online at any given time. (Thus restricting the `units_on` variable). The default value for this parameter is 1. It is possible to allow the model to increase the `number_of_units` itself, through [Investment Optimization](#)
- [unit\\_availability\\_factor](#): (number value or time series). Is the fraction of the time that this unit is considered to be available, by acting as a multiplier on the capacity. A time series can be used to indicate the intermittent character of renewable generation technologies.
- [min\\_up\\_time](#): (duration value). Sets the minimum time that a unit has to stay online after a startup. Inclusion of this parameter will trigger the creation of the constraint on [Minimum up time \(basic version\)](#)

- `min_down_time`: (duration value). Sets the minimum time that a unit has to stay offline after a shutdown. Inclusion of this parameter will trigger the creation of the constraint on [Minimum down time \(basic version\)](#)
- `minimum_operating_point`: (number value) limits the minimum value of the `unit_flow` variable for a unit which is currently online. Inclusion of this parameter will trigger the creation of the [Constraint on minimum operating point](#)
- `start_up_cost`: "number value". Cost associated with starting up a unit.
- `shut_down_cost`: "number value". Cost associated with shutting down a unit.

## Illustrative unit commitment examples

### Step 1: defining the number of members of a unit type

A spine unit can represent multiple members. This can be incorporated in a model by setting the `number_of_units` parameter to a specific value. For example, if we define a single unit in a model as follows:

- `unit_1`
  - `number_of_units: 2`

And we link the unit to a certain `node_1` with a `unit_to_node` relationship.

- `unit_1_to__node_1`

The single spine unit defined here, now represents two members. This means that a single `unit_flow` variable will be created for this unit, but the restrictions as imposed by the [Ramping and Reserves](#) framework will be adapted to reflect the fact that there are two members present, thus doubling the total capacity.

### Step 2: choosing the online\_variable\_type

Next, we have to decide the `online_variable_type` for this unit, which will restrict the kind of values that the `units_on` variable can take. This basically comes down to deciding if we are working in a classical UC framework (`unit_online_variable_type_binary`), a clustered UC framework (`unit_online_variable_type_integer`), or a relaxed clustered UC framework (`unit_online_variable_type_linear`), in which a non-integer number of units can be online.

The classical UC framework can only be applied when the `number_of_units` equals 1.

### Step 3: imposing a minimum operating point

The output of an online unit to a specific node can be restricted to be above a certain minimum by choosing a value for the `minimum_operating_point` parameter. This parameter is defined for the `unit_to_node` relationship, and is given as a fraction of the `unit_capacity`. If we continue with the example above, and define the following objects, relationships, and parameters:

- `unit_1`
  - `number_of_units`: 2
  - `unit_online_variable_type`: "unit\_online\_variable\_type\_integer"
- `unit_1_to_node_1`
  - `minimum_operating_point`: 0.2
  - `unit_capacity`: 200

It can be seen that in this case the `unit_flow` from `unit_1` to `node_1` must for any timestep  $t$  be larger than  $units\_on(t) * 0.2 * 200$

## Step 4: imposing a minimum up or down time

Spine units can also be restricted in their commitment status with minimum up- or down times by choosing a value for the `min_up_time` or `min_down_time` respectively. These parameters are defined for the `unit` object, and should be duration values. We can continue the example and add a minimum up time for the unit:

- `unit_1`
  - `number_of_units`: 2
  - `unit_online_variable_type`: "unit\_online\_variable\_type\_integer"
  - `min_up_time`: 2h
- `unit_1_to_node_1`
  - `minimum_operating_point`: 0.2
  - `unit_capacity`: 200

Whereas the `units_on` variable was restricted (before inclusion of the `min_up_time` parameter) to be smaller than or equal to the `number_of_units` for any timestep  $t$ , it now has to be smaller than or equal to the `number_of_units` decremented with the `units_started_up` summed over the timesteps that include  $t - min\_up\_time$ . This implies that a unit which has started up, has to stay online for at least the `min_up_time`

To consider a simple example let's assume that we have a model with a resolution of 1h. Suppose that before  $t$ , there is no member of the unit online and in timestep  $t \rightarrow t + 1h$ , one member starts up. Another member starts up in timestep  $t + 1h \rightarrow t + 2h$ . The first startup, along with the minimum up time of 2 hours implies that the `units_on` variable of this unit has now changed to 1 in timestep  $t - > t + 1h$  and can not go back to 0 in timestep  $t \rightarrow t + 1h \rightarrow t + 2h$ . The second startup further

restricts the number of units that are allowed to be online, it can be seen that the following restrictions apply when both startups are combined with the minimum up time of 2h:

- $t \rightarrow t + 1h: units\_on = 1$
- $t + 1h \rightarrow t + 2h: units\_on = 2$
- $t + 2h \rightarrow t + 3h: units\_on \in 1, 2$
- $t + 3h \rightarrow t + 4h: units\_on \in 0, 1, 2$

The minimum down time restrictions operate in very much the same way, they simply impose that units that have been shut down, have to stay offline for the chosen period of time.

## Step 5: allocationg a cost to startups or shutdowns

Costs can be allocated to startups or shutdowns by choosing a value for the [start\\_up\\_cost](#) or [shut\\_down\\_cost](#) respectively.

## Step 6: defining unit availabilities

By defining a [unit\\_availability\\_factor](#), the fact that typical members are not available all the time can be reflected in the model.

Typically, units are not available 100% of the time, due to scheduled maintenance, unforeseen outages, or other things. This can be incorporated in the model by setting the [unit\\_availability\\_factor](#) to a fractional value. For each timestep in the model, an upper bound is then imposed on the [units\\_on](#) variable, equal to [number\\_of\\_units](#) \* [unit\\_availability\\_factor](#). This parameter can not be used when the [online\\_variable\\_type](#) is binary. It should also be noted that when the [online\\_variable\\_type](#) is of integer type, the aforementioned product must be integer as well, since it will determine the value of the [units\\_available](#) parameter which is restricted to integer values. The default value for this parameter is 1.

The [unit\\_availability\\_factor](#) can also be taken as a timeseries. By allowing a different availability factor for each timestep in the model, it can perfectly be used to represent intermittent technologies of which the output cannot be fully controlled.

# Ramping and Reserves

To enable the representation of units with a high level of technical detail, the ramping ability of units can be constrained in SpineOpt. This means that the user has the freedom to impose restrictions on the change in output of units between consecutive timesteps, for online (spinning) units, units starting up and units shutting down. In this section, the concept of ramps in SpineOpt will be introduced. Furthermore, the use of reserves will be explained.

## Relevant objects, relationships and parameters

Everything that is related to ramping is defined in parameters of either the [unit\\_to\\_node](#), [unit\\_from\\_node](#), or [unit\\_to\\_node\\_group](#) relationship. Generally speaking, the ramping constraints will impose restrictions on the change in the [unit\\_flow](#) variable between two consecutive timesteps.

All parameters that limit the ramping abilities of a unit are expressed as a fraction of the unit capacity. This means that a value of 1 indicates the full capacity of a unit.

The discussion here will be kept conceptual, for the mathematical formulation the reader is referred to the [Ramping and reserve constraints](#)

## Constraining spinning ramps

- [unit\\_capacity](#): limit the maximum value of the `unit_flow` variable for a unit which is currently online. Inclusion of this parameter will trigger the creation of the [Define unit/technology capacity](#) constraint.
- [ramp\\_up\\_limit](#): limit the maximum increase in the `unit_flow` variable between two consecutive timesteps for which the unit is online. The parameter is given as a fraction of the [unit\\_capacity](#) parameter. Inclusion of this parameter will trigger the creation of the [Constraint on spinning upwards ramp\\_up](#)
- [ramp\\_down\\_limit](#): limit the maximum decrease in the `unit_flow` variable between two consecutive timesteps for which the unit is online. The parameter is given as a fraction of the [unit\\_capacity](#) parameter. Inclusion of this parameter will trigger the creation of the [Constraint on spinning downward ramps](#)
- [ramp\\_up\\_cost](#): cost associated with upward ramping
- [ramp\\_down\\_cost](#): cost associated with downward ramping

## Constraining shutdown ramps

- `max_shutdown_ramp` : limit the maximum of the `unit_flow` variable for the timestep right before a shutdown. The parameter is given as a fraction of the `unit_capacity` parameter. Inclusion of this parameter will trigger the creation of the [Constraint on maximum downward shut down ramps](#)
- `min_shutdown_ramp` : limit the minimum of the `unit_flow` variable for the timestep right before a shutdown. The parameter is given as a fraction of the `unit_capacity` parameter. Inclusion of this parameter will trigger the creation of the [Constraint on minimum downward shut down ramps](#)

## Constraining startup ramps

- `max_startup_ramp` : limit the maximum of the `unit_flow` variable for the timestep right after a start-up. The parameter is given as a fraction of the `unit_capacity` parameter. Inclusion of this parameter will trigger the creation of the [Constraint on maximum upward start up ramp\\_up](#)
- `min_startup_ramp` : limit the minimum of the `unit_flow` variable for the timestep right after a start-up. The parameter is given as a fraction of the `unit_capacity` parameter. Inclusion of this parameter will trigger the creation of the [Constraint on minimum upward start up ramp\\_up](#)

## General principle and example use cases

The general principle of the Spine modelling ramping constraints is that all of these parameters can be defined separately for each unit. This allows the user to incorporate different units (which can either represent a single unit or a technology type) with different flexibility characteristics.

It should be noted that it is perfectly possible to omit all of the constraining parameters mentioned above. However, once either of the ramping parameters is defined, it is necessary to also assign values to the other parameters. E.g. if a user only wants to restrict the spinning ramp up capability of a unit, one also has to assign values to the `max_startup_ramp`, `min_Shutdown_Ramp` etc.

## Illustrative examples

### Step 1: Simple case of unrestricted unit

When none of the ramping parameters mentioned above are defined, the unit is considered to have full ramping flexibility. This means that in any given timestep, its output can be any value between 0 and its capacity, regardless of what the output of the unit was in the previous timestep, and regardless of the on- or offline status or the unit in the previous timestep. Provided that this does not conflict with the [Unit commitment](#) restrictions that are defined for this unit. Parameter values for a `unit__node` relationship are illustratively given below.

- `max_shutdown_ramp` : 1
- `min_shutdown_ramp` : 0
- `max_start_up_ramp` : 1
- `min_start_up_ramp` : 0
- `ramp_up_limit` : 1



- `ramp_down_limit` : 1
- `unit_capacity` : 200

## Step 2: Spinning ramp restriction

A unit which is only restricted in spinning ramping can be created by changing the `ramp_up/down_limit` parameters:

- `ramp_up_limit` : 0.2
- `ramp_down_limit` : 0.4

This parameter choice implies that the unit's output between two consecutive timesteps can change with no more than  $0.2 * 200$  and no less than  $0.4 * 200$ . For example, when the unit is running at an output of 100 in some timestep  $t$ , its output for the next timestep must be somewhere in the interval  $[20, 140]$ . Unless it shuts down completely.

## Step 3: Shutdown restrictions

By changing the parameter `max_shutdown_ramp` in the previous example, an additional restriction is imposed on the maximum output of the unit from which it can go offline.

- `max_shutdown_ramp` : 0.5
- `min_shutdown_ramp` : 0.3

When this unit goes offline in a given timestep  $t + 1$ , the output of the unit must be below  $0.5 * 200 = 100$  in the timestep  $t$  before that. Similarly, the parameter `min_shutdown_ramp` can be used to impose a minimum output value in the timestep before a shutdown. For example, a value of 0.3 in this example would mean that the unit can not be running below an output of 60 in timestep  $t$ .

## Step 4: Startup restrictions

The startup restrictions are very similar to the shutdown restrictions, but of course apply to units that are starting up. Consider for example the same unit as in the example above, but now with a `max_start_up_ramp` equal to 0.4 and `min_start_up_ramp` equal to 0.2:

- `max_start_up_ramp` : 0.4
- `min_start_up_ramp` : 0.2

When the unit is offline in timestep  $t$  and comes online in timestep  $t + 1$ , its output in timestep  $t + 1$  will be restricted to the interval  $[40, 80]$ .

# Reserve concept

To include a requirement of reserve provision in a model, SpineOpt offers the possibility of creating reserve nodes. Of course reserve provision is different from regular operation, because the reserved

capacity does not actually get activated. In this section, we will take a look at the things that are particular for a reserve node.

## Defining a reserve node

To define a reserve node, the following parameters have to be defined for the relevant node:

- `is_reserve_node` : this boolean parameter indicates that this node is a reserve node.
- `upward_reserve` : this boolean parameter indicates that the demand for reserve provision of this node concerns upward reserves.
- `downward_reserve` : this boolean parameter indicates that the demand for reserve provision of this node concerns downward reserves.
- `reserve_procurement_cost`: (optional) this parameter indicates the procurement cost of a unit for a certain reserve product and can be define on a `unit__to_node` or `unit__from_node` relationship.

## Defining a node group

SpineOpt allows the user to constrain ramping abilities of units that are linked to multiple nodes by defining node groups. This is especially relevant for reserve provision because a unit that provides reserves is linked to a regular, as well as a reserve node. It is then possible to constrain the unit's ramping for the combination of regular operation and reserve provision.

Since reserve provision in fact literally reserves part of the capacity of a unit, the demand of the reserve node will be subtracted from the part that is available for regular operation. The section below will discuss how this works in SpineOpt by means of an example.

Since the demand of the nodes is defined on the individual node level (and the node group has no demand), the balance type of the group node should be set to `balance_type_none`.

## Ramping constraints on a node group with one reserve node

### Reserves step 1: simple case of unrestricted unit

Let's assume that we have one unit and two nodes in a model, one for reserves and one for regular demand. The unit is then linked by the `unit__to_node` relationships to each node individually, and on top of that, it is linked to a node group containing both nodes.

The ramping of the unit can now be constrained by defining the same parameters as before, but now for the node group. As before, the simplest case is a unit that is only restricted by its capacity:

- `max_shutdown_ramp` : 1

- `min_shutdown_ramp` : 0
- `max_start_up_ramp` : 1
- `min_start_up_ramp` : 0
- `ramp_up_limit` : 1
- `ramp_down_limit` : 1
- `unit_capacity` : 200

The capacity restriction now implies that the sum of the reserve demand and regular demand cannot exceed the capacity of the unit. For example: when the reserve node has a demand of 10 in timestep  $t$ , the `unit_flow` variable to the regular node must be smaller than or equal to 190.

## Reserves step 2: Spinning ramp restriction

The unit can be restricted only in spinning ramping, as in the previous example, by defining the `ramp_up/down_limit` parameters in the `unit__to_node` relationship **for the node group**:

- `ramp_up_limit` : 0.2
- `ramp_down_limit` : 0.4

This parameter choice implies that the unit's flow to the regular demand node between two consecutive timesteps can change with no more than  $0.2 * 200 - \text{upward\_reserve\_demand}$  and no less than  $0.4 * 200 - \text{downward\_reserve\_demand}$ . For example, when the unit is running at an output of 100 in some timestep  $t$ , and there is an upward reserve demand of 10 its output for the next timestep must be somewhere in the interval  $[20, 130]$ .

It can be seen in this example that the demand for reserves is subtracted from both the generation capacity, and the ramping capacity of the unit that is available for regular operation. This stems from the fact that in providing reserve capacity, the unit is expected to be able to provide the demanded reserve within one timestep.

## Reserves Step 3: Non-spinning reserves

Units can also be allowed to provide non-spinning reserves, through shutdowns and startups. This can be done by using the following parameters in the `unit__to_node` relationship **for the reserve node** :

- `max_res_startup_ramp`
- `min_res_startup_ramp`
- `max_res_shutdown_ramp`
- `min_res_shutdown_ramp`
- `unit_capacity`

These parameters are constraining reserve provision in exactly the same way as their equivalents for regular operation. Note that it is now necessary to define a capacity of the unit with respect to the

reserve node. The ramping parameters will then be interpreted as fractions of this specific capacity. The unit's overall capacity can be different than its capacity for reserve provision.

A unit which can provide both spinning and non-spinning reserves can be defined as follows:

#### Parameters to be defined for unit to node group relationship

- `max_shutdown_ramp` : 1
- `min_shutdown_ramp` : 0
- `max_start_up_ramp` : 1
- `min_start_up_ramp` : 0
- `ramp_up_limit` : 0.2
- `ramp_down_limit` : 0.4
- `unit_capacity` : 200

#### Parameters to be defined for unit to reserve node relationship

- `max_res_startup_ramp` : 0.5
- `min_res_startup_ramp` : 0.1
- `unit_capacity` : 150

The spinning reserve and ramping restrictions now remain the same as above, but on top of that the unit is able to provide non-spinning upward reserves when it is offline. In this particular example, the contribution of the offline unit to upward reserves can be anything in the interval `[15, 75]`.

## Using node\_groups for both combined and individual restrictions

It can be seen from the example above that when a node group is defined, ramping restrictions can be imposed both on the group level (thus for the unit as a whole) as well as for the individual nodes. If, for example a `ramp-up-limit` is defined for the node group, the sum of upward ramping of the two nodes will be restricted by this parameter, but it is still possible to limit the individual flows to the nodes as well. We will now discuss an example of this for the `ramp_up_limit`, but this also holds for other parameters.

Let's continue with the example above, where an online unit is capable of ramping up by 20% of its capacity and down by 40%. We might want to impose tighter restrictions for upward reserve provision than the ramping in overall operation (e.g. because the reserved capacity has to be available in a shorter time than the [duration\\_unit](#)). One can then simply define an additional parameter for the unit to reserve node relationship as follows.

- `ramp_up_limit` : 0.15

Which now restricts the spinning upward ramping provision of the unit to 15% of its capacity, as defined for the reserve node. In this case, the change in the unit's flow to the regular demand node between two consecutive timesteps is still limited to the interval  $[0.2 * 200 - \text{upward\_reserve\_demand}, 0.4 * 200 - \text{downward\_reserve\_demand}]$ . But the upward reserves that it can provide has an upper bound of  $150 * 0.15$ .

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# Investment Optimization

SpineOpt offers numerous ways to optimise investment decisions energy system models and in particular, offers a number of methodologies for capturing increased detail in investment models while containing the impact on run time. The basic principles of investments will be discussed first and this will be followed by more advanced approaches.

## Key concepts for investments

**Investment Decisions** These are the investment decisions that SpineOpt currently supports. At a high level, this means that the activity of the entities in question is controlled by an investment decision variable. The current implementation supports investments in :

- **unit:**
- **connection**
- **Storage** - Note: while the above investment decisions correspond to an object class (i.e.) an investment in a **unit** or a **connection**, **Storages** are not an object class in themselves and are rather a property of a **node**. As such, a storage investment controls whether a particular node has a state variable or not.

**Investment Variable Types** In all cases the capacity of the **unit** or **connection** or the maximum node state of a **node** is multiplied by the investment variable which may either be continuous, binary or integer. This is determined, for units, by setting the **unit\_investment\_variable\_type** parameter accordingly. Similarly, for connections and node storages where the **connection\_investment\_variable\_type** and **storage\_investment\_variable\_type** are specified.

**Identifying Investment Candidate Units, Connections and Storages** The parameter **candidate\_units** represents the number of units of this type that may be invested-in. **candidate\_units** determines the upper bound of the investment variable and setting this to a value greater than 0 identifies the unit as an investment candidate unit in the optimisation. If the **unit\_investment\_variable\_type** is set to `:variable_type_integer`, the investment variable can be interpreted as the number of discrete units that may be invested in. However, if **unit\_investment\_variable\_type** is `:variable_type_continuous` and the **unit\_capacity** is set to unity, the investment decision variable can then be interpreted as the capacity of the unit rather than the number of units with **candidate\_units** being the maximum capacity that can be invested in. Finally, we can invest in discrete blocks of capacity by setting **unit\_capacity** to the size of the investment capacity blocks and have **unit\_investment\_variable\_type** set to `:variable_type_integer` with **candidate\_units** representing the maximum number of capacity blocks that may be invested in. The key points here are:

- The upper bound on the relevant flow variables are determined by the product of the investment variable and the [unit\\_capacity](#) or [connection\\_capacity](#) for connections or [node\\_state\\_cap](#) for storages.
- [candidate\\_units](#) sets the upper bound on the investment variable, [candidate\\_connections](#) for connections and [candidate\\_storages](#) for storages
- [unit\\_investment\\_variable\\_type](#) determines whether the investment variable is integer, binary or continuous ([connection\\_investment\\_variable\\_type](#) for connections and [storage\\_investment\\_variable\\_type](#) for storages).

**Investment Costs** Investment costs are specified by setting the appropriate `*_investment\_cost` parameter. The investment cost for [units](#) are specified by setting the [unit\\_investment\\_cost](#) parameter. This is currently interpreted as the full cost over the investment period for the unit. See the section below on **investment temporal structure** for setting the investment period. If the investment period is 1 year, then the corresponding [unit\\_investment\\_cost](#) is the annualised investment cost. For connections and storages, the investment cost parameters are [connection\\_investment\\_cost](#) and [storage\\_investment\\_cost](#), respectively.

**Temporal and Stochastic Structure of Investment Decisions** SpineOpt's flexible stochastic and temporal structure extend to investments where individual investment decisions can have their own temporal and stochastic structure independent of other investment decisions and other model variables. A global temporal resolution for all investment decisions can be defined by specifying the relationship [model\\_default\\_investment\\_temporal\\_block](#). If a specific temporal resolution is required for specific investment decisions, then one can specify the following relationships: - [unit\\_investment\\_temporal\\_block](#) for [unit](#) - [connection\\_investment\\_temporal\\_block](#) for [connection](#) - [node\\_investment\\_temporal\\_block](#) for storages. Specifying any of the above relationships will override the corresponding [model\\_default\\_investment\\_temporal\\_block](#).

Similarly, a global stochastic structure can be defined for all investment decisions by specifying the relationship [model\\_default\\_investment\\_stochastic\\_structure](#). If a specific stochastic structure is required for specific investment decisions, then one can specify the following relationships: - [unit\\_investment\\_stochastic\\_structure](#) for [unit](#) - [connection\\_investment\\_stochastic\\_structure](#) for [connection](#) - [node\\_investment\\_stochastic\\_structure](#) for storages Specifying any of the above relationships will override the corresponding [model\\_default\\_investment\\_stochastic\\_structure](#).

## Creating an Investment Candidate Unit Example

If we have a model that is not currently set up for investments and we wish to create an investment candidate unit, we can take the following steps.

- Create the unit object with all the relationships and parameters necessary to describe its function.
- Ensure that the [number\\_of\\_units](#) parameter is set to zero so that the unit is unavailable unless invested-in

- Set the `candidate_units` parameter for the unit to 1 to specify that a maximum of 1 new unit of this type may be invested-in by the model.
- Set the `unit_investment_variable_type` to `unit_investment_variable_type_integer` to specify that this is a discrete `unit` investment decision.
- Specify the `unit_investment_lifetime` of the unit to, say, 1 year to specify that this is the minimum amount of time this new unit must be in existence after being invested-in.
- Specify the investment period for this `unit`'s investment decision in one of two ways
  - Define a default investment period for all investment decisions in the model as follows:
    - create a `temporal_block` with the appropriate `resolution` (say 1 year)
    - link this to your `[model]{@ref}` object by creating the appropriate `model_temporal_block` relationship
    - set it as the default investment temporal block by setting `model_default_investment_temporal_block`
  - Or, define an investment period unique to this investment decision as follows:
    - creating a `temporal_block` with the appropriate `resolution` (say 1 year)
    - link this to your model object by creating the appropriate `model_temporal_block` relationship
    - specify this as the investment period for your `unit`'s investment decision by setting the appropriate `unit_investment_temporal_block` relationship
- Similarly to the above, define the stochastic structure for the `unit`'s investment decision by specifying either `model_default_investment_stochastic_structure` or `unit_investment_stochastic_structure`
- Specifying your `unit`'s investment cost by setting the `unit_investment_cost` parameter. Since we have defined the investment period above as 1 year, this is therefore the `unit`'s annualised investment cost.

## Model Reference

### Variables for investments

Variable Name	Indices	Description
<code>units_invested_available</code>	<code>unit, s, t</code>	The number of invested in <code>units</code> that are available at a given (s, t)
<code>units_invested</code>	<code>unit, s, t</code>	The point-in-time investment decision corresponding to the number of <code>units</code> invested in at (s,t)
<code>units_mothballed</code>	<code>unit, s, t</code>	"Instantaneous" decision variable to



		mothball a unit
connections_invested_available	connection, s, t	The number of invested-in connectionss that are available at a given (s, t)
connections_invested	connection, s, t	The point-in-time investment decision corresponding to the number of connectionss invested in at (s,t)
connections_decommissioned	connection, s, t	"Instantaneous" decision variable to decommission a connection
storages_invested_available	node, s, t	The number of invested-in storages that are available at a given (s, t)
storages_invested	node, s, t	The point-in-time investment decision corresponding to the number of storages invested in at (s,t)
storages_decommissioned	node, s, t	"instantaneous" decision variable to decommission a storage

## Relationships for investments

Relationship Name	Related Object Class List	
model__default_investment_temporal_block	model, temporal_block	Default te unit_investm
model__default_investment_stochastic_structure	model, stochastic_structure	Default st unit_inve
unit__investment_temporal_block	unit, temporal_block	Set t model_c
unit__investment_stochastic_structure	unit, stochastic_structure	Set st model_defau

# Parameters for investments

Parameter Name	Object Class List	Description
candidate_units	unit	The number of additional units of this type that can be invested in
unit_investment_cost	unit	The total overnight investment cost per candidate unit over the model horizon
unit_investment_lifetime	unit	The investment lifetime of the unit - once invested-in, a unit must exist for at least this amount of time
unit_investment_variable_type	unit	Whether the units_invested_available variable is continuous, integer or binary
fix_units_invested	unit	Fix the value of units_invested
fix_units_invested_available	unit	Fix the value of connections_invested_available
candidate_connections	connection	The number of additional connections of this type that can be invested in
connection_investment_cost	connection	The total overnight investment cost per candidate connection over the model horizon
connection_investment_lifetime	connection	The investment lifetime of the connection - once invested-in, a connection must exist for at least this amount of time
connection_investment_variable_type	connection	Whether the connections_invested_available variable is continuous, integer or binary
fix_connections_invested	connection	Fix the value of connections_invested

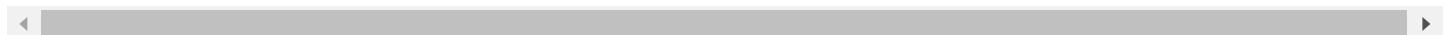
<code>fix_connections_invested_available</code>	<code>connection</code>	Fix the value of <code>connection_invested_available</code>
<code>candidate_storages</code>	<code>node</code>	The number of additional storages of this type that can be invested in at <code>node</code>
<code>storage_investment_cost</code>	<code>node</code>	The total overnight investment cost per candidate storage over the model horizon
<code>storage_investment_lifetime</code>	<code>node</code>	The investment lifetime of the storage - once invested-in, a storage must exist for at least this amount of time
<code>storage_investment_variable_type</code>	<code>node</code>	Whether the <code>storages_invested_available</code> variable is continuous, integer or binary
<code>fix_storages_invested</code>	<code>node</code>	Fix the value of <code>storages_invested</code>
<code>fix_storages_invested_available</code>	<code>node</code>	Fix the value of <code>storages_invested_available</code>

## Related Model Files

Filename	Relative Path	Description
<code>constraintunitsinvested_available.jl</code>	<code>\constraints</code>	constrains <code>units_invested_available</code> to be less than <code>candidate_units</code>
<code>constraintunitsinvested_transition.jl</code>	<code>\constraints</code>	defines the relationship between <code>units_invested_available</code> , <code>units_invested</code> and <code>units_mothballed</code> . Analogous to <code>units_on</code> , <code>units_started</code> and <code>units_shutdown</code>
<code>constraintunitlifetime.jl</code>	<code>\constraints</code>	once a <code>unit</code> is invested-in, it must remain in existence for at least <code>unit_investment_lifetime</code> -

analogous to `min_up_time`.

<code>constraintunitsavailable.jl</code>	<code>\constraints</code>	Enforces <code>units_available</code> is the sum of <code>number_of_units</code> and <code>units_invested_available</code>
<code>constraintconnectionsinvested_available.jl</code>	<code>\constraints</code>	constrains <code>connections_invested_available</code> to be less than <code>candidate_connections</code>
<code>constraintconnectionsinvested_transition.jl</code>	<code>\constraints</code>	defines the relationship between <code>connections_invested_available</code> , <code>connections_invested</code> and <code>connections_decommissioned</code> . Analogous to <code>units_on</code> , <code>units_started</code> and <code>units_shutdown</code>
<code>constraintconnectionlifetime.jl</code>	<code>\constraints</code>	once a <code>connection</code> is invested-in, it must remain in existence for at least <code>connection_investment_lifetime</code> - analogous to <code>min_up_time</code> .
<code>constraintstoragesinvested_available.jl</code>	<code>\constraints</code>	constrains <code>storages_invested_available</code> to be less than <code>candidate_storages</code>
<code>constraintstoragesinvested_transition.jl</code>	<code>\constraints</code>	defines the relationship between <code>storages_invested_available</code> , <code>storages_invested</code> and <code>storages_decommissioned</code> . Analogous to <code>units_on</code> , <code>units_started</code> and <code>units_shutdown</code>
<code>constraintstoragelifetime.jl</code>	<code>\constraints</code>	once a <code>storage</code> is invested-in, it must remain in existence for at least <code>storage_investment_lifetime</code> - analogous to <code>min_up_time</code> .



# Unit Constraints

Unit constraints allow the user to define arbitrary linear constraints involving most of the problem variables. This section describes this function and how to use it.

## Key Unit Constraint Concepts

1. **The basic principle:** The basic steps involved in forming a unit constraint are:

- Creating a user constraint object: One creates a new `unit_constraint` object which will be used as a unique handle for the specific constraint and on which constraint-level parameters will be defined.
- Specify which variables are involved in the constraint: this generally involves creating a relationship involving the `unit_constraint` object. For example, specifying the relationship `unit_from_node_unit_constraint` specifies that the corresponding `unit_flow` variable is involved in the constraint. The table below contains a complete list of variables and the corresponding relationships to set.
- Specify the variable coefficients: this will generally involve specifying a parameter named `*_coefficient` on the relationship defined above to specify the coefficient on that particular variable in the constraint. For example, to define the coefficient on the `unit_flow` variable, one specifies the `unit_flow_coefficient` parameter on the appropriate `unit_from_node_unit_constraint` relationship. The table below contains a complete list of variables and the corresponding coefficient parameters to set.
- Specify the right-hand-side constant term: The constraint should be formed in conventional form with all constant terms moved to the right-hand side. The right-hand-side constant term is specified by setting the `right_hand_side_unit_constraint` parameter.
  - Specify the constraint sense: this is done by setting the `constraint_sense_unit_constraint` parameter. The allowed values are `==`, `>=` and `<=`.
  - Coefficients can be defined on some parameters themselves. For example, one may specify a coefficient on a node's demand parameter. This is done by specifying the relationship `node_unit_constraint` and specifying the `demand_coefficient` parameter on that relationship

1. **Piecewise unit\_flow coefficients:** As described in [operating\\_points](#), specifying this parameter decomposes the `unit_flow` variable into a number of sub operating segment variables named `unit_flow_op` in the model and with an additional index, `i` for the operating segment. The intention of this functionality is to allow `unit_flow` coefficients to be defined individually per segment to define a piecewise linear function. To accomplish this, the steps are as described above with the exception that one must define `operating_points` on the appropriate `unit_from_node` or `unit_to_node` as an array type with the dimension corresponding to the number of operating points and then set the `unit_flow_coefficient` for the appropriate `unit_from_node_unit_constraint`

relationship, also as an array type with the same number of elements. Note that if operating points is defined as an array type with more than one elements, [unit\\_flow\\_coefficient](#) may be defined as either an array or non-array type. However, if [operating\\_points](#) is of non-array type, corresponding [unit\\_flow\\_coefficients](#) must also be of non-array types.

**2. Variables, relationships and coefficient guide for unit constraints** The table below provides guidance regarding what relationships and coefficients to set for various problem variables and parameters.

Problem variable / Parameter Name	Relationship	Parameter
<code>unit_flow</code> (direction=from_node)	<a href="#">unit_from_node_unit_constraint</a>	<a href="#">unit_flow_coefficient</a> (non-array type)
<code>unit_flow</code> (direction=to_node)	<a href="#">unit_to_node_unit_constraint</a>	<a href="#">unit_flow_coefficient</a> (non-array type)
<code>unit_flow_op</code> (direction=from_node)	<a href="#">unit_from_node_unit_constraint</a>	<a href="#">unit_flow_coefficient</a> (array type)
<code>unit_flow_op</code> (direction=to_node)	<a href="#">unit_to_node_unit_constraint</a>	<a href="#">unit_flow_coefficient</a> (array type)
<code>connection_flow</code> (direction=from_node)	<a href="#">connection_from_node_unit_constraint</a>	<a href="#">connection_flow_coefficient</a>
<code>connection_flow</code> (direction=to_node)	<a href="#">connection_to_node_unit_constraint</a>	<a href="#">connection_flow_coefficient</a>
<code>node_state</code>	<a href="#">node_unit_constraint</a>	<a href="#">node_state_coefficient</a>
<code>demand</code>	<a href="#">node_unit_constraint</a>	<a href="#">demand_coefficient</a>

# Decomposition

Decomposition approaches take advantage of certain problem structures to separate them into multiple related problems which are each more easily solved. Decomposition also allows us to do the inverse, which is to combine independent problems into a single problem, where each can be solved separately but with communication between them (e.g. investments and operations problems)

Decomposition thus allows us to do a number of things

- Solve larger problems which are otherwise intractable
- Include more detail in problems which otherwise need to be simplified
- Combine related problems (e.g. investments/operations) in a more scientific way (rather than ad-hoc).
- Employ parallel computing methods to solve multiple problems simultaneously.

## High-level Decomposition Algorithm

The high-level algorithm is described below. For a more detailed description please see [Benders decomposition](#)

- Model initialisation (preprocess *data* structure, generate temporal structures etc.)
- For each benders\_iteration
  - Solve master problem
  - Process master-problem solution:
    - set `units_invested_bi(unit=u, benders_iteration=bi)` equal to a timeseries representing the investment variables solution from the master problem
  - Rewind and update operations problem
  - Solve operations problem loop
  - Process operations sub-problem
    - set `units_available_mv(unit=u, benders_iteration=bi)` equal to a timeseries representing the marginal value of the `units_on bound` constraint
  - Test for convergence
  - Update master problem
    - Add [Benders cuts constraints](#)
  - Next benders iteration

## Duals calculation for decomposition

The `optimize_model!()` function has been updated to optionally include an additional step for the calculation of duals. The dual solution to a MIP problem is not well defined. The standard approach to obtaining marginal values from a MIP model is to relax the integer variables, fix them to their last solution value and re-solve the problem as an LP. This is the standard approach in energy system modelling to obtain energy prices. However, although this is the standard approach, it does need to be used with caution (see here for example). The main hazard associated with inferring duals in this way is that the impact on costs of an investment may be overstated. However, since these duals are used in Benders decomposition to obtain a lower bound on costs (i.e. the maximum potential value from an investment), this is ok and can be "corrected" in the next iteration. And finally, the benders gap will tell us how close our decomposed problem is to the optimal global solution.

This additional relaxed LP solve is done as follows:

- `add_variable!()` stores the list of integer and binary variables in `m.ext[:integer_variables]`
- the `fix_value` for integer variables is set to the last MIP solution value
- the integer constraints on the integer variables are `unset()`
- A final LP is solved
- required dual values are saved
- integer constraints on integer variables are `set()`

This final fixed LP solve is triggered by specifying `calculate_duals=true` in the call to `optimize_model!()`

## Reporting dual values:

To report the dual of a constraint, one can add an output item with the corresponding constraint name (e.g. `constraint_nodal_balance`) and add that to a report. This will cause the corresponding constraint's relaxed problem marginal value will be reported in the output DB. When adding a constraint name as an output we need to preface the actual constraint name with `constraint_` to avoid ambiguity with variable names (e.g. `units_available`). So to report the marginal value of `units_available` we add an output object called `constraint_units_available`.

To report the `reduced_cost()` for a variable which is the marginal value of the associated active bound or fix constraints on that variable, one can add an output object with the variable name prepended by `bound_`. So, to report the unit *on reduced cost* value, one would create an output item called `bound_units_on`. If added to a report, this will cause the reduced cost of unit *on in the final fixed LP to be written to the output db*. Finally, if any constraint duals or reduced cost values are requested via a report, `calculate_duals` is set to true and the final fixed LP solve is triggered.

## Using Decomposition

The decomposition framework creates a master problem where the investment variables are optimised. The decomposition framework is invoked when a model object with the parameter `model_type` set to



:spineopt\_operations is found and a second model object with model\_type set to :spineopt\_master. Once these conditions are met, all investment decisions in the model are automatically decomposed and optimised in the master problem. This behaviour may change in the future to allow some investment decisions to be optimised in the operations problem and some optimised in the master problem as desired.

**Steps to involve decomposition in an investments problem** Assuming one has set up a conventional investments problem as described in [Investment Optimization](#) the following additional steps are required to utilise the decomposition framework:

- Create a new [model](#) object to represent the benders master problem
- Set the [model\\_type](#) parameter for the master problem model to spineopt\_master.
- Set the [model\\_type](#) parameter for the existing conventional, operations problem model to spineopt\_operations.
- Specify the master problem [model](#) parameter, max\_gap - This determines the master problem convergence criterion for the relative benders gap. A value of 0.05 will represent a relative benders gap of 5%.
- Specify the master problem [model](#) parameter max\_iterations - This determines the master problem convergence criterion for the number of iterations. A value of 10 could be appropriate but this is highly dependent on the size and nature of the problem
- Specify appropriate [model\\_report](#) relationships to determine which reports are written for which model

# Power transfer distribution factors (PTDF) based DC power flow

There are two main methodologies for directly including DC powerflow in unit commitment/energy system models. One method is to directly include the bus voltage angles as variables in the model. This method is described in [Nodal lossless DC Powerflow](#).

Here we discuss the method of using power transfer distribution factors (PTDF) for DC power flow and line outage distribution factors (lodf) for security constrained unit commitment.

## Key concepts

1. **ptdf**: The power transfer distribution factors are a property of the network reactances and their derivation may be found [here](#). `ptdf(n, c)` represents the fraction of an injection at [node](#) `n` that will flow on [connection](#) `c`. The flow on [connection](#) `c` is then the sum over all nodes of `ptdf(n, c)*net_injection(c)`. The advantage of this method is that it introduces no additional variables into the problem and instead, introduces only one constraint for each connection whose flow we are interested in monitoring.
2. **lodf**: Line outage distribution factors are a function of the network ptdfs and their derivation is also found [here](#). `lodf(c_contingency, c_monitored)` represents the fraction of the pre-contingency flow on connection `c_contingency` that will flow on `c_monitored` if `c_contingency` is disconnected. Therefore, the post contingency flow on connection `c_monitored` is the `pre_contingency_flow` plus `lodf(c_contingency, c_monitored)*pre_contingency_flow(c_contingency)`. Therefore, consideration of `N` contingencies on `M` monitored lines introduces `N x M` constraints into the model. Usually one wishes to contain this number and methods are given below to achieve this.
3. **Defining your network** To identify the network for which ptdfs, lodfs and connection\_flows will be calculated according to the ptdf method, one does the following:
  - Create [node](#) objects for each bus in the model.
  - Create [connection](#) objects representing each line of the network: For each connection specify the `connection_reactance` parameter and the `connection_type` parameter. Setting `connection_type=connection_type_lossless_bidirectional` simplifies the amount of data that needs to be specified for an electrical network. See [connection\\_type](#) for more details
  - Set the `connection_to_node` and `connection_from_node` relationships to define the topology of each connection along with the `connection_capacity` parameter on one or both of these relationships.
  - Set the `connection_emergency_capacity` parameter to define the post contingency rating if lodf-based N-1 security constraints are to be included

- Create a `commodity` object and `node__commodity` relationships for all the nodes that comprise the electrical network for which PTDFs are to be calculated.
- Specify the `commodity_physics` parameter for the commodity to `:commodity_physics_ptdf` if ptdf-based DC load flow is desired with no N-1 security constraints or to `:commodity_physics_lodf` if it is desired to include lodf-based N-1 security constraints
- To identify the reference bus(`node`) specify the `node_opf_type` parameter for the appropriate `node` with the value `node_opf_type_reference`.

#### 4. Controlling problem size

- The lines to be monitored are specified by setting the `connection_monitored` property for each connection for which a flow constraint is to be generated
- The contingencies to be considered are specified by setting the `connection_contingency` property for the appropriate connections. For N contingencies and M monitored lines, N x M constraints will be generated.
- If the `lodf(c_contingency, c_monitored)` is very small, it means the outage of `c_contingency` has a small impact on the flow on `c_monitored` and there is little point in including this constraint in the model. This can be achieved by setting the `commodity_lodf_tolerance` `commodity` parameter. Contingency / Monitored line combinations with lodfs below this value will be ignored, reducing the size of the model.
- If `ptdf(n, c)` is very small, it means an injection at n has a small impact on the flow on c and there is little point in considering it. This can be achieved by setting the `commodity_ptdf_threshold` `commodity` parameter. Node / Monitored line combinations with ptdfs below this value will be ignored, reducing the number of coefficients in the model.

# Pressure driven gas transfer

The generic formulation of SpineOpt is based on a trade based model. However, network physics can be different depending on the traded commodity. This chapter specifically addresses the use of pressure driven gas transfer models and enabling linepack flexibility in SpineOpt. To this date, investments in pressure driven pipelines are not yet supported within SpineOpt. The use of multiple feed-in nodes, e.g. to represent multiple commodity flows through a pipeline is not yet supported.

For the representation of pressure driven gas transfer, we use the MILP formulation, as described in [Schwele - Coordination of Power and Natural Gas Systems: Convexification Approaches for Linepack Modeling](#). Here, the non-linearities associated with the Weymouth equation are convexified through an outer approximation of the Weymouth equation through fixed pressure points.

## Key concept

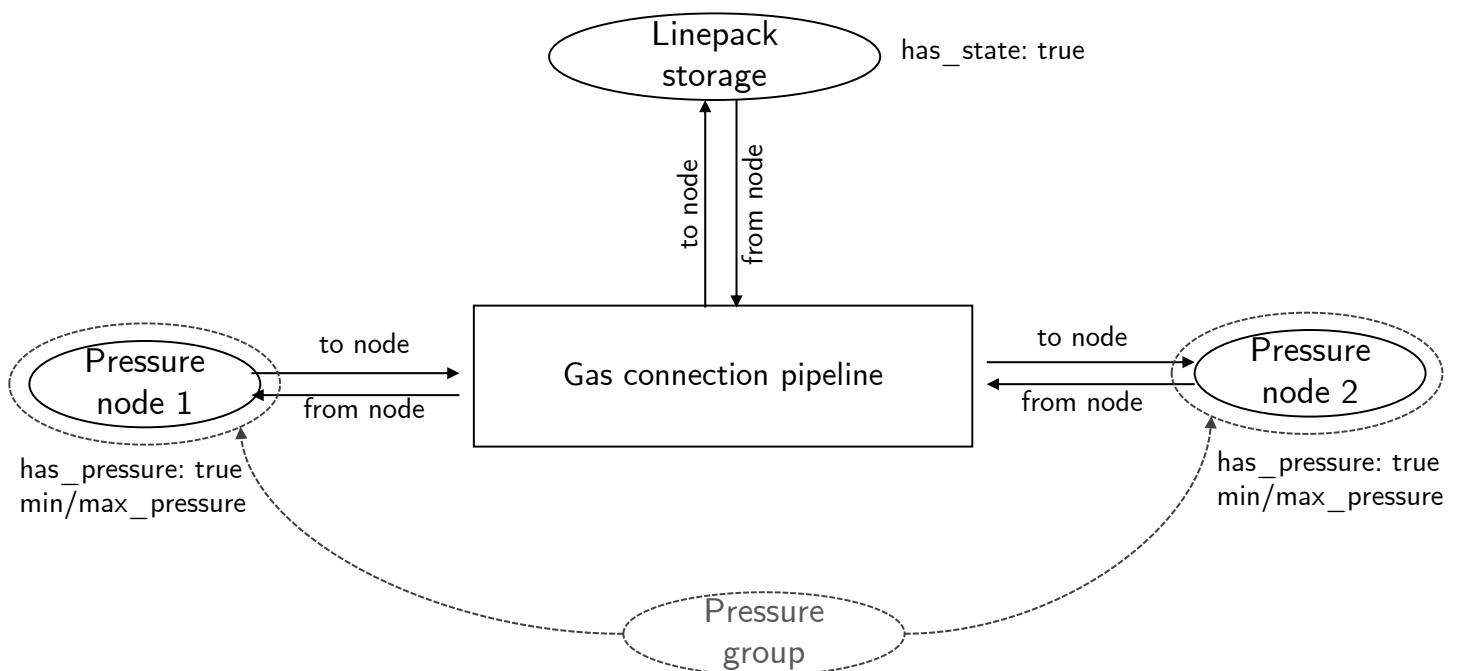
Here, we briefly describe the key objects and relationships required to model pressure driven gas transfers in SpineOpt.

1. **connection**: A connection represents the gas pipeline being modelled. Usually the direction of flow is not known a priori. To ensure that the flow through the gas pipeline is unidirectional, the parameter `has_binary_gas_flow` needs to be set to `true`.
2. **node**: Nodes with different characteristics are used for the representation of pressure driven gas transfer.
  - For each connection, there will be two nodes representing the start and end point of the pipeline. Associated with these nodes are the following parameters: the `has_pressure` parameter, which needs to be set to `true`, in order to create the variable `node_pressure`; the `max_node_pressure` and `min_node_pressure` to constrain the pressure variable.
  - To leverage linepack flexibility, a third node is introduced representing the linepack storage of the pipeline. To trigger the storage linepack and hence, `node_state` variables, the `has_state` parameter needs to be set to `true`.
3. **connection\_to\_node** and **connection\_from\_node** To enable flows through the pipeline and into the linepack storage, each `node` has to have both these relationships in common with the connection pipeline. These relationships will trigger the generation of `connection_flow` variables in all possible directions.
4. **connection\_node\_node** This relationship is key to the pressure driven gas transfer, holding the information about the pipeline characteristics and bringing the elements into interaction.
  - The parameter `connection_linepack_constant` holds the linepack constant and triggers the generation of the `line pack storage constraint`. Note that the first node should be the linepack

storage node, while the second node should be a [node\\_group](#) of both, the start and the end node of the pipeline.

- The linearization of the Weymouth equation through outer approximation relies on the use of fixed pressure points. For this purpose, the two parameters [fixed\\_pressure\\_constant\\_1](#) and [fixed\\_pressure\\_constant\\_0](#) hold the fixed pressure constants and trigger the generation of the [constraint\\_fix\\_node\\_pressure\\_point](#). The constraint introduces the relationship between pressure and gas flows. Note, that the pressure constants should be entered in a way, that the first node represents the origin node, the second node the destination node. Each connection should have a [connection\\_node\\_node](#) to each combination of its start and end nodes (and associated parameters). (See [Schwele - Coordination of Power and Natural Gas Systems: Convexification Approaches for Linepack Modeling](#))
- By default, pipelines are considered to be passive. However, a compression station between two pipeline pressure nodes can be represented by defining a [compression\\_factor](#). The relationship should be defined in such a manner, that the first node represents the sending node, the second node represents the receiving node, which pressure is equal or smaller to the pressure at the sending node times the compression factor.
- Lastly, to ensure the balance between incoming/outgoing flows and flows into the linepack, the ratio between the flows need to be fixed. The average incoming flows of the node group (of the pressure start and end nodes) have to equal the flows into the linepack storage, and vice versa. Therefore, the [fix\\_ratio\\_out\\_in\\_connection\\_flow](#) needs to be set to a value (typically 1) for the (pressure group, linepack storage) node pair, and for the (linepack storage, pressure group) node pair.

A gas pipeline and its connected nodes are illustrated below. A complete mathematical formulation can be found [here](#).



# Lossless nodal DC power flows

Currently, there are two different methods to represent lossless DC power flows. In the following the implementation of the nodal model is presented, based of node voltage angles.

## Key concepts

In the following, it is described how to set up a connection in order to represent a nodal lossless DC power flow network. Therefore, key object - and relationship classes as well as parameters are introduced.

1. **connection**: A connection represents the electricity line being modelled. A physical property of a connection is its [connection\\_reactance](#), which is defined on the connection object. Furthermore, if the reactance is given in a p.u. different from the standard unit used (e.g. p.u. = 100MVA), the parameter [connection\\_reactance\\_base](#) can be used to perform this conversion.
2. **node**: In a lossless DC power flow model, nodes correspond to buses. To use voltage angles for the representation of a lossless DC model, the [has\\_voltage\\_angle](#) needs to be `true` for these nodes (which will trigger the generation of the [node\\_voltage\\_angle](#) variable). Limits on the voltage angle can be enforced through the [max\\_voltage\\_angle](#) and [min\\_voltage\\_angle](#) parameters. The reference node of the system should have a voltage angle equal to zero, assigned through the parameter [fix\\_node\\_voltage\\_angle](#).
3. **connection\_to\_node** and **connection\_from\_node**: These relationships need to be introduced between the connection and each node, in order to allow power flows (i.e. [connection\\_flow](#)). Furthermore, a capacity limit on the connection line can be introduced on these relationships through the parameter [connection\\_capacity](#).
4. **connection\_node\_node**: To ensure energy conservation across the power line, a fixed ratio between incoming and outgoing flows should be given. The [fix\\_ratio\\_out\\_in\\_connection\\_flow](#) parameter enforces a fixed ratio between outgoing flows (i.e. to\_node) and incoming flows (i.e. from\_node). This parameter should be defined for both flow direction.

The mathematical formulation of the lossless DC power flow model using voltage angles is fully described [here](#).

# Representative days with seasonal storages

In order to reduce computational times, representative periods are often used in optimization models. However, this often limits the ability to properly account for seasonal storages.

In SpineOpt, we provide functionality to use representative days with seasonal storages in combination with the package [SpinePeriods.jl](#).

## General idea

The general idea is to mimick the seasonal effects throughout a non-representative period, e.g. a year of optimization, by introducing a specific sequence of the representative periods. Taking the example of one year to be optimized with representative days and seasonal storages, [SpinePeriods.jl](#) provides a mapping of each day of the year to its corresponding representative day. This information is stored in the mapping parameter [representative\\_periods\\_mapping](#) and is defined on the [temporal\\_block](#) for the whole year. The [representative\\_periods\\_mapping](#) parameter is a timeseries, pointing the beginning of each day to its corresponding representative day [temporal\\_block](#), which can also be automatically be generated through [SpinePeriods.jl](#).

In SpineOpt, this is interpreted in the following way:

- All operational variables, with the exception of the [node\\_state](#) variable, are created for each representative period. For each non-representative period, the variables are mapped to their corresponding variable of the representative periods according to the [representative\\_periods\\_mapping](#) parameter.
- Only the [node\\_state](#) variables and all investment [variables](#) are created for both, representative and non-representative period (of course, depending on the existence of relationships to [temporal\\_blocks](#)).

## Usage of representative days and seasonal storages for investment problems

To make use of representative days with seasonal storages concept, multiple [temporal\\_block](#) objects need to be created and connected to the system components, holding information about the resolutions in different parts of the model that come into play. As described in the section [Temporal Framework](#), every temporal block needs to be connected to a [model](#) object.

- **[temporal\\_block](#) for investments:** In order to define the resolution of the investment decisions, a [temporal\\_block](#) reflecting the frequency of investment decisions should be introduced. For yearly investment, the [resolution](#) of this temporal block would be equal to 1Y. In order to link [nodes](#), [units](#)

or `connections` to this investment resolution, the `node_investment_temporal_block`, `unit_investment_temporal_block`, or `connection_investment_temporal_block` relationships need to be defined, respectively. For more details on investments, see also section [Investment Optimization](#)

- **temporal\_block for representative days:** For each representative day, one `temporal_block` needs to be created, indicating the `block_start` and `block_end` of the representative day. The use of disconnected periods is also described in the section [Disconnected time periods](#). The resolution of the representative days corresponds to the resolution of the operational variables, e.g. 1h. In order to associate operational variables with the representative periods, `node_temporal_block` and `units_on_temporal_block` relationships need to be created. For convenience, it is also possible to create a group of all representative `temporal_blocks` and link this group to these relationships. Note that, when using `SpinePeriods.jl`, the representative temporal blocks are auto-generated.
- **temporal\_block for non-representative days:** To introduce `node_state` variables for the entire operational period, a temporal block overarching the entire horizon is created. Note that currently, this temporal block needs to have the same resolution as the representative days, e.g. 1h. In order to associate operational variables with the representative periods, `node_temporal_block` and `units_on_temporal_block` relationships need to be created. Note that, as described above, the non-representative variables, will be mapped to their corresponding representative days. To manually introduce the mapping between non-representative and representative periods, instead of using the recommended `SpinePeriods.jl`, the user must define the mapping parameter `representative_periods_mapping` by hand, consisting of `DateTime` indices (indicating the start of each non-representative period, e.g. for a daily mapping 2021-01-01T00:00:00, 2021-01-02T00:00:00 etc.) and the name of the corresponding representative `temporal_block` as a value.